

## **Research Article**

# **Smoothing Techniques and Augmented Lagrangian Method for Recourse Problem of Two-Stage Stochastic Linear Programming**

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The augmented Lagrangian method can be used for solving recourse problems and obtaining their normal solution in solving two-stage stochastic linear programming problems. The augmented Lagrangian objective function of a stochastic linear problem is not twice differentiable which precludes the use of a Newton method. In this paper, we apply the smoothing techniques and a fast Newton-Armijo algorithm for solving an unconstrained smooth reformulation of this problem. Computational results and comparisons are given to show the effectiveness and speed of the algorithm.

## 1. Introduction

In stochastic programming, some data are random variables with specific possibility distribution [1], which was first introduced by the designer of linear programming problems, Dantzig, in [2].

In this paper, we consider the following two-stage stochastic linear program (slp) with recourse which involves the calculation of an expectation over a discrete set of scenarios:

$$\min_{x \in X} f(x) = c^T x + \phi(x),$$

$$X = \{x \in \mathbb{R}^n : Ax = b, x \ge 0\},$$
(1)

where

$$\phi(x) = E(Q(x,\omega)) = \sum_{i=1}^{N} Q(x,\omega^{i}) \rho(\omega^{i})$$
(2)

and *E* shows the expectation of function  $Q(x, \omega)$  which depend on the random variable  $\omega$ . The function *Q* is defined as follows:

$$Q(x,\omega) = \min_{y \in \mathbb{R}^{n_2}} \left\{ q(\omega)^T y \mid W(\omega)^T y \ge h(\omega) - T(\omega) x \right\},$$
(3)

where  $A \in \mathbb{R}^{m \times n}$ ,  $c \in \mathbb{R}^n$ , and  $b \in \mathbb{R}^m$ . Also, in the problem (3) vector of coefficients  $q(\cdot) \in \mathbb{R}^{n_2}$ , matrix of coefficients

 $W^T(\cdot) \in \mathbb{R}^{m_2 \times n_2}$ , demand vector  $h(\cdot) \in \mathbb{R}^{m_2}$ , and matrix  $T(\cdot) \in \mathbb{R}^{m_2 \times n}$  depend on the random vector  $\omega$  with support space  $\Omega$ . The problems (1) and (3) are called master and recourse problems of stochastic programming, respectively.

We assume that the problem (3) has a solution for each  $x \in X$  and  $\omega \in \Omega$ .

In general, the recourse function  $\phi(x)$  is not differentiable everywhere. Therefore, the traditional methods use nonsmooth optimization techniques [3-5]. However, in the last decade, it is proposed smoothing method for recourse function in standard form of recourse problem [6-11]. In this paper, we apply a smooth approximation technique to smooth recourse function that the recourse problem has inequality linear constrained. For more explanation see Section 2. The approximated problem is based on the least two-norm solution of recourse problem. This paper considers the augmented Lagrangian method to obtain least two-norm solution (Section 3). For convenience, Euclidean least twonorm solution of linear programming problem is named normal solution. This effective method contains solving an unconstrained quadratic problem which its objective function is not twice differentiable. To apply a fast Newton method we use the soothing technique and replace plus function by an accurate smooth approximation [12, 13]. In Section 4, the smoothing algorithm and the numerical results are presented. Also, concluding remarks are given in Section 5.

We now describe our notation. Let  $a = [a_i]$  be a vector in  $\mathbb{R}^n$ . By  $a_+$  we mean a vector in  $\mathbb{R}^n$  whose *i*th entry is 0 if  $a_i < 0$  and equals  $a_i$  if  $a_i \ge 0$ . By  $A^T$  we mean the transpose of matrix A, and  $\nabla f(x_0)$  is the gradient of f at  $x_0$ . For  $x \in \mathbb{R}^n$ , || x || and  $|| x ||_{\infty}$  denote 2-norm and infinity norm, respectively.

#### 2. Approximation of Recourse Function

As mentioned the objective function of (1) is nondifferentiable. This disadvantage property occurs on the recourse function. In this section, there is an attempt to approximate it to a differentiable function.

Using dual of the problem (3), function  $Q(x, \omega)$  can be written as follows:

$$Q(x, \omega) = \max_{z \in \mathbb{R}^{m_2}} (h(\omega) - T(\omega)x)^T z$$
s.t.  $W(\omega) z = q(\omega), \quad z \ge 0.$ 
(4)

Unlike the linear recourse function, the quadratic recours function is differentiable. Thus in this paper, the approximation is based on the following quadratic problem with helpful properties:

$$Q_{\epsilon}(x,\omega) = \max_{z \in \mathbb{R}^{m_2}} (h(\omega) - T(\omega)x)^T z - \frac{\epsilon}{2} \|z\|^2$$
(5)
s.t.  $W(\omega) z = q(\omega), \quad z \ge 0.$ 

The next theorem shows that, for the sufficiently small  $\epsilon > 0$ , the solution of this problem is the normal solution of the problem (4).

