

## A NOTE ON THE VERTEX-SWITCHING RECONSTRUCTION

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**ABSTRACT.** Bounds on the maximum and minimum degree of a graph establishing its reconstructibility from the vertex switching are given. It is also shown that any disconnected graph with at least five vertices is reconstructible.

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### 1. INTRODUCTION.

A switching  $G_v$  of a graph  $G$  at vertex  $v$  is a graph obtained from  $G$  by deleting all edges incident to  $v$  and inserting all possible edges to  $v$  which are not in  $G$ . Since switching is a commutative operation, i.e.,  $(G_v)_u = (G_u)_v$ , the definition can be naturally extended to arbitrary subsets of the vertex set  $V(G)$ . Thus,  $G_A$  is defined for all  $A \subseteq V(G)$ .

The Vertex-Switching Reconstruction Problem, proposed by Stanley [1], asks: Is  $G$  uniquely determined up to isomorphism by the set (deck),  $\{G_v\}_{v \in V(G)}$ ? If the answer is "yes" then  $G$  is called reconstructible.

It was shown in [1] that any graph  $G$  with  $n = |V(G)| \neq 0 \pmod{4}$  is reconstructible. It seems that a little is known about the case  $n = 0 \pmod{4}$ . However, Stanley pointed out [1], that the degree sequence of a graph, and consequently, the number of edges easily reconstructible, provided  $n \neq 4$ . Bounds on the number of edges in a graph,  $e(G)$ , establishing its reconstructibility was given [2]. Namely:

$$e(G) \notin \left[ \frac{n(n-2)}{4}, \frac{n^2}{4} \right], \quad n \neq 4.$$

As might be expected, in virtue of the last result,  $G$  is reconstructible if it has a vertex of degree not close to  $n/2$  or if  $G$  is disconnected. Here we will prove the last claim (Theorem 2) and show that for sufficiently large  $n$  a graph is reconstructible if  $\max(\Delta, n - \delta) > 0.9n$ , where  $\Delta$  and  $\delta$  are the maximum and the minimum degree of  $G$  respectively. Actually, we prove a little more, namely:

## 2. MAIN RESULTS.

THEOREM 1. If  $\min \left( n \binom{n-1}{\Delta}, n \binom{n-1}{\delta} \right) < 2^{n/2-3}$ ,

then  $G$  is reconstructible.

PROOF. In virtue of the quoted result of Stanley, we may assume  $n = 0 \pmod{4}$ . We will consider a graph  $G$  as a spanning subgraph of a fixed copy of the complete graph  $K_n$ . The switching equivalence class  $G^*$  of  $G$  is the set of all  $H \subset K_n$  isomorphic to  $G$  such that  $H = G_A$  for some switching  $A \subseteq V(G)$ .

For each subgraph  $g \subset G$ , let  $\mu(G^* \supset g)$  be the number of those elements of  $G^*$  which contain a fixed copy of  $g$ .

First we show that  $G$  is reconstructible if

$$\frac{\mu(G^* \supset g) s(g \rightarrow K_n)}{s(g \rightarrow G)} < 2^{n/2-2}, \quad (2.1)$$

where  $s(H \rightarrow F)$  is the number of the subgraphs of  $F$  isomorphic to  $H$ .

Observe that

$$|G^*| s(g \rightarrow G) \leq \mu(G^* \supset g) s(g \rightarrow K_n). \quad (2.2)$$

On the other hand, consider the set  $S_i = \{A : G_A \in G^*, |A| = i\}$ .

Observe that  $\Sigma |S_i| = 2|G^*|$  since  $G_A$  and  $G_{\bar{A}}$ ,  $\bar{A} = V(G) \setminus A$ , are identical. It is known that for a nonreconstructible graph  $|S_{4i}| \geq \binom{n/2}{2i}$  ([2], Corollary 2.4). Thus,

if  $G$  is not reconstructible then

$$2|G^*| \geq \Sigma \binom{n/2}{2i} = 2^{n/2-1}. \quad (2.3)$$

Comparing (2.2) and (2.3), we get that (2.1) is enough for the reconstructibility of  $G$ .

