

FRAGMENTS OF COMPLETE EXTENSIONS OF PA  
AND McDOWELL-SPECKER'S THEOREM

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**Abstract.** We generalise Theorem 1.4 of [2] and prove that for every complete extension  $\mathbf{T}$  of  $\mathbf{PA}$  and any  $n \in \omega$  there exists a model for  $\Sigma_n$ -fragment of  $\mathbf{T}$  that is not extendable (that is, a model with no proper strong elementary end-extension.) This is accomplished using a model called  $\Sigma_n$ -atomic. This result can be interpreted as “McDowell-Specker's Theorem does not hold for  $\Sigma_n$ -fragments of  $\mathbf{PA}$ ”.

**Basic definitions and notation.** The notation is the same as in [2].  $\mathbf{PA}$  stands for the axiom system of Peano arithmetic (e.g. as described in [1, p. 40]). A formula is  $\Sigma_n$  ( $\Pi_n$ ) iff the string of quantifiers in one of its prenex normal forms, begins with  $\exists$  ( $\forall$ ), and has no more than  $n - 1$  quantifier alternations. A sentence is  $\Delta_n$  iff it is both  $\Sigma_n$  and  $\Pi_n$ .  $\mathbf{T}_n$  stands for the  $\Sigma_n$ -fragment of the theory  $\mathbf{T}$ , that is a theory consisting of all the consequences of  $\mathbf{T}$  that are  $\Sigma_n$  sentences.

The notations  $\mathfrak{A}$  ( $A$ ),  $\mathfrak{B}$  ( $B$ ),  $\dots$  denote models (their universes), and the notations  $\mathfrak{M}$  ( $M$ ) and  $\mathfrak{N}$  ( $N$ ) denote models of  $\mathbf{PA}$  (their universes.) The letters  $x, y, z, \dots$  denote variables, while the letters  $\mathbf{a}, \mathbf{b}, \mathbf{c}, \dots$  denote constants. For a model  $\mathfrak{A}$  of some language  $\mathcal{L}$ , the theory of  $\mathfrak{A}$  (denoted  $\text{Th}(\mathfrak{A})$ ) is the set of all the sentences  $\varphi$  of  $\mathcal{L}$  such that  $\mathfrak{A} \models \varphi$ . Models  $\mathfrak{A}$  and  $\mathfrak{B}$  are elementarily equivalent iff  $\text{Th}(\mathfrak{A}) = \text{Th}(\mathfrak{B})$  (denoted  $\mathfrak{A} \equiv \mathfrak{B}$ .) For some set of sentences  $\Phi$  (some model  $\mathfrak{A}$ ),  $\mathcal{L}_\Phi$  ( $\mathcal{L}_\mathfrak{A}$ ) denotes the language of  $\Phi$  (of  $\mathfrak{A}$ .) If  $z$  codes an ordered pair  $\langle x, y \rangle$  we write  $(z)_0$  for  $x$  and  $(z)_1$  for  $y$ .

*Definition 1.* A model  $\mathfrak{A}$  is said to be a  $\Sigma_n$ -elementary extension of a model  $\mathfrak{B}$  ( $\mathfrak{B} \prec_n \mathfrak{A}$ ) iff for any  $\Sigma_n$ -formula  $\varphi$  with  $m$  free variables and any  $m$ -tuple  $\mathbf{a} \in B^m$ ,

$$\mathfrak{A} \models \varphi(\mathbf{a}) \quad \text{iff} \quad \mathfrak{B} \models \varphi(\mathbf{a}).$$

A  $\Sigma_n$ -elementary extension is a  $\Sigma_n$ -elementary end-extension iff it is also an end-extension. We say that a complete theory  $\mathbf{T}$  contains some formula schemata iff  $\mathbf{T}$  contains it as a set of formulas.

In [2, Theorem 1.6] it is shown that for some complete theory  $\mathbf{T}$  with binary relational symbol  $\rho$  the existence of  $\omega_1$ -like models for a first-order theory  $\mathbf{T}$  and extendability of all (or any) countable models for  $\mathbf{T}$  are equivalent to the same first-order property, certain scheme denoted by  $\mathcal{R}^+$ . It consists of the following sentences:

- C1.**  $\forall x \exists y \neg \rho(x, y)$   
**C2.**  $\forall x \forall y \exists z (\rho(x, z) \wedge \rho(y, z)),$

and for all formulas  $\varphi(x, u)$  of  $\mathcal{L}$ :

- C3.**  $\forall v [\forall x \exists y \forall u (\rho(x, v) \rightarrow (\varphi(x, u) \rightarrow \rho(u, y))) \rightarrow$   
 $\exists y \forall x \forall u (\rho(x, v) \rightarrow (\varphi(x, u) \rightarrow \rho(u, y)))].$

These axioms first appeared in [3], where a proof is given that every countable model satisfying **C1**, **C2** and **C3** is extendable. In [2, Theorem 1.4] it is shown that for every  $n \in \omega$  the theory  $\mathbf{PA}_n$  does not contain the scheme  $\mathcal{R}^+$ . Actually, a somewhat stronger result is given — any  $\Sigma_n$ -fragment of True Arithmetic (that is,  $\text{Th}(\omega, S, +, \cdot, 0)$ ) does not contain  $\mathcal{R}^+$ .

