

Geometrical Methods for Adaptive Approximation of Image and Video Data

Laurent Demaret and <u>Armin Iske</u>

based on joint work with $\ensuremath{\mathrm{NIRA}}$ Dyn and $\ensuremath{\mathrm{WAHID}}$ Khachabi



1 Introduction

Digital Image Compression: Basic Steps

- (1) Data reduction from input image;
- (2) Encoding of the reduced data at the sender;
- (3) Transmission of the encoded data from the sender to the receiver;
- (4) Decoding of the transmitted data at the receiver;
- (5) Data reconstruction.



Original Image.

0101100011010110010 ...





Reconstruction.



Image Representation.

- A digital image I is a rectangular grid of pixels, X.
- Each pixel $x \in X$ bears a greyscale luminance I(x).
- We regard the image as a function, $I:[X] \to [0,1,\ldots,2^r-1]$, where the convex hull [X] of X is the image domain.





Image Approximation.

INPUT: The image $I = \{(x, I(x)) : x \in X\}$ is given by discrete pixel values in X. **OUTPUT:** Reconstructed image $\tilde{I} = \{(x, \tilde{I}(x)) : x \in X\}$.



AIM. Increase *Peak Signal to Noise Ratio* (PSNR)

$$\mathsf{PSNR} = 10 * \mathsf{log}_{10} \left(rac{2^{\mathrm{r}} imes 2^{\mathrm{r}}}{ar\eta^2(\mathrm{I}, \widetilde{\mathrm{I}})}
ight),$$

as much as possible, where

$$\bar{\eta}^2(I,\tilde{I}) = \frac{1}{|X|} \sum_{x \in X} |I(x) - \tilde{I}(x)|^2$$

denotes the *mean square error* (MSE).



2 Methods for Image Compression

Wavelets: The standard (EBCOT, JPEG2000)

Wavelet Image Approximation Scheme.

- The image is expanded in a fixed orthonormal basis of wavelets.
- The expansion coefficients below a given threshold are set to zero.

A mildly nonlinear approximation scheme.



Some recent highly nonlinear approximation schemes ...

... for capturing the image geometry.

- Bandelets: LEPENNEC & MALLAT (2005);
- Brushlets: COIFMAN & MEYER (1997);
- Curvelets: CANDÈS & DONOHO (2000, 2004/2005);
- **Contourlets:** Do & VETTERLI (2005);
- Directionlets: Velisavljević, Beferull-Lozano, Vetterli & Dragotti (2006);
- Shearlets: GUO, KUTYNIOK, LABATE, LIM (2006);
- Wedgelets: DONOHO (1999); ROMBERG, WAKIN & BARANIUK (2002);
- The Easy Path Wavelet Transform (EPWT): PLONKA(2009), PLONKA, TENORTH & I.(2010), PLONKA, TENORTH & ROŞCA (2009);
- Nonlinear edge-adapted multiscale decomposition: COHEN & MATEI (2001);
- Adaptive approximation by anisotropic triangulations:
 - Generic triangulations and simulated annealing: LEHNER, UMLAUF, HAMANN (2007)
 - Adaptive thinning algorithms: DEMARET, DYN & I. (2006), DEMARET & I. (2006)
 - Anisotropic geodesic triangulations: BOUGLEUX, PEYRÉ & L. COHEN (2009)
 - Greedy triangle bisections: A. COHEN, DYN, HECHT & MIREBEAU (2010)

3 Linear Splines over Triangulations

Definition. A triangulation of a planar point set $Y = \{y_1, \ldots, y_N\}$ is a collection $\mathcal{T}(Y) = \{T\}_{T \in \mathcal{T}(Y)}$ of triangles in the plane, such that

(T1) the vertex set of $\mathcal{T}(Y)$ is Y;

- (T2) any pair of two distinct triangles in $\mathcal{T}(Y)$ intersect at most at one common vertex or along one common edge;
- (T3) the convex hull [Y] of Y coincides with the area covered by the union of the triangles in $\mathcal{T}(Y)$.







Linear Splines over Triangulations.



Triangulation of pixels.



Linear spline over triangulation.



Approximation Spaces.

 \bullet Given any triangulation $\mathcal{T}(Y)$ of Y, we denote by

$$\mathcal{S}_{Y} = \left\{s: s \in C([Y]) \text{ and } s \big|_{T} \text{ linear for all } T \in \mathcal{T}(Y) \right\},$$

the spline space containing all continuous functions over [Y] whose restriction to any triangle in $\mathcal{T}(Y)$ is linear.