**Theorem 1.** For functions  $Q(x, \omega)$  and  $Q_{\epsilon}(x, \omega)$  introduced in (4) and (5), the following can be presented:

- (a) ∃ε > 0 such that, for each ε ∈ (0, ε], the solution for the problem (5) is the normal solution for the problem (4).
- (b) For each ε > 0, function Q<sub>ε</sub>(x, ω) is differentiable with respect to x.
- (c) The gradient of function  $Q_{\epsilon}(x, \omega)$  at point x is

$$\nabla Q_{\epsilon}(x,\omega) = -T^{T}(\omega) z_{\epsilon}^{*}(x,\omega)$$
(6)

in which  $z_{\epsilon}^{*}(x, \omega)$  is the solution of the problem (5).

*Proof.* To prove (a), refer to [14, 15].

Also, (b) and (c) can be easily proved considering that function  $Q_{\epsilon}(x, \omega)$  is the conjugate of function

$$p(z) = \begin{cases} \frac{\epsilon}{2} \|z\|^2, & z \in Z, \\ \infty, & z \notin Z, \end{cases}$$
(7)

where

$$Z = \left\{ z \in \mathbb{R}^{m_2} : W(\omega) \, z = q(\omega) \, , \, z \ge 0 \right\}, \tag{8}$$

Using the approximated recourse function  $Q_{\epsilon}(x, \omega)$ , we can define a differentiable approximation function to the objective function of (1):

$$f_{\epsilon}(x) = c^{T}x + \sum_{i=1}^{N} Q_{\epsilon}(x, \omega^{i}) \rho(\omega^{i}).$$
(9)

By (6), the gradient of above function exists and is obtained by

$$\nabla f_{\epsilon}(x) = c + \sum_{i=1}^{N} \nabla Q_{\epsilon}(x, \omega^{i}) \rho(\omega^{i})$$

$$= c - \sum_{i=1}^{N} T^{T}(\omega^{i}) z_{\epsilon}^{*}(x, \omega^{i}) \rho(\omega^{i}).$$
(10)

This approximation has paved the way to use the optimization algorithm for master problem (1) in which the objective function is substituted by  $f_{\epsilon}(x)$ 

$$\min_{x \in X} f_{\epsilon}(x) \,. \tag{11}$$

In [7], it is considered slp problem with inequality constrained in master problem and equality constrained in recourse problem. Also, in Theorem 2.3 of [7], it is shown that a solution of the approximated problem is a good approximation to a solution of master problem. Here we can express a similar theorem for the problem (1) by using the similar technique in the proof of Theorem 2.3 in [7].

**Theorem 2.** Consider the problem (1). Then, for any  $x \in X$ , there exists an  $\overline{\epsilon}(x) > 0$  such that for any  $\epsilon \in (0, \overline{\epsilon}(x)]$ 

$$\left|f\left(x\right) - f_{\epsilon}\left(x\right)\right| \le \frac{\epsilon}{2}M,$$
 (12)

where *M* is defined as follows:

$$M = \max_{i=1,2,...,N} \left\| z_{\epsilon}^{*}(x,\omega^{i}) \right\|^{2}.$$
 (13)

Let  $x^*$  be a solution of (1) and  $x^*_{\epsilon}$  a solution of (11). Then, there exists an  $\overline{\epsilon} > 0$  such that for any  $0 < \epsilon \leq \overline{\epsilon}$ 

$$\max\left\{f\left(x_{\epsilon}^{*}\right) - f\left(x^{*}\right), f_{\epsilon}\left(x^{*}\right) - f_{\epsilon}\left(x_{\epsilon}^{*}\right)\right\} \le \frac{\epsilon}{2}M.$$
(14)

*Further, one assumes that* f or  $f_{\epsilon}$  *are strongly convex on* X *with modulus*  $\mu > 0$ *. Then,* 

$$\|\boldsymbol{x}^* - \boldsymbol{x}_{\varepsilon}^*\| \le M \frac{\epsilon}{\mu}.$$
(15)

According to Theorem 1, it can be found that for obtaining the gradient of function  $f_{\epsilon}(x)$  in each iteration, we need the normal solution of N linear programming problems (4). In this paper, the augmented Lagrangian method [17] is used for this purpose.

## 3. Smooth Approximation and Augmented Lagrangian Method

In the augmented Lagrangian method, the unconstrained maximization problem is solved which gives the project of a point on the solution set of the problem (4).

Assume that  $\hat{z}$  is an arbitrary vector. Consider the problem of finding the least 2-norm projection  $\hat{z}_*$  of  $\hat{z}$  on the solution set  $Z_*$  of the problem (4)

$$\frac{1}{2} \| \hat{z}_* - \hat{z} \|^2 = \min_{z \in \mathbb{Z}_*} \frac{1}{2} \| z - \hat{z} \|^2,$$

$$Z_* = \left\{ z \in \mathbb{R}^{m_2} : W(\omega) \, z = q(\omega) \,, \ \xi^T z = Q(x, \omega) \,, \ z \ge 0 \right\}.$$
(16)

In this problem, vector x and random variable  $\omega$  are constants; therefore, for simplicity, this is assumed to be  $\xi = h(\omega) - T(\omega)x$ , and function  $\widehat{Q}(\xi)$  is defined in a way that  $\widehat{Q}(\xi) = Q(x, \omega)$ .