Let now  $g$  be a star  $K_{1,\Delta}$ . Observe that  $\mu(G^* \supset K_{1,\Delta}) \leq 2$  since the only proper switching, possibly preserving a fixed copy of  $K_{1,\Delta}$ , is  $A = V(K_{1,\Delta})$ . Furthermore,

$s(g \rightarrow K_n) = n \binom{n-1}{\Delta}$ . Hence, by (2.1),  $G$  is reconstructible if  $n \binom{n-1}{\Delta} < 2^{n/2-3}$ .

Now, to complete the proof, one has to consider the complementary graph  $\bar{G}$ , which is reconstructible iff  $G$  is.  $\square$

Now we will prove that disconnected graphs are reconstructible. First we need the following simple lemma:

LEMMA 1. Suppose that nonisomorphic graphs  $G$  and  $H$  have the same deck. Then for any  $v \in V(G)$  there is  $u \in V(G)$ ,  $v \neq u$ , such that  $G_{vu} \cong H$ .

PROOF. Since the decks of  $G$  and  $H$  are equal then there is a bijection  $\phi: V(G) \rightarrow V(H)$  such that  $G_v \cong H_{\phi(v)}$ . Let  $h_v: H_{\phi(v)} \rightarrow G_v$  be an isomorphism. Choosing  $u = h(\phi(v))$  we obtain  $G_{vu} \cong H$ . Moreover, since  $G_{vv} = G$ , then  $v \neq \phi(v)$ .  $\square$

COROLLARY 1. Let  $n \neq 4$ . If  $G_{vu}$  and  $G$ ,  $v \neq u$ , have the same deck then  $\deg(v) + \deg(u) = n$  or  $n - 2$ , depending on whether  $v$  and  $u$  are adjacent in  $G$  or are not.

PROOF. Let  $e(v,u)$  be the number of edges between  $v$  and  $u$ . Since  $e(G) = e(H)$  then

$$\deg(v) + \deg(u) - 2e(v,u) = \frac{1}{2} \cdot 2(n - 2) = n - 2. \quad \square$$

COROLLARY 2. If  $G$  is not reconstructible and  $n \neq 4$  then  $n - 2 \leq \delta + \Delta \leq n$ .

PROOF. This easily follows from Lemma 1 and Corollary 1. We omit the details.  $\square$

THEOREM 2. Any disconnected graph is reconstructible, provided  $n \neq 4$ .

PROOF. Assume the contrary. Then there are two nonisomorphic graphs  $G$  and  $H$  with the same deck,  $n \neq 4$ , and, say,  $G$  is disconnected. Denote by  $C$  a minimal connected component of  $G$ . First we show that  $G$  has exactly two connected components and  $C \cong K_{\delta+1}$ .

Let  $v$  be a vertex of the minimal degree in  $C$ , and let  $u$  be such a vertex that  $G_{vu} \cong H$ . We claim that either  $u = \phi(v) \in \bar{C}$  or  $G$  is regular of degree  $\frac{n-2}{2}$ . Indeed, otherwise,

$$|C| \geq \max(\deg(v) + 1, \deg(u) + 1) > n/2,$$

which contradicts the minimality of  $C$ . Furthermore, if  $G$  is regular then again  $v$  and  $u$  are in different components since, otherwise, the degree sequences of  $G$  and  $G_{vu}$  are different. Now it follows by Corollary 1,  $\deg(v) + \deg(u) = n - 2$ . Therefore,  $G$  has exactly two components,  $C$  is regular, and  $\Delta \geq n/2$ .

Let us show that  $C$  is just  $K_{\delta+1}$ . Since all vertices of degree  $\Delta$  are in  $\bar{C}$ , we have

$$\deg(v) + 1 \leq |C| \leq n - \Delta - 1.$$

Hence, applying Corollary 2, we get

$$n - 2 \leq \delta + \Delta \leq \deg(v) + \Delta \leq n - 2.$$

Thus,  $\deg(v) = \delta$ ,  $\deg(u) = \Delta$ , and  $C \cong K_{\delta+1}$ .

Finally,  $G_{vu} \cong G$  since  $\deg(v) = |C| - 1$ ,  $u \in \bar{C}$  and  $\deg(u) = \Delta = |\bar{C}| - 1$ ,

which is a contradiction. This completes the proof.  $\square$

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