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A GRAPH AND ITS COMPLEMENT WITH SPECIFIED PROPERTIES I:

CONNECTIVITY

JIN AKIYAMA*

Mathematics Department Nippon Ika University Kawasaki, Japan

FRANK HARARY**

Mathematics Department University of Michigan Ann Arbor, Michigan 48109

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Dedicated to Karl Menger

<u>ABSTRACT</u>. We investigate the conditions under which both a graph G and its complement \overline{G} possess a specified property. In particular, we characterize all graphs G for which G and \overline{G} both (a) have connectivity one, (b) have line-connectivity one, (c) are 2-connected, (d) are forests, (e) are bipartite, (f) are outerplanar and (g) are eulerian. The proofs are elementary but amusing. KEY WORDS AND PHRASES. Graphs, Complement.

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Visiting Scholar 1978-79 at The University of Michigan. ** Vice-President, Calcutta Mathematical Society, 1978 and 1979.

1. CONNECTIVITY.

The <u>connectivity</u> (or line-connectivity) $\kappa = \kappa(G)$ (or $\lambda = \lambda(G)$) of a graph G is the minimum number of points (or lines) whose removal results in a disconnected or a trivial graph. We write $\bar{\kappa}$ (or $\bar{\lambda}$) for $\kappa(\bar{G})$ (or $\lambda(\bar{G})$) where \bar{G} is the complement of G. We follow the graph theoretic terminology and notation of the book [1]. Recall that Δ denotes the maximum degree among all points of G.

LEMMA 1. The complement \overline{G} of a connected graph G is connected if and only if G has no spanning complete bipartite subgraph.

PROOF. If G has a spanning complete bipartite subgraph, then \overline{G} clearly contains no line joining the two parts, hence must be disconnected. Conversely, if \overline{G} is disconnected, then any bipartition of V(G) in which one part consists of the points of precisely one component of \overline{G} gives a spanning complete bipartite subgraph of G.

The next statement is an easy consequence of the lemma.

THEOREM 1. A graph G with p points satisfies the condition $\kappa = \overline{\kappa} = 1$ if and only if G is a graph with either

(1) $\kappa = 1$ and $\Delta = p - 2$, or

(2) $\kappa = 1$, $\Delta \leq p - 3$ and G has a cutpoint v with endline e and endpoint u such that G - u contains a spanning complete bipartite subgraph.

PROOF. We note that if $\kappa = \bar{\kappa} = 1$, then the degree of each point of G is at most p - 2, since otherwise \bar{G} would contain an isolated point which would make $\bar{\kappa} = 0$.

(1) Let G be a graph with $\Delta = p - 2$ and $\kappa = 1$, as in Figure 1a.



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The removal of any cutpoint v from G results in a disconnected graph, so that $\overline{G-v}$ is connected. Since $\Delta = p - 2$ by hypothesis, v is adjacent in \overline{G} to a point of \overline{G} - v. Thus \overline{G} is connected. Furthermore \overline{G} has an endline since Δ = p - 2, and hence \overline{G} has a cutpoint (as illustrated in Figure 1b), so that $\overline{\kappa}$ = 1. (2) Let G be a graph with $\kappa = \bar{\kappa} = 1$ and $\Delta \leq p - 3$. By the definition of κ , G is connected and has a cutpoint v. We see that H = G - v has just two components, since otherwise every two points of \overline{G} would lie on a common cycle of \overline{G} and thus \overline{G} would have no cutpoint, contradicting $\overline{\kappa}$ = 1. Denote by H₁ and H_2 the two components of H, with p_1 and p_2 points respectively. Assume that both p_1 , $p_2 \ge 2$. Then \overline{G} would have no cutpoint since every two points of $ar{\mathsf{G}}$ would lie on a common cycle of $ar{\mathsf{G}}_{m{\bullet}}$. Thus it is sufficient to consider only a connected graph G which has a cutpoint with endline e and endpoint u. We now show that G - u contains a spanning complete bipartite subgraph. If G - u does not contain such a subgraph, then $\overline{G-u}$ is connected by Lemma 1. Moreover, the endpoint u of e is adjacent in \overline{G} to every point of \overline{G} lie on a common cycle and so $\overline{\mathsf{G}}$ has no cutpoint, which again contradicts $\overline{\mathsf{K}}$ = 1. Thus G - u contains a spanning complete bipartite subgraph.