Considering that the objective function of the problem (16) is strictly convex, its solution is unique. Let us introduce the Lagrangian function for the problem (16) as follow:

$$L(z, p, \beta, \widehat{z}, \xi, \omega) = \frac{1}{2} \|z - \widehat{z}\|^{2} + p^{T} (W(\omega) z - q(\omega)) + \beta (\xi^{T} z - \widehat{Q}(\xi)),$$
(17)

where  $p \in \mathbb{R}^{n_2}$  and  $\beta \in \mathbb{R}$  are Lagrangian multipliers and  $\xi$ ,  $\hat{z}$  are constant values. Therefore, the dual problem of (16) becomes

$$\max_{\beta \in \mathbb{R}} \max_{p \in \mathbb{R}^{n_2}} \min_{z \in \mathbb{R}^{m_2}_+} L(z, p, \beta, \hat{z}, \xi, \omega).$$
(18)

By solving the inner minimization of the problem (18), duality of the problem (16) is obtained:

$$\max_{\beta \in \mathbb{R}} \max_{p \in \mathbb{R}^{n_2}} \widehat{L}\left(p, \beta, \widehat{z}, \xi\right), \tag{19}$$

where duality function is

$$\widehat{L}(p,\beta,\widehat{z},\xi) = q^{T}(\omega) p - \frac{1}{2} \left\| \left( \widehat{z} + W^{T}(\omega) p + \beta \xi \right)_{+} \right\|^{2} + \beta \widehat{Q}(\xi) + \frac{1}{2} \|\widehat{z}\|^{2}.$$
(20)

The following theorem states that if  $\beta$  is sufficiently large, solving the inner maximization of (19) gives the solution of the problem (16).

**Theorem 3** (see [17]). Consider the following maximization problem

$$\max_{p \in \mathbb{R}^{n_2}} S\left(p, \beta, \hat{z}, \xi, \omega\right) \tag{21}$$

*in which*  $\beta$ ,  $\hat{z}$ , *and*  $\xi$  *are constants, and function*  $S(p, \beta, \hat{z}, \xi)$  *is introduced as follows:* 

$$S(p,\beta,\widehat{z},\xi,\omega) = q^{T}(\omega) p - \frac{1}{2} \left\| \left( \widehat{z} + W^{T}(\omega) p + \beta \xi \right)_{+} \right\|^{2}.$$
(22)

Also, assume that the set  $Z_*$  is nonempty, and the rank of submatrix  $W_l$  of W corresponding to nonzero components of  $\hat{z}_*$  is  $n_2$ . In such a case, there is  $\beta^*$  which for all  $\beta \ge \beta^*$ ,  $\hat{z}_* = (\hat{z} + W^T p(\beta) + \beta\xi)_+$  is the unique and exact solution for the problem (16), where  $p(\beta)$  is the point obtained from solving the problem (21).

Also, in special conditions, the solution for the problem (3) can be also obtained and the following theorem expresses this issue.

**Theorem 4** (see [17]). Assume that the solution set  $Z_*$  is nonempty. For each  $\beta > 0$  and  $\hat{z} \in Z_*$ ,  $y_* = p(\beta)/\beta$  is one exact solution for the linear programming problem (3), where  $p(\beta)$  is the solution for the problem (21).

According to the theorems mentioned above, augmented Lagrangian method presents the following iteration process for solving the problem (16):

$$p_{k+1} \in \arg \max_{p \in \mathbb{R}^{n_2}} \left\{ q^T(\omega) p - \frac{1}{2} \left\| \left( z_k + W(\omega)^T p + \beta \xi \right)_+ \right\|^2 \right\},$$
$$z_{k+1} = \left( z_k + W^T(\omega) p_{k+1} + \beta \xi \right)_+,$$
(23)

where  $z_0$  is an arbitrary vector and here we can use zero vector as initial vector for obtaining normal solution of the problem (4).

We note that the problem (23) is a concave problem and its objective function is piecewise quadratic and is not twice differentiable. Applying the smoothing techniques [18, 19] and replacing  $x_+$  by a smooth approximation, we transform this problem to a twice continuously differentiable problem.

