

66

6. Standard declarations

This section describes the ordered set SD of standard declarations supporting the use of the language. The description takes the form of orderly enumeration of a family of disjoint subsets in SD .[†] The description of each member of this family is headed by a reference numbering of the form $(SDi.j.k)$. The ordering in SD is such that one element under $(SDi_1.j_1.k_1)$ precedes another element under $(SDi_2.j_2.k_2)$ if $(i_1.j_1.k_1)$ precedes $(i_2.j_2.k_2)$ in the lexical order.

Each member description begins either with a '{' (not a '{') or with an 'f'.

(a) If a member description begins with a '{', then the matching '}' terminates that description, and the member has a single element, which is the enclosed text.

(b) If a member description begins with an 'f', then the part enclosed by the first following '{' and the matching '}', with which the member description is terminated, gives an element of the member when asterisked symbols (T^* for example) in that part are consistently replaced by certain objects as prescribed after the leading 'f' in terms of common mathematical notations ($T^*e T$ for example).

[†] In the present document, instead of the whole family covering SD , only some illustrating subfamily will be presented. Other members, including those for input/output operations, will be found elsewhere.

6.1. Basic operations

6.1.1. Copying and enproceduring operations

(SD1.1.1) {let copy operate before all left after copy right}(SD1.1.2) {let enproc operate before all left after all right}(SD1.1.3) $\int T^* \in T$ {let copy() represent procedure(T*)T* by
((original) code (c:original, d:T*) T* by
(core(Q) let w(Q)={<c:,Q₁>, <d:,Q₂>};
m(Q₂) \leftarrow m(Q₁); w(Q₂) \leftarrow w(Q₁); \Rightarrow Q₂ end of core)))}(SD1.1.4) $\int T^* \in T$ {let enproc() represent
procedure(T*)procedure()T* by
((body) procedure()T* by
(() body))}

6.1.2. Simple assignment operations

(SD1.2.1) {let \leftarrow , := operate before \leftarrow , := left after all right}(SD1.2.2) $\int T^* \in \{\text{real, bits, } \overset{\text{string}}{\text{reference}}\} \cup T$ procedure, $\Delta^* \in \{\leftarrow, :=\}$
{let () Δ^* () represent procedure(T*, T*)effect by
((destination, source)
code (d:destination, s:source) effect by
(core(Q) let w(Q)={<d:,Q₁>, <s:,Q₂>};
w(Q₁) \leftarrow p(m(Q₁), w(Q₂)); \Rightarrow Q₀ end of core))}

6.1.3. Simple comparison operations

(SD1.3.1) {let =, \neq operate before left after \leftarrow , := right}(SD1.3.2) {let <, \leq , \geq , > operate before =, \neq left
after \leftarrow , := right}(SD1.3.3) $\int T^* \in T, \Delta^* \in \{=, \neq\}; T^* = \text{real}, \Delta^* \in \{<, \leq, \geq, >\}$
{let () Δ^* () represent procedure(T*, T*)bits by
((left, right) code (l:copy left, r:right, t:l) bits by
(core(Q) let w(Q)={<l:,Q₁>, <r:,Q₂>, <t:,Q₃>};
if w(Q₁) Δ^* w(Q₂) then \Rightarrow Q₃, else \rightarrow next;
w(Q₃) \leftarrow 0; \Rightarrow Q₃ end of core))}

6.1.4. Conditional operations

- (SD1.4.1) {let if operate before all left}
- (SD1.4.2) {let then operate}
- (SD1.4.3) {let else, do operate after all right}
- (SD1.4.4) $\int T^*eT$ {let if()then()else() represent
procedure(bits, T*, T*) T* by
 ((condition, then, else)
 code (c:condition, t:enproc then, e:enproc else)
enproc T* by
 (core(Q) let w(Q)={<c:,Q₁>,<t:,Q₂>,<e:,Q₃>};
 if the bit string w(Q₁) contains at least one 1
 then $\Rightarrow Q_2$, else $\Rightarrow Q_3$ end of core) ())}
- (SD1.4.5) {let if()do() represent procedure(bits, effect)effect
by((condition, statement) if condition then statement
else dummy)}