Conversely, let G satisfy the condition (2) as shown in Figure 2a. Then \overline{G} is connected and the removal of the endpoint u from \overline{G} results in at least two components by Lemma 1. Hence we see that $\kappa = \overline{\kappa} = 1$.

A graph G is a <u>block</u> if G is connected and has no cutpoint. From Theorem 1 and Lemma 1, we obtain two consequences whose proofs are ommitted or outlined.

COROLLARY 1a. If G is a block, then \overline{G} is also a block if and only if

(1) 2 \leq deg v \leq p - 3 for every point v of G, and

(2) G has no spanning complete bipartite subgraph.

CORLLARY 1b. A graph G with p points satisfies the condition λ = $\overline{\lambda}$

= 1 if and only if G is a connected graph with a bridge and Δ = p - 2.

PROOF. Let G be a graph with $\lambda = \overline{\lambda} = 1$. Then G satisfies the condition $\kappa = \overline{\kappa} = 1$ by the relation $\kappa \leq \lambda$. Hence the graph G satisfies either (1) or (2) of Theorem 1. It is clear that (2) cannot hold, since \overline{G} can possess an endline only if the spanning bipartite subgraph of G - u is a star, in which case $\Delta = p - 2$, and so (1) must obtain.

Conversely, if G is a graph with $\lambda = 1$ and $\Delta = p - 2$, then \overline{G} is connected and has an endline, that is, $\overline{\lambda} = 1$.

2. BIPARTITE GRAPHS AND OUTERPLANAR GRAPHS.

A graph G is a <u>forest</u> if G has no cycles. An <u>outerplanar graph</u> is planar and can be embedded in the plane so that all its points lie on the same face.

THOEREM 2. All the graphs G such that both G and \overline{G} are bipartite are: are shown in Figure 3.



PROOF. The number k of components of G is at most two, since otherwise $\bar{\rm G}$ would contain a triangle.

CASE 1: k = 2. Let G have components G_1 and G_2 . Both G_1 and G_2 are complete, since otherwise \overline{G} would contain a triangle. Furthermore, the order of each of the complete graphs G_1 and G_2 is at most two, since otherwise G would contain a triangle. Hence we obtain $G = \overline{K}_2$, $K_1 \cup K_2$ and $2K_2$.

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CASE 2: k = 1. Since G is bipartite, the point set of G can be partitioned into two subsets V_1 and V_2 such that every line of G joins V_1 with V_2 . The cardinalities of V_1 and V_2 are at most two, since otherwise \overline{G} would contain a triangle. Furthermore, each subgraph induced by any three points of G contains one or two lines. Hence we get $G = K_1$, K_2 , P_3 , P_4 , and C_4 .

COROLLARY 2a. All the graphs G such that both G and \overline{G} are forests are:

G = K_1 , K_2 , \overline{K}_2 , K_1 U K_2 , P_3 and P_4

We have determined in Theorem 2 all eight graphs such that both G and \overline{G} are bipartite, and note that for none of these graphs G is both G and \overline{G} have even cycles. We now show that for just two graphs G, both G and \overline{G} have an odd cycle.

THEOREM 3. The two self-complementary graphs of order 5, A and C₅, are the only G such that both G and \overline{G} have only odd cycles (Figure 4).

PROOF. If the number of points of G is at least 6, either G or \overline{G} contains C_4 since the ramsey number $r(C_4) = 6$. It is easily verified that the two self-complementary graphs of order 5, A and C_5 shown in Figure 4, are the only G such that both G and \overline{G} have odd cycles.

THEOREM 4. All the graphs G such that neither G nor \overline{G} are forests but both are outerplanar are the following 32 graphs:

(1) the two self-complementary graphs A and C_5 of order 5 (Figure 4), and

(2) the 15 graphs shown in Figure 5 and their complements.

THEOREM 5. Both G and \overline{G} are eulerian if and only if both are connected, p is odd, and G is eulerian.

Of course p must be odd so that the degree of each point in both G and \overline{G} is even. Lemma 1 already gives a simple condition for both G and \overline{G} to be connected. The result follows at once.







Figure 5.

REFERENCE

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