Chen and Mangasarian [19] introduced a family of smoothing functions, which is built as follows. Let  $\rho : R \rightarrow [0, \infty)$  be a piecewise continuous density function satisfying

$$\int_{-\infty}^{+\infty} \rho(s) \, ds = 1, \qquad \int_{-\infty}^{+\infty} |s| \, \rho(s) \, ds < \infty.$$
 (24)

It is obvious that the derivative of plus function is step function, that is,  $(x)_{+} = \int_{-\infty}^{x} \delta(t) dt$ , where the step function  $\delta(x)$  is defined 1 if x > 0 and equals 0 if  $x \le 0$ . Therefore, a smoothing approximation function of the plus function is defined by

$$\varphi(x,\alpha) = \int_{-\infty}^{x} \psi(t,\alpha) dt, \qquad (25)$$

where  $\psi(x, \alpha)$  is smoothing approximation function of step function and is defined as

$$\psi(x,\alpha) = \int_{-\infty}^{x} \alpha \rho(\alpha t) dt.$$
 (26)

By choosing

$$\rho(s) = \frac{e^{-s}}{\left(1 + e^{-s}\right)^2},$$
(27)

specific cases of these approaches are obtained as follows:

$$\psi(x,\alpha) = \frac{1}{1+e^{-\alpha x}} \approx \delta(x),$$

$$\varphi(x,\alpha) = x + \frac{1}{\alpha} \log(1+e^{-\alpha x}) \approx (x)_{+}.$$
(28)

The function  $\varphi$  with a smoothing parameter  $\alpha$  is used here to replace the plus function of (22) to obtain a smooth reformulation of function (22):

$$\widehat{S}(p,\beta,\widehat{z},\xi,\beta,\omega,\alpha) := q^{T}(\omega) p - \left\|\varphi(\widehat{z} + W^{T}(\omega)p + \beta\xi,\alpha)\right\|^{2}.$$
(29)

Therefore, we have the following iterative process instead of (23) and (28):

$$p_{k+1} \in \arg \max_{p \in \mathbb{R}^{n_2}} \left\{ q^T(\omega) p - \left\| \varphi(z_k + W^T(\omega)p + \beta\xi, \alpha) \right\|^2 \right\},$$
$$z_{k+1} = \varphi\left( z_k + W^T(\omega) p_{k+1} + \beta\xi, \alpha \right).$$
(30)

It can be shown that as the smoothing parameter  $\alpha$  approaches infinity any solution of smooth problem (29) approaches the solution of the equivalent problem (22) (see [19]).

We begin with a simple lemma that bounds the square difference between the plus function  $x_+$  and its smooth approximation  $\varphi(x, \alpha)$ .

**Lemma 5** (see [13]). For  $x \in \mathbb{R}$  and  $|x| < \omega$ 

$$\varphi^{2}(x,\alpha) - (x_{+})^{2} \leq \left(\frac{\log(2)}{\alpha}\right)^{2} + \frac{2\overline{\alpha}}{\alpha}\log(2),$$
 (31)

where  $\varphi(x, \alpha)$  is the  $\varphi$  function of (28) with smoothing parameter  $\alpha > 0$ .

Theorem 6. Consider the problems (21) and

$$\max_{p \in \mathbb{R}^{n_2}} \widehat{S}\left(p, \beta, \widehat{z}, \xi, \omega, \alpha\right).$$
(32)

*Then, for any*  $p \in \mathbb{R}^{n_2}$  *and*  $\alpha > 0$ 

$$\left| S\left(p,\beta,\widehat{z},\xi,\omega\right) - \widehat{S}\left(p,\beta,\widehat{z},\xi,\omega,\alpha\right) \right| \leq \left(\frac{\log\left(2\right)}{\alpha}\right)^{2} + 2Mm_{2}\frac{\log\left(2\right)}{\alpha},$$
(33)

where M is defined as follows:

$$M = \max_{1 \le i \le m_2} \left| \widehat{z}_i + W_i^T(\omega) p + \beta \xi_i \right|.$$
(34)

Let  $p^*$  be a solution of (21) and  $p^*_{\alpha}$  a solution of (32). Then

$$\max \left\{ S\left(p^{*}, \beta, \widehat{z}, \xi, \omega\right) - S\left(p^{*}_{\alpha}, \beta, \widehat{z}, \xi, \omega\right), \\ \widehat{S}\left(p^{*}_{\alpha}, \beta, \widehat{z}, \xi, \omega, \alpha\right) - \widehat{S}\left(p^{*}, \beta, \widehat{z}, \xi, \omega, \alpha\right) \right\}$$

$$\leq \left(\frac{\log\left(2\right)}{\alpha}\right)^{2} + 2Mm_{2}\frac{\log\left(2\right)}{\alpha}.$$
(35)

Further, one assumes that  $W^T$  is a full rank matrix. Then,

$$\|p^* - p^*_{\alpha}\| \le \mu \left(\frac{2\log(2)}{\alpha}\right)^2 + 8Mm_2\mu \frac{\log(2)}{\alpha}.$$
 (36)

*Proof.* For any  $\alpha > 0$  and  $p \in \mathbb{R}^{n_2}$ 

$$\varphi\left(\widehat{z}_{i}+W_{i}^{T}\left(\omega\right)p+\beta\xi_{i},\alpha\right)\geq\left(\widehat{z}_{i}+W_{i}^{T}\left(\omega\right)p+\beta\xi_{i},\alpha\right)_{+}.$$
(37)