6.1.5. Basic arithmetic operations

- (SD1.5.1) {let +,~ operate before =,≠,<,≤,≥,> left
after +, := right}
- (SD1.5.2) {let ×,/ operate before =,≠,<,≤,≥,>,+,- left
after +, := right}
- (SD1.5.3) $\int \Delta^*e\{+,-,\times,/ \}$ {let () Δ^* () represent
procedure(real,real)real by ((left,right)
code (a:copy left, b:right) real by
 (core(Q) let w(Q)={<a:,Q₁>,<b:,Q₂>};
 if the arithmetic operation meant by w(Q₁) Δ^* w(Q₂)
 can be performed then \rightarrow next, else $\Rightarrow L0$;
 let W₃ be the resulting value of that operation
 (possibly with some implementation dependent deviation);
 w(Q₁) \leftarrow p(m(Q₁),W₃); $\Rightarrow Q_1$ end of core)}
- (SD1.5.4) {let -() represent procedure(real)real by
 ((right) begin (right)b; (copy b)a; a:=0; a-b end)}
- (SD1.5.5) {let /() represent procedure(real)real by
 ((right) begin (right)b; (copy b)a; a:=1; a/b end)}

6.1.6. Some enquiry operations

- (SD1.6.1) {let mode, length, bd operate before all left
after right}
- (SD1.6.2) {let mode() represent procedure(real)structure
(fix:bits, min:real, step:real, max:real) by ((real)
begin (real)x; code (fix:l, min:copy x, step:real,
max:copy x) (mode real) by
(core(Q) let w(Q)={<fix:,Q₁>, <min:,Q₂>, <step:,Q₃>, <max:,Q₄>};
let M stand for m(Q₂);
w(Q₂)← the minimum value in W_M;
w(Q₄)← the maximum value in W_M;
if M=real [precision R] with a real number R,
then →next, else →Kfixed;
w(Q₃)← p(m(Q₃), R); w(Q₁)← 0; ⇒Q;
Kfixed: let M=real [R₁:R₂:R₃] with real numbers R₁,
R₂, R₃;
w(Q₃)← p(m(Q₃), R₂); ⇒Q end of core) end)}}
- (SD1.6.3) {T*e{bits,string} {let mode() represent
procedure(T*)structure (exact:bits, length:real) by ((some)
code (s:some, t:(exact:l, length:)) (mode T*) by
(core(Q) let w(Q)={<s:,Q₁>, <t:,Q₂>}; → integer
let w(Q₂)={<exact:,Q₃>, <length:,Q₄>};
let m(Q₁)=T* [Y I] ,where I is an integer and Y
is either exact or varying ;
w(Q₄)← p(m(Q₄), I);
if Y=exact then ⇒Q₂, else →next;
w(Q₃)← 0; ⇒Q₂ end of core) } }}
- (SD1.6.4) {T*eT {let mode() represent procedure(array T*)
structure (lbd:real, ubd:real) by ((array)
code (a:array, t:(lbd:l, ubd:)) (mode array T*) by
(core(Q) let w(Q)={<a:,Q₁>, <t:,Q₂>}; → integer
let w(Q₂)={<lbd:,Q₃>, <ubd:,Q₄>};
let m(Q₁)=array [I₁:I₂] T* with integers I₁, I₂;
w(Q₃)← p(m(Q₃), I₁); w(Q₄)← p(m(Q₄), I₂);
⇒Q₂ end of core) } }}

6.1.6 continued

(SD1.6.5) $\int T^*e\{\text{bits}, \text{string}\}$ {let length() represent
procedure(T*)real by ((string) code (s:string, t:) real by
 (core(Q) let w(Q)={<s:,Q₁>, <t:,Q₂>};
 let I be the integral length of w(Q₁);
 w(Q₂)← p(m(Q₂),I); \Rightarrow Q₂ end of core)}

integer

(SD1.6.6) $\int T^*eT$ {let bd() represent procedure(array T*)real
by ((array) (mode array)[lbd:])}

(SD1.6.7) $\int T^*eT$ {let ()bd represent procedure(array T*)real
by ((array) (mode array)[ubd:])}