- Any element in \mathcal{S}_Y is referred to as a linear spline over $\mathcal{T}(Y)$.
- For given function values $\{I(y): y \in Y\}$, there is a unique linear spline, $L(Y, I) \in S_Y$, which interpolates I at the points of Y, i.e.,

$$L(Y,I)(y)=I(y), \qquad \text{ for all } y\in Y.$$



Example 1: Geometrical Image PQuad.



Image PQuad of size (512×512) .



Adaptive Triangulation with 800 vertices.



Reconstruction at PSNR 42.85 db.



Example 2: Geometrical Image Game.



Image Game of size (512×512) .



Adaptive Triangulation with 6000 vertices.



Reconstruction at PSNR 36.54 db.



Example 3: Multiscale Image Aerial.



Image Aerial of size (512×512) .



Adaptive Triangulation with 16000 vertices.



Reconstruction at PSNR 30.33 db.



Example 4: Multiscale Image Boat.



Image Boat of size (512×512) .



Adaptive Triangulation with 7000 vertices.



Reconstruction at PSNR 31.83 db.

4 Approximation over Anisotropic Triangulations

Goal: On input image $I = \{(x, I(x)) : x \in X\}$,

- determine a *good* adaptive spline space S_Y , where $Y \subset X$;
- determine from \mathcal{S}_Y the unique best approximation $L^*(Y\!,I)\in \mathcal{S}_Y$ satisfying

$$\sum_{x \in X} |L^*(Y, I)(x) - I(x)|^2 = \min_{s \in S_Y} \sum_{x \in X} |s(x) - I(x)|^2.$$

- Encode the linear spline $L^* \in \mathcal{S}_Y$;
- Decode $L^* \in S_Y$, and so obtain the reconstructed image $\tilde{I} = \{(x, L(Y, \tilde{I})(x)) : x \in X\}$, where $L(Y, \tilde{I}) \approx L^*(Y, I)$.

OBS! Key Step: Construction of *anisotropic* triangulation $\mathcal{T}(Y)$ for $Y \subset X$.

- One possible approach is by *adaptive thinning* (AT).
- In AT, we take the Delaunay triangulation $\mathcal{D}(Y)$ of Y for $\mathcal{S}_Y,$

UH



The Bramble-Hilbert Lemma.

Recall classical error estimates from finite element methods (FEM).

Bramble-Hilbert: For any image f from Sobolev space $W^{2,2}(T)$, $T \in \mathcal{T}(Y)$, we obtain the basic error estimate

$$\|f - \Pi_{\mathcal{S}_{Y}} f\|_{L^{2}(T)} \le |f|_{W^{2,2}(T)}, \quad \text{for } f \in W^{2,2}(T),$$

where $\Pi_{S_Y} f$ is the orthogonal L²-projection of f onto S_Y .



5 Delaunay Triangulations

Definition. The Delaunay triangulation $\mathcal{D}(X)$ of a discrete planar point set X is a triangulation of X, such that the circumcircle for each of its triangles does not contain any point from X in its interior.



Two triangulations of a convex quadrilateral, \mathcal{T} (left) and $\tilde{\mathcal{T}}$ (right).

Properties of Delaunay Triangulations.

• Uniqueness.

Delaunay triangulation $\mathcal{D}(X)$ is *unique*, if no four points in X are co-circular.

• Complexity.

For any point set X, its Delaunay triangulation $\mathcal{D}(X)$ can be computed in $\mathcal{O}(N \log N)$ steps, where N = |X|.

• Local Updating.

For any X and $x \in X$, the Delaunay triangulation $\mathcal{D}(X \setminus x)$ of the point set $X \setminus x$ can be computed from $\mathcal{D}(X)$ by retriangulating the *cell* $\mathcal{C}(x)$ of x.



Removal of the node y, and retriangulation of its cell $\mathcal{C}(y)$.

UН



6 Adaptive Thinning

Popular Example: Test Image Fruits.



Original Image (512×512) .



Delaunay Triangulation.



4044 significant pixels.



Image Reconstruction.



Adaptive Thinning Algorithm.

INPUT. I = $\{0, 1, ..., 2^r - 1\}^X$, pixels and luminances, where X set of pixels, r number of bits for representation of luminances.

- (1) Let $X_N = X$;
- (2) FOR k = 1, ..., N n

(2a) Find a least significant pixel $x \in X_{N-k+1}$;

- (2b) Let $X_{N-k} = X_{N-k+1} \setminus x$;
- **OUTPUT:** Data hierarchy

$$X_n \subset X_{n+1} \subset \cdots \subset X_{N-1} \subset X_N = X$$

of nested subsets of X.



Controlling the Mean Square Error.