Hence

$$\begin{aligned} \left| S\left(p,\beta,\widehat{z},\xi,\omega\right) - \widehat{S}\left(p,\beta,\widehat{z},\xi,\omega,\alpha\right) \right| \\ &= S\left(p,\beta,\widehat{z},\xi,\omega\right) - \widehat{S}\left(p,\beta,\widehat{z},\xi,\omega,\alpha\right) \\ &= \left\| \varphi(\widehat{z} + W^{T}(\omega)p + \beta\xi,\alpha) \right\|^{2} \\ &- \left\| \left(\widehat{z} + W^{T}(\omega)p + \beta\xi,\alpha\right)_{+} \right\|^{2} \end{aligned} \tag{38}$$
$$&= \sum_{i=1}^{m_{2}} \left( \varphi^{2}\left(\widehat{z}_{i} + W_{i}^{T}(\omega)p + \beta\xi_{i},\alpha\right) \\ &- \left(\widehat{z}_{i} + W_{i}^{T}(\omega)p + \beta\xi_{i},\alpha\right)^{2}_{+} \right). \end{aligned}$$

By using Lemma 5, we get that

$$\begin{split} \left| S\left(p,\beta,\widehat{z},\xi,\omega\right) - \widehat{S}\left(p,\beta,\widehat{z},\xi,\omega,\alpha\right) \right| \\ &\leq \sum_{i=1}^{m_2} \left( \left(\frac{\log\left(2\right)}{\alpha}\right)^2 + 2\left|\widehat{z}_i + W_i^T\left(\omega\right)p + \beta\xi_i,\alpha\right|\frac{\log\left(2\right)}{\alpha}\right) \right. \\ &\leq \left(\frac{\log\left(2\right)}{\alpha}\right)^2 + 2\frac{\log\left(2\right)}{\alpha} \\ &\qquad \times \sum_{i=1}^{m_2} \max_{1 \leq i \leq m_2} \left|\widehat{z}_i + W_i^T\left(\omega\right)p + \beta\xi_i\right| \\ &= \left(\frac{\log\left(2\right)}{\alpha}\right)^2 + 2Mm_2\frac{\log\left(2\right)}{\alpha}. \end{split}$$
(39)

From above inequality, we have

$$\widehat{S}(p,\beta,\widehat{z},\xi,\omega,\alpha) \leq S(p,\beta,\widehat{z},\xi,\omega)$$

$$\leq \widehat{S}(p,\beta,\widehat{z},\xi,\omega,\alpha) + \left(\frac{\log(2)}{\alpha}\right)^{2}$$

$$+ 2Mm_{2}\frac{\log(2)}{\alpha}.$$
(40)

 TABLE 1: Comparative between smooth augmented Lagrangian Newton method (SALN) and CPLEX solver.