In that proof a kind of definable ultraproduct is used. It is  $\Sigma_n$ -definable ultrapower of a model  $\mathfrak{M}$  of  $\mathbf{PA}$ , a model that consists of  $\Sigma_n$ -definable functions modulo some ultrafilter  $G$  of  $\Sigma_n$ -definable sets. This model is denoted by  $F_{\Sigma_n}(\mathfrak{M})/G$ , and for such models a variant of Fundamental Theorem for Ultraproducts holds, namely  $\mathfrak{M} \prec_n F_{\Sigma_n}(\mathfrak{M})/G$ .

We will prove that the scheme  $\mathcal{R}^+$  is not contained in any  $\Sigma_n$ -fragment of  $\mathbf{T}$  (from now on,  $\mathbf{T}$  stands for some (fixed) complete extension of  $\mathbf{PA}$ .) From this we have our

**MAIN THEOREM.** *For any theory  $\mathbf{T}_n$  ( $\mathbf{T}$  is some complete extension of  $\mathbf{PA}$ ) there is a model that is not extendable.*

Let  $\mathfrak{M}$  be a model for  $\mathbf{PA}$ . An element  $\mathbf{a} \in M$  is said to be  $\Sigma_n$ -definable in the model  $\mathfrak{M}$  iff there exists a  $\Sigma_n$ -formula  $\varphi_{\mathbf{a}}$  of  $\mathcal{L}_{\mathbf{PA}}$  such that the following holds:

$$\mathfrak{M} \models \varphi_{\mathbf{a}}(\mathbf{a}) \wedge \forall x (\varphi_{\mathbf{a}}(x) \rightarrow x = \mathbf{a})$$

(We will usually say “ $\Sigma_n$ -definable” instead of “ $\Sigma_n$ -definable in a model  $\mathfrak{M}$ ” when no ambiguity occurs.)

By  $\Sigma_n^{\mathfrak{M}}$  we denote a countable submodel of  $\mathfrak{M}$  that consists of exactly those elements  $\mathbf{a}$  that are  $\Sigma_n$ -definable in  $\mathfrak{M}$ . It is easily verified that  $\Sigma_n^{\mathfrak{M}}$  is closed under the operations  $+$  and  $\cdot$ .

**LEMMA 1.** *For any model  $\mathfrak{M} \models \mathbf{PA}$  the following holds:*

$$\mathfrak{M} \prec_n \Sigma_n^{\mathfrak{M}}$$

so  $\Sigma_n^{\mathfrak{M}} \models \text{Th}_n(\mathfrak{M})$ .

*Proof.* Suppose that  $\varphi(x, y)$  is  $\Sigma_n$ -formula and  $\mathbf{a} \in \Sigma_n^{\mathfrak{M}}$  is such that

$$\mathfrak{M} \models \exists x \varphi(x, \mathbf{a}).$$

It is enough to show that there is a  $\Sigma_n$ -definable  $\mathbf{b} \in M$  such that  $\mathfrak{M} \models \varphi(\mathbf{b}, \mathbf{a})$ . Note that for some  $\Sigma_n$ -formula  $\varphi_{\mathbf{a}}(x)$  the following holds

$$\mathfrak{M} \models \exists x \exists y (\varphi(x, y) \wedge \varphi_{\mathbf{a}}(y))$$

If we encode the pair  $\langle x, y \rangle$  by  $z$ , then the formula above becomes:

$$\mathfrak{M} \models \exists z (\varphi((z)_0, (z)_1) \wedge \varphi_{\mathbf{a}}((z)_1))$$

It is clear that the formula  $\psi(z)$  defined as

$$\varphi((z)_0, (z)_1) \wedge \forall x < z \neg \varphi((x)_0, (x)_1)$$

is again  $\Sigma_n$ , and that  $\mathfrak{M} \models \exists y \psi(y)$ . So there is some  $\Sigma_n$ -definable  $\mathbf{c} \in M$  such that  $\mathfrak{M} \models \psi(\mathbf{c})$  and  $\mathfrak{M} \models \varphi((\mathbf{c})_1, \mathbf{a})$ , and we just set  $\mathbf{b} = (\mathbf{c})_1$ .  $\square$

*Remark.* In  $\Sigma_n^{\mathfrak{M}}$  every element is  $\Sigma_n$ -definable, so we may say that  $\Sigma_n^{\mathfrak{M}}$  is  $\Sigma_n$ -atomic. It can easily be shown that this model is also  $\Sigma_n$ -prime (that is,  $\Sigma_n$ -elementarily embeddable in every model for  $\mathbf{T}_n$ ), so this construction might be of interest in its own right.