- For a given mean square error (MSE), $\bar{\eta}^*$, the adaptive thinning algorithm can be changed in order to terminate when for the first time, the MSE value corresponding to the current linear spline $L(X_p, I)$ is above $\bar{\eta}^*$, for some X_p in the data hierarchy, n = p a posteriori.
- We take as the final approximation to the image the linear spline $L^*(X_{p+1}, I)$, and so we let $Y = X_{p+1}$.
- \bullet Observe that $L^*(X_{p+1}, I)$ satisfies

$$\sum_{x \in X} |L^*(X_{p+1}, I)(x) - I(x)|^2 / |X_{p+1}| \le \bar{\eta}^*,$$

as desired.



7 Pixel Significance Measures

Quality Measure: Current ℓ_2 -Square Error.

$$\eta(Y;X) = \sum_{x \in X} |L(I,Y)(x) - I(x)|^2, \qquad \text{ for } Y \subset X.$$

Anticipated Error for the Greedy Removal of one Pixel.

$$e(y)=\eta(Y\setminus y;X),\qquad \text{ for }y\in Y.$$

Definition. (Adaptive Thinning Algorithm AT). For $Y \subset X$, a point $y^* \in Y$ is said to be least significant in Y, iff it satisfies

$$e(\mathbf{y}^*) = \min_{\mathbf{y}\in\mathbf{Y}} e(\mathbf{y}).$$



Aim: Compute anticipated error *locally*.

$$\begin{split} \eta(Y \setminus y; X) &= \eta(Y \setminus y; X \setminus \mathcal{C}(y)) + \eta(Y \setminus y; X \cap \mathcal{C}(y)) \\ &= \eta(Y; X \setminus \mathcal{C}(y)) + \eta(Y \setminus y; X \cap \mathcal{C}(y)) \\ &= \eta(Y; X) + \eta(Y \setminus y; X \cap \mathcal{C}(y)) - \eta(Y; X \cap \mathcal{C}(y)). \end{split}$$

where $\mathcal{C}(y)$ is the cell of y in the Delaunay triangulation $\mathcal{D}(Y)$ of Y.

Therefore, minimizing e(y) is equivalent to minimizing

$$e_{\delta}(\mathbf{y}) = \eta(\mathbf{Y} \setminus \mathbf{y}; \mathbf{X} \cap \mathcal{C}(\mathbf{y})) - \eta(\mathbf{Y}; \mathbf{X} \cap \mathcal{C}(\mathbf{y})), \quad \text{for } \mathbf{y} \in \mathbf{Y}.$$

Proposition. For $Y \subset X$, a point $y^* \in Y$ is, according to the criterion **AT**, **least significant** in Y, iff it satisfies

$$e_{\delta}(\mathbf{y}^*) = \min_{\mathbf{y}\in\mathbf{Y}} e_{\delta}(\mathbf{y}).$$



Greedy Two-Point-Removal.

Anticipated Error for the Removal of two Points.

 $e(y_1,y_2) = \eta(Y \setminus \{y_1,y_2\};X) \qquad \text{ for } y_1,y_2 \in Y,$

can be rewritten as $e(y_1,y_2)=\eta(Y;X)+e_{\delta}(y_1,y_2),$ where

 $e_{\delta}(y_1, y_2) = \eta(Y \setminus \{y_1, y_2\}; X \cap (\mathcal{C}(y_1) \cup \mathcal{C}(y_2))) - \eta(Y; X \cap (\mathcal{C}(y_1) \cup \mathcal{C}(y_2))),$

which can for $[y_1, y_2] \notin \mathcal{D}(Y)$ be simplified as

 $e_{\delta}(y_1, y_2) = e_{\delta}(y_1) + e_{\delta}(y_2).$

Definition. (Adaptive Thinning Algorithm AT²).

For $Y \subset X$, a point pair $y_1^*, y_2^* \in Y$ is said to be least significant in Y, iff

$$e_{\delta}(\mathbf{y}_1^*,\mathbf{y}_2^*) = \min_{\mathbf{y}_1,\mathbf{y}_2\in\mathbf{Y}} e_{\delta}(\mathbf{y}_1,\mathbf{y}_2).$$

8 Implementation of Adaptive Thinning. Efficient Implementation of Algorithm AT. Initialization.

- Compute Delaunay triangulation $\mathcal{D}(X)$;
- Compute $e_{\delta}(x)$ for all $x \in X$ and store nodes of $\mathcal{D}(X)$ in a Heap.

Removal Step. For current $Y \subset X$

- Pop root $y^* \in Y$ from Heap, update Heap;
- Remove y^* from $\mathcal{D}(Y)$ and compute $\mathcal{D}_{Y \setminus y^*}$;
- Reattach *historical points* in $C(y^*) \cap (X \setminus Y)$;
- Attach y^* to new triangle in $\mathcal{C}(y^*)$;
- Update $e_{\delta}(y)$ for neighbours of y^* and update Heap.