N. P	Recourse problem $n_2 \times m_2 \times d$	Solver	$\ Wz - q\ _{\infty}$	$ \widehat{Q}(\xi) - \xi^T z $	<i>z</i>	Time (second)
P1	2	SALN	9.9135 <i>e</i> - 011	7.2760 <i>e</i> - 012	41.0362	0.3292
	$50 \times 50 \times 0.68$	CPLEX	1.0717e - 007	5.2589e - 007	41.0362	0.3182
P2		SALN	1.8622 <i>e</i> – 010	4.3074e - 009	57.1793	0.1275
	$100 \times 105 \times 0.4$	CPLEX	6.5591 <i>e</i> – 008	1.6105e – 006	57.1826	0.1585
	$150 \times 150 \times 0.5$	SALN	3.1559e - 010	5.6361 <i>e</i> - 009	74.9098	0.6593
P3		CPLEX	5.9572e - 010	1.2014e - 0.09	74 9098	0 1605
P4	$200 \times 200 \times 0.5$	SALN	53819e - 010	1.0506e - 0.08	85 6646	0.2530
		CPLEX	11734e - 007	1.3964e - 0.06	85 6646	0.1820
P5	$300 \times 300 \times 0.5$	SALN	9.4178e = 010	2.7951e - 0.08	102 4325	2 1356
		CPLEX	4.1638e - 010	4.4456e = 0.09	102.4325	0.1830
		SALN	1.2787e = 0.09	2.6226a = 0.08	110 / 189	3 2102
P6	$350 \times 350 \times 0.5$	CDLEY	7.7398a = 0.10	2.0220e = 000	110.4189	0.2116
		CPLEA	7.7398e - 010	5.4452e - 009	124 1204	0.2110
P7	$450\times500\times0.05$	SALN	1.0504e - 010	4.7094e - 009	124.1204	0.7807
		CPLEX	9.0949 <i>e</i> - 013	1.63/1e - 011	124.1205	0.2606
P8	$500 \times 550 \times 0.04$	SALN	1.0425e - 010	1.0241e - 009	134.3999	0.7567
		CPLEX	6.1618 <i>e</i> – 011	1.9081e - 009	134.3999	0.2660
Р9	$700 \times 800 \times 0.6$	SALN	4.6139e - 009	2.3908e - 007	153.9782	7.1364
		CPLEX	9.1022e - 009	1.3768e - 007	153.9906	0.8435
P10	$900 \times 1100 \times 0.4$	SALN	5.3396 <i>e</i> – 009	1.5207e - 007	178.1151	7.9383
		CPLEX	4.5020e - 011	2.0373e - 010	178.1163	1.3643
P11	$1500 \times 2000 \times 0.1$	SALN	1.1289e - 008	2.0696e - 007	231.5284	13.2493
		CPLEX	7.5886 <i>e</i> – 011	1.1059e - 009	231.5286	2.1343
P12	$1000 \times 2000 \times 0.01$	SALN	3.6398e - 010	7.9162e - 009	198.4905	5.2620
		CPLEX	6.8212e - 013	1.1642e - 010	198.4922	0.6752
P13	$1000\times5000\times0.001$	SALN	2.7853e - 010	1.9281e - 010	190.0141	0.2709
1 15		CPLEX	4.8203e - 010	1.5425e - 009	190.0142	0.2162
D1/	$1000 \times 10000 \times 0.001$	SALN	1.1221e - 010	1.4988e - 009	212.2416	0.4867
1 14		CPLEX	7.9094e - 009	2.6691 <i>e</i> - 007	212.3453	0.3353
D15	$1000 \times 1e5 \times 0.001$	SALN	9.7702 <i>e</i> - 010	2.3283e - 009	231.7930	3.7511
P13		CPLEX	2.7285 <i>e</i> - 012	1.8044e - 009	231.8763	1.2472
DIC	$1000 \times 1e6 \times 0.0002$	SALN	9.9432 <i>e</i> - 013	2.0464e - 012	121.3937	9.0948
P10		CPLEX	9.9098 <i>e</i> - 006	1.1089e - 004	121.3940	3.7848
D15	$100 \times 1e6 \times 0.001$	SALN	3.9563 <i>e</i> - 010	5.6607 <i>e</i> - 009	73.8493	6.8537
PI/		CPLEX	2.0082e - 003	1.1777 <i>e</i> – 002	73.9412	2.5582
	$10 \times 1e4 \times 0.001$	SALN	2.2737 <i>e</i> – 013	7.9581 <i>e</i> – 013	19.5735	0.0166
P18		CPLEX	2.2737 <i>e</i> – 013	1.1369 <i>e</i> – 013	19.5739	0.1386
	$10 \times 1e6 \times 0.001$	SALN	1.2478e - 009	1.0710e - 008	18.8192	5.6399
P19		CPLEX	5.9615e - 0.04	64811e - 005	18 8863	2 5623
	$10 \times 5e6 \times 0.01$	SALN	2.0425e - 0.08	1.1816e - 0.08	20.8470	28 7339
P20		CPLEX	4.9966e + 004	3.3500e + 0.05	0.0000	1 9482
		SALN	3.8654a = 0.12	7.7875e = 0.12	42.0698	7 8895
P21	$100 \times 1e6 \times 0.01$	CDIEV	1.7094a = 004	7.7875e = 012	42.0098	9.8324
			1.7774t = 0.04	0.2712e = 003	42.0731	0.0324
P22	$100 \times 1e4 \times 0.1$	ODLEY	$1.1004\ell = 012$	2.1000e - 012	37.4303	0.1440
		CPLEA	1.2318e - 003	1.31/4e = 003	37.4303	0.3099
P23	$100 \times 1e5 \times 0.05$	SALN	1.70556 - 012	0.3121e - 012	45.5944	1.13/9
		CPLEX	3.982/e - 006	2.8309 <i>e</i> - 006	45.5944	1.5729

TABLE 1: Continued.

N. P	Recourse problem $n_2 \times m_2 \times d$	Solver	$\left\ Wz - q\right\ _{\infty}$	$ \widehat{Q}(\xi) - \xi^T z $	<i>z</i>	Time (second)
P24	$1000 \times 5e4 \times 0.1$	SALN	1.3074e - 012	2.4102 <i>e</i> - 011	117.4693	15.1065
		CPLEX	6.7455e - 007	6.7379 <i>e</i> – 006	117.4699	20.3967
P25	$1000 \times 1e5 \times 0.08$	SALN	2.5011 <i>e</i> - 012	1.0516e – 011	116.6964	23.2540
		CPLEX	9.4298e - 006	5.5861e - 004	116.6964	33.4319