6.2. Extended operations

6.2.1. Repetitive operations

- (SD2.1.1) {let while,until operate before left after all right}
- (SD2.1.2) {let succ,step operate before left after $\leftarrow, :=, =, \neq, <, \leq, \geq, >, +, -, \times, /$ right}
- (SD2.1.3) $\int T^* \epsilon T$ {let ()succ() represent procedure(T^*, T^*)
structure (init:enproc T^* , succ:enproc T^*) by
((a, s) (init:enproc a , succ:enproc s)) }
- (SD2.1.4) $\int T^* \epsilon T$ {let ()while() represent procedure($T^* \text{succ} T^*, \text{bits}$)
structure (init:enproc T^* , succ:enproc T^* , while:enproc bits) by
((r, t) begin (r rl; (init:copy $rl[\text{init:}]$,
succ:copy $rl[\text{succ:}]$, while:enproc t) end) }
- (SD2.1.5) {let ()step() represent procedure(real, real)
structure (init:enproc real , step:enproc real) by
((a, b) (init:enproc a , step:enproc b)) }
- (SD2.1.6) {let step() represent procedure(real)(real step real)
by ((b) 1 step b) }
- (SD2.1.7) {let ()until() represent procedure(real step real, real)
structure (init:enproc real , step:enproc real ,
until:enproc real) by
(progression, limit)
begin (progression) ab; (init:copy $ab[\text{init:}]$,
step:copy $ab[\text{step:}]$, until:enproc limit) end}
- (SD2.1.8) {let ()until() represent procedure(real, real)
(step real until real) by ((a, c) a step 1 until c) }
- (SD2.1.9) {let until() represent procedure (real)(real until real)
by ((c) 1 until c) }

6.2.1 continued

(SD2.1.10) {let from operate before all left}

(SD2.1.11) $\int T^*eT$ {let ()from()do^()represent
procedure(T*, T*succT*while bits, effect) effect by
 ((cvar, domain, statement)
begin (cvar)cv; (domain)dom; cv← dom[init:]();
 next: if dom[while:]() do
begin statement; cv← dom[succ:](); go to next end
end)}

(SD2.1.12) {let ()from()do() represent
procedure(real, until real, effect) effect by
 ((cvar, domain, statement)
begin (cvar)cv; (copy cv)v1; (copy domain)dom;
 (copy dom[init:]())a1; (copy dom[step:]())b1;
 (copy dom[until:]())c1;
 v1 from a1 succ v1+b1 while if b1>0 then v1<=c1
else if b1<0 then v1>=c1 else v1≠c1
do begin cv:=v1; statement end
end)}

6.2.2. Assign-and-hold operations

(SD2.2.1) {let the operate before all left after right}

(SD2.2.2) $\int T^*eT$ {let the() represent
procedure(T*) procedure(effect) T* by
 ((expression) procedure(effect) T* by ((dummy) expression))}

(SD2.2.3) $\int T^*eT, \Delta^*e\{\leftarrow, :=\}$ {let () Δ^* () represent
procedure(the T*, T*) T* by ((the destination, source)
begin (the destination(dummy))destination;
 destination Δ^* source; destination
end)}

6.2.3. Reference handling operations

(SD2.3.1) {let ref, has type, as type operate before all left
after right }

(SD2.3.2) $\int T^* \in T$ {let ref() represent procedure(T*) reference
by ((referent) code (q:referent, r:reference) reference
by (core(Q) let w(Q)={<q:,Q₁>,<r:,Q₂>} ;
w(Q₂)← {Q₁}; \Rightarrow Q₂ end of core))}

(SD2.3.3) $\int T^* \in T$ {let () has type() represent
procedure(reference, T*) bits by ((ref,type)
code (r:ref, t:type, b:0) bits by
(core(Q) let w(Q)={<r:,Q₁>,<t:,Q₂>,<b:,Q₃>};
if w(Q₁)= empty, then \Rightarrow Q₃, else \rightarrow next;
let w(Q₁)={Q₄};
if t(Q₂)=t(Q₄) then \rightarrow next, else \Rightarrow Q₃;
w(Q₃)← 1; \Rightarrow Q₃ end of core))}

(SD2.3.4) $\int T^* \in T, T^* \neq \text{reference}$ {let () as type() represent
procedure(reference, T*) T* by ((ref, type)
code (r:ref, t:type) T* by
(core(Q) let w(Q)={<r:,Q₁>,<t:,Q₂>};
if w(Q₁)=empty, then \Rightarrow Q₂, else \rightarrow next;
let w(Q₁)={Q₄};
if t(Q₂)=t(Q₄) then \Rightarrow Q₄, else \Rightarrow Q₂ end of core))}