Total Complexity. $\mathcal{O}(N \log(N))$ operations.





Efficient Implementation of Algorithm AT².

• Due to the representation

 $e_{\delta}(y_1,y_2)=e_{\delta}(y_1)+e_{\delta}(y_2), \qquad \text{ for } [y_1,y_2] \notin \mathcal{D}(Y),$

the maintenance of significances $\{e_{\delta}(y_1, y_2): \{y_1, y_2\} \subset Y\}$ can be reduced to maintenance of $\{e_{\delta}(y_1, y_2): [y_1, y_2] \in \mathcal{D}(Y)\}$ and $\{e_{\delta}(y): y \in Y\}$.

- For efficient implementation of AT^2 we use two different priority queues,
 - HeapY for significances $e_{\delta}(y)$ of pixels $y \in Y$;
 - HeapE for significances $e_{\delta}(y_1; y_2)$ of edges $[y_1; y_2] \in \mathcal{D}(Y)$.
- Each priority queue, HeapY and HeapE, contains a least significant element at its head, and is updated after each pixel removal.
- The resulting algorithm has also complexity $\mathcal{O}(N \log N).$



Further Computational Details.

• We do not remove corner points from X, so that the image domain [X] is invariant during the performance of adaptive thinning.

Uniqueness of Delaunay triangulation.

- Recall that the Delaunay triangulation $\mathcal{D}(Y)$ of $Y \subset X$, is unique, provided that no four points in Y are co-circular.
- Since neither X nor its subsets satisfy this condition, we apply an efficient method, termed *simulation of simplicity* (Edelsbrunner & Mücke, 1990), which assures uniqueness (by using lexicographical order of vertices).
- Unlike in previous perturbation methods, the simulation of simplicity method allows us to work with integer arithmetic rather than with floating point arithmetic.

9 Local Optimization by Exchange

Definition: For any $Y \subset X$, let $Z = X \setminus Y$. A point pair $(y, z) \in Y \times Z$ satisfying

 $\eta((Y\cup z)\setminus y;X) < \eta(Y;X)$

is said to be exchangeable. A subset $Y \subset X$ is said to be locally optimal in X, iff there is no exchangeable point pair $(y, z) \in Y \times Z$.

Algorithm (Exchange) INPUT: $Y \subset X$;

(1) Let $Z = X \setminus Y$;

(2) WHILE (Y not locally optimal in X)

(2a) Locate an exchangeable pair $(y, z) \in Y \times Z$;

(2b) Let $Y = (Y \setminus y) \cup z$ and $Z = (Z \setminus z) \cup y$;

OUTPUT: $Y \subset X$, locally optimal in X.

UH

Characterization of Exchangeable Point Pairs.

Let $Z=X\setminus Y,$ for any $Y\subset X,$ and recall

$$\eta(Y\setminus y;X)=\eta(Y;X)+e_{\delta}(y;Y), \qquad \text{ for } y\in Y,$$

where $e_{\delta}(y;Y) = \eta(Y \setminus y; X \cap \mathcal{C}(y;Y)) - \eta(Y; X \cap \mathcal{C}(y;Y)).$

Letting first $Y = Y \cup z$, and then y = z, this implies

$$\begin{split} \eta((Y \cup z) \setminus y; X) &= \eta(Y \cup z; X) + e_{\delta}(y; Y \cup z) \\ \eta(Y; X) &= \eta(Y \cup z; X) + e_{\delta}(z, Y \cup z). \end{split}$$

Therefore, $(y, z) \in Y \times Z$ are exchangeable, iff

$$e_{\delta}(z; Y \cup z) > e_{\delta}(y; Y \cup z),$$

which simplifies to

 $e_{\delta}(z; Y \cup z) > e_{\delta}(y; Y),$

in case $C(y; Y) = C(y; Y \cup z)$, i.e., $[y; z] \notin D(Y \cup z)$.



Efficient Implementation of Exchange.

• Due to the swapping criterion

 $e_{\delta}(z; Y \cup z) > e_{\delta}(y; Y), \quad \text{for } [y; z] \notin \mathcal{D}(Y \cup z),$

the localization of exchangeable point pairs can efficiently be accomplished by maintenance of three different priority queues,

- HeapY for significances $e_{\delta}(y; Y)$ of pixels $y \in Y$;
- HeapZ for significances $e_{\delta}(z; Y \cup z)$ of pixels $z \in Z$;
- HeapE for significances $\sigma(y, z) = e_{\delta}(z; Y \cup z) e_{\delta}(y; Y \cup z)$ of edges $[y; z] \in \mathcal{D}(Y \cup z).$
- The priority queue HeapY contains a least significant element at its head; the head of HeapZ and HeapE contains a most significant element.
- Each of the three priority queues is updated after each pixel exchange.
- The resulting complexity for *one* pixel exchange is $\mathcal{O}(\log N)$.