Therefore

$$S(p^{*}, \beta, \widehat{z}, \xi, \omega) - S(p^{*}_{\alpha}, \beta, \widehat{z}, \xi, \omega)$$

$$\leq \widehat{S}(p^{*}, \beta, \widehat{z}, \xi, \omega, \alpha) - \widehat{S}(p^{*}_{\alpha}, \beta, \widehat{z}, \xi, \omega, \alpha)$$

$$+ \left(\frac{\log(2)}{\alpha}\right)^{2} + 2Mm_{2}\frac{\log(2)}{\alpha}$$

$$\leq \left(\frac{\log(2)}{\alpha}\right)^{2} + 2Mm_{2}\frac{\log(2)}{\alpha},$$

$$\widehat{S}(p^{*}_{\alpha}, \beta, \widehat{z}, \xi, \omega, \alpha) - \widehat{S}(p^{*}, \beta, \widehat{z}, \xi, \omega, \alpha)$$

$$\leq S(p^{*}_{\alpha}, \beta, \widehat{z}, \xi, \omega) - S(p^{*}, \beta, \widehat{z}, \xi, \omega)$$

$$+ \left(\frac{\log(2)}{\alpha}\right)^{2} + 2Mm_{2}\frac{\log(2)}{\alpha}$$

$$\leq \left(\frac{\log(2)}{\alpha}\right)^{2} + 2Mm_{2}\frac{\log(2)}{\alpha}.$$
(41)

Suppose that  $W^T$  is full rank. Then the Hessian of  $\hat{S}$  is negative definite, and  $\hat{S}$  is strongly concave on bounded sets. By the definition of strong concavity, for any  $\gamma \in (0, 1)$ ,

$$\widehat{S}(\gamma p_{\alpha}^{*} + (1 - \gamma) p^{*}, \beta, \widehat{z}, \xi, \omega, \alpha) - \gamma \widehat{S}(p_{\alpha}^{*}, \beta, \widehat{z}, \xi, \omega, \alpha)$$
$$- (1 - \gamma) \widehat{S}(p^{*}, \beta, \widehat{z}, \xi, \omega, \alpha)$$
$$\geq \frac{1}{2} \mu \gamma (1 - \gamma) \| p_{\alpha}^{*} - p^{*} \|^{2}.$$
(42)

Let  $\gamma = 1/2$ , then

$$\begin{aligned} \frac{1}{8}\mu \|p_{\alpha}^{*} - p^{*}\|^{2} \\ &\leq \widehat{S}\left(\frac{1}{2}\left(p_{\alpha}^{*} + p^{*}\right), \beta, \widehat{z}, \xi, \omega, \alpha\right) \\ &- \frac{1}{2}\left(\widehat{S}\left(p_{\alpha}^{*}, \beta, \widehat{z}, \xi, \omega, \alpha\right) + \widehat{S}\left(p^{*}, \beta, \widehat{z}, \xi, \omega, \alpha\right)\right) \\ &\leq \widehat{S}\left(p_{\alpha}^{*}, \beta, \widehat{z}, \xi, \omega, \alpha\right) \\ &- \frac{1}{2}\left(\widehat{S}\left(p_{\alpha}^{*}, \beta, \widehat{z}, \xi, \omega, \alpha\right) + \widehat{S}\left(p^{*}, \beta, \widehat{z}, \xi, \omega, \alpha\right)\right) \\ &\leq \frac{1}{2}\left(\widehat{S}\left(p_{\alpha}^{*}, \beta, \widehat{z}, \xi, \omega, \alpha\right) - \widehat{S}\left(p^{*}, \beta, \widehat{z}, \xi, \omega, \alpha\right)\right) \\ &\leq \frac{1}{2}\left(\frac{\log\left(2\right)}{\alpha}\right)^{2} + Mm_{2}\frac{\log\left(2\right)}{\alpha}. \end{aligned}$$

43)

Considering the advantage of the twice differentiability of the objective function of the problem (32) allows us to use a quadratically convergent Newton algorithm with an Armijo stepsize [20] that makes the algorithm globally convergent.

#### 4. Numerical Results and Algorithm

In each iteration of the process (30), one concave, quadratic, unconstrained maximization problem is solved. For solving it, the fast Newton method can be used.

In the algorithm, the Hessian matrix may be singular, thus we use a modified Newton. The direction in each iteration for solving (30) is obtained through the following relation:

$$d_{s} = -\left(\nabla_{p}^{2}\widehat{S}\left(p,\beta,\widehat{z},\xi,\omega,\alpha\right) - \delta I_{n_{2}}\right)^{-1}\left(\nabla_{p}\widehat{S}\left(p,\beta,\widehat{z},\xi,\omega,\alpha\right)\right),$$
$$p_{s+1} = p_{s} + \lambda_{s}d_{s},$$
(44)

where  $\delta$  is a small positive number,  $I_{n_2}$  is the identity matrix of order  $n_2$ , and  $\lambda_s$  is the suitable step length that Armijo algorithm is used for determining it (see Algorithm 1).

The proposed algorithm was applied to solve some recourse problems. Table 1 compares this algorithm with CPLEX v. 12.1 solver for quadratic convex programming problems (5). As is evident from Table 1, most of recourse problems could be solved more successful by the algorithm which is based on smooth augmented Lagrangian Newton method (SALN) than CPLEX package (for illustration see the problems 21–25 in Table 1). This algorithm gives us high accuracy and the solution with minimum norm in suitable time (see last column of Table 1). Also, we can find that CPLEX is better than the algorithm proposed for some recourse problems in which the matrices are approximately square (Ex. line 5–12).