LEMMA 2 (cf. [2, Lemma 1.1]). *For any  $n \in \omega$  there exists a  $\Delta_{n+1}$ -formula  $\varphi(x, y)$  in  $\mathcal{L}_{\mathbf{PA}}$  and a model  $\mathfrak{M}_1 \models \mathbf{T}_n \cup \{\neg \mathcal{R}(\varphi)\}$ .*

*Proof.* Let  $\mathfrak{M}$  be a model for  $\mathbf{T}$ . By Lemma 1 we have  $\Sigma_n^{\mathfrak{M}} \models \mathbf{T}_n$ . Let  $G$  be a nonprincipal ultrafilter in  $\mathcal{D}_n(\Sigma_n^{\mathfrak{M}})$ , the set of  $\Sigma_n$ -definable subsets of  $\Sigma_n^{\mathfrak{M}}$ . Now we construct a model  $\mathfrak{M}_1 = F_{\Sigma_n}(\Sigma_n^{\mathfrak{M}})/G$ . By Lemma 1 this is also a model for  $\mathbf{T}_n$ . Fix some  $\mathbf{b} \in M_1$ . It is a  $=_G$ -equivalence class of some function  $f$  that is  $\Sigma_n$ -definable without parameters (remember that  $\Sigma_n^{\mathfrak{M}}$  is  $\Sigma_n$ -atomic) in  $\Sigma_n^{\mathfrak{M}}$ . So we have a  $\Sigma_n$ -formula  $\psi_f(x, y)$  such that  $fm = n$  iff  $\Sigma_n^{\mathfrak{M}} \models \psi_f(m, n)$  for all  $m, n \in \Sigma_n^{\mathfrak{M}}$ . And now,

$$\begin{aligned} B &= \{n \in \Sigma_n^{\mathfrak{M}} \mid \Sigma_n^{\mathfrak{M}} \models \text{SAT}_{\Sigma_n}(\ulcorner \theta \urcorner, n, fn)\} \\ &= \{n \in \Sigma_n^{\mathfrak{M}} \mid \Sigma_n^{\mathfrak{M}} \models \theta(n, fn)\} \\ &= \Sigma_n^{\mathfrak{M}}, \end{aligned}$$

and  $\mathfrak{M}_1 \models \text{SAT}_{\Sigma_n}(\ulcorner \theta \urcorner, i_G, \mathbf{b})$ , where  $i_G$  stands for the  $=_G$ -equivalence class of the diagonal  $i$  of  $\Sigma_n^{\mathfrak{M}}$ . We conclude that for every  $\mathbf{b}$  in  $M_1$  there exists some  $\mathbf{e} \in \omega$  such that

$$\mathfrak{M}_1 \models \text{SAT}_{\Sigma_n}(\mathbf{e}, i_G, \mathbf{b}).$$

It remains to show that wanted  $\Delta_{n+1}$  formula is  $\psi(x, y)$ , defined as

$$\text{SAT}_{\Sigma_n}(x, i_G, y) \wedge (\forall z < x) \neg \text{SAT}_{\Sigma_n}(z, i_G, y),$$

“ $x$  is the least Gödel’s number of a formula that defines  $y$ ”. From the previous discussion it is evident that  $\mathfrak{M}_1 \models \psi(x, y)$  only if  $x$  is standard, and that the set of all  $x \in \omega$  such that  $\mathfrak{M}_1 \models \exists y \psi(x, y)$  is cofinal in  $\omega$ . Now we check that  $\psi$  is not regular in  $\mathfrak{M}_1$ , i.e. that the following holds:

$$\begin{aligned} \mathfrak{M} \models \exists v [\forall x \exists y \forall u (x < v \rightarrow (\phi(x, u) \rightarrow u < y)) \wedge \\ \forall y \exists x \exists u (x < v \wedge \phi(x, u) \wedge y \leq u)]. \end{aligned}$$

For  $v$  we fix some nonstandard element  $\mathbf{v}$  of  $M_1$ . To prove the first part of the statement, fix any  $\mathbf{x} < \mathbf{v}$ . The set  $\{u \in M_1 \mid \mathfrak{M}_1 \models \psi(\mathbf{x}, u)\}$  has at most one element, thus it is bounded by some  $\mathbf{y}$ . To check the second part, note that for any  $\mathbf{y} \in M_1$  there is an  $\mathbf{x} \in \omega$  (thus  $\mathbf{x} < \mathbf{v}$ ) and  $\mathbf{u} > \mathbf{y}$  such that  $\mathfrak{M}_1 \models \psi(\mathbf{x}, \mathbf{u})$ .  $\square$

Note that the minor modification of the proof that  $\psi$  is not regular in  $\mathfrak{M}_1$  gives the following semantical characterization of **C3** for every model  $\mathfrak{M}$  with built-in Skolem functions:

$\mathfrak{M} \models \mathbf{C3}$  iff there is no definable (in  $\mathfrak{M}$ ) function mapping a bounded subset of  $M$  cofinally into  $M$ .

*Proof of the Main Theorem.* The model  $\mathfrak{M}_1$  from Lemma 2 is not extendable.

$\square$

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