UH



10 Image Compression

- Our compression method replaces the image I by its linear spline approximation $L^*(Y, I)$, where $Y \subset X$ are the significant pixels.
- $\bullet\,$ In order to code $L^*(Y\!,I),$ we code the information

 $\{(y, I^*(y)): y \in Y\}.$

Quantization.

- Apply uniform quantization to the *optimal* luminances $I^*(y) = L^*(Y, I)(y)$,
- so obtain quantized symbols $\{Q(I^*(y)): y \in Y\}$,
- corresponding to quantized luminances $\{\tilde{I}(y)): y \in Y\}$.



11 Theoretical Coding Costs

OBSERVE! Due to the uniqueness of the Delaunay triangulation, no connectivity coding is required!

• We are only concerned with coding the elements of the set

 $\{(y,Q(I^*(y))): y \in Y\} \in \mathcal{I}_n^s,$

where, with n = |Y|,

$$\mathcal{I}_n^s = \left\{\{0,1,\ldots,2^s-1\}^Z : Z \subset X \text{ and } |Z| = n\right\}.$$

- The number of elements in \mathcal{I}_n^s is $\binom{|X|}{n} \times 2^{s \times n}$.
- If we assume that every element of \mathcal{I}_n^s has the same probability of occurrence, then the theoretical coding cost is

$$\log_2\left(\binom{|X|}{n}\right) + s \times n.$$



12 Scattered Data Coding

OBSERVE! We can reduce the theoretical coding costs by taking advantage of the *geometric* structure of the image as follows.

The elements of $\{(i,j,Q(I^*(i,j))):(i,j)\in Y\}$ are coded by decomposing their bounding cell

 $\Omega = [0..2^p - 1] \times [0..2^q - 1] \times [0..2^s - 1]$

recursively, where $[0..2^{s} - 1]$ is the range for the quantized symbols.



Splitting of the cell Ω into eight subcells in three stages.



(1) Coding of Scattered Pixels.

- Coding of pixels in Y relies on a recursive splitting of the pixel domain $\Omega = [X].$
- For the sake of simplicity, let us assume that Ω is a square domain of the form $\Omega = [0, 2^q 1] \times [0, 2^q 1]$.
- In the splitting, a square subdomain ω ⊂ Ω (initially ω = Ω) is split horizontally into two rectangular subdomains of equal size. A rectangular subdomain is split vertically into two square subdomains of equal size.
- The splitting terminates at subdomains which are either *empty*, i.e., not containing any pixel from Y, or *atomic*, i.e., of size 1 × 1.



(1) Coding of Scattered Pixels.

- This recursive splitting can be represented by a binary tree, whose nodes correspond to the subdomains. The root of the tree corresponds to Ω, and its leaves correspond to empty or atomic subdomains.
- In each node of the tree, with a corresponding subdomain ω, we store the number |ω| of pixels from Y contained in ω, i.e., |ω| = |Y ∩ ω|.
- Note that for a parent node ω , and its two children nodes, ω_1 and ω_2 , we have the relation $|\omega| = |\omega_1| + |\omega_2|$. This relation allows a non-redundant representation of the binary tree.
- The bitstream, representing the tree, is constructed by a Huffman code.



(2) Coding of Quantized Symbols.

- To code the quantized symbols in Q_Y , we first split the image domain Ω into a small number of square subdomains of equal size.
- For each subdomain, the pixels from Y contained in it are ordered linearly, such that close pixels in the image domain are close in this ordering.
- The quantized symbol of any pixel in this ordering is coded relative to the quantized symbol of its predecessor, except for that of the first pixel.
- The coding is done by using a Huffman code.

13 Image Reconstruction at the Decoder

Reconstruction of the compressed image from information

```
\{(\mathtt{y}, Q(I^*(\mathtt{y}))) : \mathtt{y} \in Y\}
```

in four steps:

- (1) Compute Delaunay triangulation $\mathcal{D}(Y)$ of Y;
- (2) Construct unique linear spline $L(Y,\widetilde{I})\in \mathcal{S}_Y$ satisfying

 $L(Y,\widetilde{I})(y)=\widetilde{I}(y), \quad \text{ for all } y\in Y,$

from quantized luminance values $\{\tilde{I}(y) : y \in Y\};$

(3) Obtain reconstructed image by

$$\tilde{I} = \{(x, L(Y, \tilde{I})(x)) : x \in X\}.$$

UH


14 First Comparisons with JPEG2000

Preliminary Remarks.