The test generator generates recourse problems. These problems are generated using the MATLAB code show in Algorithm 2.

The algorithm considered for solving several recourse problems was run on a computer with 2.5 dual-core CPU and 4 GB memory in MATLAB 7.8 programming environment. Also, in the generated problems, recourse matrix W is the Sparse matrix  $(n_2 \times m_2)$  with the density d. The constants  $\beta$  and  $\delta$  in the above algorithm in (44) were selected 1 and 10<sup>-8</sup>, respectively.

In Table 1, the second column indicates the size and density of matrix W, the forth column indicates the feasibility of the primal problem (4), and the next column indicates the error norm function of this problem (the MATLAB code of this paper is available from the authors upon request).

```
Choose a z_0 \in \mathbb{R}^{m_2}, \alpha > 1, i > 1, \epsilon > 0 be error tolerance and \delta is a small positive number.

i = 0;

While ||z_i - z_{i-1}||_{\infty} \ge \epsilon

Choose a p_0 \in \mathbb{R}^{n_2} and set k = 0.

While ||\nabla_p \widehat{S}(p_k, \beta, z_i, \xi, \omega, \alpha)||_{\infty} \ge \epsilon

Choose \lambda_k = \max\{s, s\sigma, s\sigma^2, \ldots\} such that

\widehat{S}(p_k, \beta, z_i, \xi, \omega, \alpha) - \widehat{S}(p_k + \lambda_k d_k, \beta, z_i, \xi, \omega, \alpha) \ge -\lambda_i \mu \nabla (\widehat{S}(p_k, \beta, z_i, \xi, \omega, \alpha))^T d_k,

where,

d_k = -(\nabla^2 \widehat{S}(p_k, \beta, z_i, \xi, \omega, \alpha, ) - \delta I_{n_2})^{-1} \nabla \widehat{S}(p_k, \beta, z_i, \xi, \omega, \alpha), s > 0 be a constant,

\sigma \in (0, 1) and \mu \in (0, 1).

Put p_{k+1} = p_k + \lambda_k d_k, k = k + 1 and \alpha = i\alpha.

end

Set p_{i+1} = p_{k+1}, z_{i+1} = \varphi(z_i + W^T(\omega)p_{i+1} + \beta\xi, \alpha) and i = i + 1.

end
```

ALGORITHM 1: Newton method with the Armijo rule.

```
\label{eq:spectral_series} \begin{split} & \mbox{``Sgen: Generate random solvable recourse problems:} \\ & \mbox{``Input: m,n,d(ensity); Output: W,q,\xi;} \\ & \mbox{``m=input(`Enter $n_2$:')} \\ & \mbox{``n=input(`Enter $m_2$:')} \\ & \mbox{``d=input(`Enter $d:')} \\ & \mbox{``pl=inline('(abs(x)+x)/2')} \\ & \mbox{``m=sprand($n_2$,$m_2$,$d);W=100*(W=0.5*spones(W));} \\ & \mbox{``ssparse(10*pl(rand($m_2$,1)));} \\ & \mbox{``ssparse(10*pl(rand($m_2$,1)));} \\ & \mbox{``ssparse(isign(pl(rand($n_2$,1)-rand($n_2$,1)))),0,$n_2$,$n_2$)} \\ & \mbox{``ssparse((rand($m_2$,1)-rand($n_2$,1)));} \\ & \mbox{``ssparse(isign($pl(rand($m_2$,1)-rand($n_2$,1)))),0,$m_2$,$n_2$)} \\ & \mbox{``ssparse($m_2$,1)-rand($n_2$,1));} \\ & \mbox{``ssparse($m_2$,1)-sign($z)$),0,$m_2$,$m_2$)*ones($m_2$,1));} \\ & \mbox{``ssparse($m_2$,1)-rign($z)$),0,$m_2$,$m_2$)*ones($m_2$,1));} \\ & \mbox{``ssparse($m_2$,1)-rign($z)$),0,$m_2$,$m_2$)*ones($m_2$,1),0;} \\ & \mbox{``ssparse($m_2$,1)-rign($z)$),0,$m_2$,m_2$)*ones($m_2$,1),0;} \\ & \mbox{``ssparse($m_2$,1)-rign($m_2$,1),0;} \\ & \mbox{
```



## 5. Conclusion

In this paper, a smooth reformulation process, based on augmented Lagrangian algorithm, was proposed for obtaining the normal solution of recourse problem of a stochastic linear programming. This smooth iterative process allows us to use a quadratically convergent Newton algorithm, which accelerates obtaining the normal solution.

Table 1 shows that the proposed algorithm has appropriate speed in most of the problems. This result, specifically, can be observed in recourse problems with the matrix of coefficients in which the number of constraints is noticeably more than the number of variables. The more challenging is solving the problems which their coefficient matrix is square (the numbers of constraints and variables get closer to each other), and more time is needed by the algorithm for solving the problem.

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