- We compare the performance of our compression method **AT**² with that of EBCOT, which is the basic algorithm in JPEG2000.
- In each comparison, the compression rate, in *bits per pixel* (bpp), is fixed.
- The quality of the resulting reconstructions is measured by their PSNR values, and by their visual quality.



Geometric Test Image Chessboard. AT versus AT².



Original Image.



JPEG2000



AT





Delaunay triangulation.



Delaunay triangulation.



Geometric Test Image Chessboard.





Original Image Chessboard of size (128×128) .

Reconstruction by JPEG2000 at 0.23 bpp PSNR 18.68 db.

Reconstruction byAT² at 0.23 bppPSNR 45.15 db.



Geometric Real Image Reflex.



Original Image Reflex of size (128×128) .



Reconstruction by JPEG2000 at 0.251 bpp PSNR 28.74 db.



Reconstruction byAT² at 0.251 bppPSNR 42.86 db.

UHI **Ľ**

15 More Recent Comparisons with JPEG2000

Current Version (AT2009):

• L. Demaret, A. Iske, W. Khachabi (2009)

Contextual image compression from adaptive sparse data representations. In: *Signal Processing with Adaptive Sparse Structured Representations*. Workshop Proceedings, Saint-Malo (France), 6.-9. April 2009.

Previous Version (AT2006):

• L. Demaret, A. Iske (2006)

Adaptive image approximation by linear splines over locally optimal Delaunay triangulations. IEEE Signal Processing Letters 13(5), 281-284.

 L. Demaret, N. Dyn, A. Iske (2006) Image compression by linear splines over adaptive triangulations. Signal Processing 86(7), July 2006, 1604–1616.



Comparison between JPEG2000 and AT2009.



Original Image Cameraman of size (256×256) .



Reconstruction by JPEG2000 at 3.247 kB PSNR 29.84 db.



 Reconstruction by

 AT2009 at 3.233 kB

 PSNR 30.66 db



Rate-Distortion Curves for JPEG2000 and AT.



Asymptotic Behaviour of N-term Approximations.

Theorem (BIRMAN & SOLOMJAK 1967): Let $\alpha \in (0,2]$ and $p \ge 1$ satisfy $\alpha > 2/p - 1$. Then, for any $f \in W^{\alpha,p}([0,1]^2)$ we have

$$E_N(f) = \mathcal{O}(N^{-\alpha}) \quad \text{ for } N \to \infty$$

where

$$\mathsf{E}_{\mathsf{N}}(\mathsf{f}) = \inf \left\{ \|\mathsf{f} - \widehat{\mathsf{f}}(\mathcal{Q}_{\mathsf{N}})\|_{\mathsf{L}^{2}([0,1]^{2})}^{2} \colon \mathcal{Q}_{\mathsf{N}} \in \mathfrak{Q} \text{ with } |\mathcal{Q}_{\mathsf{N}}| = \mathsf{N} \right\}.$$

Corollary (DEMARET & I. 2010): Let $\alpha \in (0,2]$ and $p \ge 1$ satisfy $\alpha > 2/p - 1$. Then, for any $f \in W^{\alpha,p}([0,1]^2)$ we have

$$E_N(f) = \mathcal{O}(N^{-\alpha}) \quad \text{ for } N \to \infty$$

where

$$\mathsf{E}_{\mathsf{N}}(f) = \inf \left\{ \|f - \widehat{f}(\mathcal{D}_{\mathsf{N}})\|_{L^{2}([0,1]^{2})}^{2} \colon \mathcal{D}_{\mathsf{N}} \in \mathcal{D} \text{ with } |\mathcal{D}_{\mathsf{N}}| = \mathsf{N} \right\}.$$

UН

Video Compression: Test Case Suzie.



Alba di Canazei, 6-9 September 2010



Ш



Video Compression: Preliminary Remarks.

- Natural videos are composed of a superposition of moving objects ...
- ... usually resulting from anisotropic motions;
- a video may be regarded as a sequence of consecutive natural still images ...
- ... or a video may be regarded as a 3d scalar field;
- it is desirable to work with sparse representations of video data;
- ...
- Adaptive Thinning (AT) extracts significant video pixels ...
- ... to obtain a sparse representation of the video ...
- ... relying on linear splines over anisotropic tetrahedralizations.



Representation of Video Data.

- A digital video V is a rectangular 3d grid of pixels, X.
- Each pixel $x \in X$ bears a greyscale luminance V(x).
- We regard the video as a trivariate function,

 $V:[X] \rightarrow \{0,1,\ldots,2^r-1\}$

where the convex hull [X] of X is the video domain.

INPUT: The video is given by its restriction to the pixels in X,

$$\mathbf{V}\big|_{\mathbf{X}} = \{(\mathbf{x}, \mathbf{V}(\mathbf{x})) : \mathbf{x} \in \mathbf{X}\}.$$

GOAL: Approximation of V from discrete data $V|_{\chi}$.



Linear Splines over Tetrahedralizations.

 \bullet Given any tetrahedralizations $\mathcal{T}(Y)$ of Y, we denote by

$$\mathcal{S}_{Y} = \left\{s: s \in C([Y]) \text{ and } s \big|_{\mathsf{T}} \text{ linear for all } \mathsf{T} \in \mathcal{T}(Y) \right\},$$

the spline space containing all continuous functions over [Y] whose restriction to any tetrahedron in $\mathcal{T}(Y)$ is linear.

- Any element in \mathcal{S}_Y is referred to as a linear spline over $\mathcal{T}(Y)$.
- For given function values $\{V(y) : y \in Y\}$, there is a unique linear spline, $L(Y, V) \in S_Y$, which interpolates V at the points of Y, i.e.,

$$L(Y,V)(y) = V(y), \qquad \text{ for all } y \in Y.$$

Basic Features of Delaunay Tetrahedralizations.

• Uniqueness.

Delaunay tetrahedralization $\mathcal{D}(X)$ is *unique*, if no five points in X are co-spherical.

• Complexity.

For any point set X, its Delaunay tetrahedralization $\mathcal{D}(X)$ can be computed in $\mathcal{O}(N \log N)$ steps, where N = |X|.

• Local Updating.

For any X and $x \in X$, the Delaunay tetrahedralization $\mathcal{D}(X \setminus x)$ of the point set $X \setminus x$ can be computed from $\mathcal{D}(X)$ by re-tetrahedralization of the *cell* $\mathcal{C}(x)$ of x.



Removal of the node x and re-tetrahedralization of its cell $\mathcal{C}(x).$

UH



Numerical Simulation for Test Case Suzie.



Alba di Canazei, 6-9 September 2010



Test Case Suzie: Frame 0000.



Original Frame Suzie.



Delaunay tetrahedralization.



708 significant pixels.



Reconstruction by AT at 34.58 dB.



Test Case Suzie: Frame 0001.



Original Frame Suzie.



Delaunay tetrahedralization.



118 significant pixels.



Reconstruction by AT at 35.15 dB.



Test Case Suzie: Frame 0002.



Original Frame Suzie.



Delaunay tetrahedralization.



287 significant pixels.



Reconstruction by AT at 35.18 dB.



Test Case Suzie: Frame 0003.



Original Frame Suzie.



Delaunay tetrahedralization.



338 significant pixels.



Reconstruction by AT at 34.91 dB.



Test Case Suzie: Frame 0004.



Original Frame Suzie.



Delaunay tetrahedralization.



398 significant pixels.



Reconstruction by AT at 34.98 dB.



Test Case Suzie: Frame 0005.



Original Frame Suzie.



Delaunay tetrahedralization.



448 significant pixels.



Reconstruction by AT at 34.99 dB.



Test Case Suzie: Frame 0006.



Original Frame Suzie.



Delaunay tetrahedralization.



424 significant pixels.



Reconstruction by AT at 34.96 dB.



Test Case Suzie: Frame 0007.



Original Frame Suzie.



Delaunay tetrahedralization.



460 significant pixels.



Reconstruction by AT at 34.92 dB.



Test Case Suzie: Frame 0008.



Original Frame Suzie.



Delaunay tetrahedralization.



534 significant pixels.



Reconstruction by AT at 35.11 dB.



Test Case Suzie: Frame 0009.



Original Frame Suzie.



Delaunay tetrahedralization.



523 significant pixels.



Reconstruction by AT at 34.82 dB.



Test Case Suzie: Frame 0010.



Original Frame Suzie.



Delaunay tetrahedralization.



539 significant pixels.



Reconstruction by AT at 34.89 dB.



Test Case Suzie: Frame 0011.



Original Frame Suzie.



Delaunay tetrahedralization.



534 significant pixels.



Reconstruction by AT at 34.95 dB.



Test Case Suzie: Frame 0012.



Original Frame Suzie.



Delaunay tetrahedralization.



513 significant pixels.



Reconstruction by AT at 35.34 dB.



Test Case Suzie: Frame 0013.



Original Frame Suzie.



Delaunay tetrahedralization.



432 significant pixels.



Reconstruction by AT at 35.30 dB.



Test Case Suzie: Frame 0014.



Original Frame Suzie.



Delaunay tetrahedralization.



364 significant pixels.



Reconstruction by AT at 35.49 dB.



Test Case Suzie: Frame 0015.



Original Frame Suzie.



Delaunay tetrahedralization.



311 significant pixels.



Reconstruction by AT at 35.68 dB.



Test Case Suzie: Frame 0016.



Original Frame Suzie.



Delaunay tetrahedralization.



285 significant pixels.



Reconstruction by AT at 35.82 dB.



Test Case Suzie: Frame 0017.



Original Frame Suzie.



Delaunay tetrahedralization.



293 significant pixels.



Reconstruction by AT at 36.32 dB.



Test Case Suzie: Frame 0018.



Original Frame Suzie.



Delaunay tetrahedralization.



289 significant pixels.



Reconstruction by AT at 36.08 dB.



Test Case Suzie: Frame 0019.



Original Frame Suzie.



Delaunay tetrahedralization.



307 significant pixels.



Reconstruction by AT at 36.25 dB.



Test Case Suzie: Frame 0020.



Original Frame Suzie.



Delaunay tetrahedralization.



292 significant pixels.



Reconstruction by AT at 36.26 dB.



Test Case Suzie: Frame 0021.



Original Frame Suzie.



Delaunay tetrahedralization.



293 significant pixels.



Reconstruction by AT at 36.02 dB.


Test Case Suzie: Frame 0022.



Original Frame Suzie.



Delaunay tetrahedralization.



326 significant pixels.



Reconstruction by AT at 36.06 dB.



Test Case Suzie: Frame 0023.



Original Frame Suzie.



Delaunay tetrahedralization.



341 significant pixels.



Reconstruction by AT at 36.08 dB.



Test Case Suzie: Frame 0024.



Original Frame Suzie.



Delaunay tetrahedralization.



311 significant pixels.



Reconstruction by AT at 36.24 dB.



Test Case Suzie: Frame 0025.



Original Frame Suzie.



Delaunay tetrahedralization.



321 significant pixels.



Reconstruction by AT at 36.16 dB.



Test Case Suzie: Frame 0026.



Original Frame Suzie.



Delaunay tetrahedralization.



320 significant pixels.



Reconstruction by AT at 35.95 dB.



Test Case Suzie: Frame 0027.



Original Frame Suzie.



Delaunay tetrahedralization.



273 significant pixels.



Reconstruction by AT at 35.60 dB.



Test Case Suzie: Frame 0028.



Original Frame Suzie.



Delaunay tetrahedralization.



179 significant pixels.



Reconstruction by AT at 35.48 dB.



Test Case Suzie: Frame 0029.



Original Frame Suzie.



Delaunay tetrahedralization.



669 significant pixels.



Reconstruction by AT at 35.00 dB.



Number of significant pixels:

Total: 11,430; minimal: 118; maximal: 708; average: 381 pixels.

PSNR value:

Overall: 35.45 dB; minimal: 34.58 dB; maximal: 36.32 dB; average: 35.49 dB.



Number of significant pixels.

PSNR values.

UН



Relevant Literature.

- L. Demaret and A. Iske (2010) Anisotropic triangulation methods in image approximation. In: *Approximation Algorithms for Complex Systems*,
 E.H. Georgoulis, A. Iske, and J. Levesley (eds.), Springer, Berlin, 47–68.
- L. Demaret, A. Iske, and W. Khachabi (2010) Sparse representation of video data by adaptive tetrahedralisations. In: *Locally Adaptive Filters in Signal and Image Processing*, L. Florack, R. Duits, G. Jongbloed, M.-C. van Lieshout, L. Davies (eds.), 197–220.
- L. Demaret, A. Iske, and W. Khachabi (2009) Contextual image compression from adaptive sparse data representations. Workshop Proceedings *Signal Processing with Adaptive Sparse Structured Representations*, 6.-9. April 2009 Saint-Malo (France).
- L. Demaret and A. Iske (2006) Adaptive image approximation by linear splines over locally optimal Delaunay triangulations. *IEEE Signal Processing Letters* **13**(5), 281–284.
- L. Demaret, N. Dyn, and A. Iske (2006) Image compression by linear splines over adaptive triangulations. *Signal Processing* **86**(7), July 2006, 281–284.
- L. Demaret, N. Dyn, M.S. Floater, and A. Iske (2005) Adaptive thinning for terrain modelling and image compression. *Advances in Multiresolution for Geometric Modelling*, N.A. Dodgson, M.S. Floater, and M.A. Sabin (eds.), Springer, 321–340.