



ENZ für Algebra

On Plethysms of type $s_\mu[ts_\nu]$ for non-integral t and a few Applications

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Integer partitions, Young diagrams / tableaux

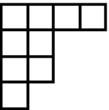


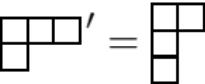
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$$\lambda = \lambda_1 + \lambda_2 + \dots + \lambda_l$$

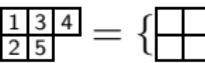
Let $\lambda \vdash n$ be a partition of the integer n

- ▶ $|\lambda| = n$ weight of the partition
- ▶ $\ell(\lambda) = l$ number of parts, length of the partition
- ▶ $\lambda = (\lambda_1, \dots, \lambda_l) = [1^{m_1}, 2^{m_2}, \dots, n^{m_n}]$
additive and multiplicative representation

- ▶  shape or (Ferrers) Young diagram of a partition

- ▶  $\lambda' =$ conjugation of partitions

- ▶  (semi) standard Young tableau, box $x = (i,j) \in SYT$
 $\begin{matrix} 1 & 2 & 3 & 4 \\ 5 & 7 \\ 6 & 8 \\ 9 \end{matrix}$
(weakly) increasing rows, strictly increasing columns

- ▶  $\{ \text{Young diagrams} \}$



$\mathcal{H}(G)$ Hopf algebra (groups)

Group

- ▶ $f, g, h, e \in G$ group elements
- ▶ $f(gh) = (fg)h$
- ▶ $fg = h, \quad ge = g = eg, \quad g^{-1}$ exists with $g^{-1}g = e$

Representative functions

- ▶ $\phi, \psi, \mathbf{1} \in \mathcal{H}(G) = \{G \rightarrow \mathbb{C}\}$ functions
- ▶ $(\psi\phi)(g) = \psi(g)\phi(g)$ pointwise product (assoc.)
- ▶ $\mathbf{1}(g) = 1$ unit element

Maps:

$$\Delta : \mathcal{H}(G) \rightarrow \mathcal{H}(G) \otimes \mathcal{H}(G); \quad \epsilon : \mathcal{H}(G) \rightarrow \mathbb{C}; \quad S : \mathcal{H}(G) \rightarrow \mathcal{H}(G)$$

- ▶ Coproduct: $\Delta(\psi)(g, h) = \psi(g)\psi(h)$
- ▶ Counit: $\epsilon(\psi) = \psi(e)$
- ▶ Antipode: $S(\psi)(g) = \psi(g^{-1})$

$\mathcal{H}(G)$ Hopf algebra, cont.



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Hopf algebra; CD version (counit; pentagon; antipode)

$$\begin{array}{ccc}
 \begin{array}{ccc}
 \mathcal{H} & \xrightarrow{\Delta} & \mathcal{H} \otimes \mathcal{H} \\
 \downarrow \Delta & \searrow \cong & \downarrow \epsilon \otimes 1 \\
 \mathcal{H} \otimes \mathcal{H} & \xrightarrow{1 \otimes \epsilon} & \mathcal{H}
 \end{array}
 &
 \begin{array}{ccc}
 \mathcal{H} \otimes \mathcal{H} & \xrightarrow{\Delta \otimes \Delta} & \mathcal{H} \otimes \mathcal{H} \otimes \mathcal{H} \otimes \mathcal{H} \\
 \downarrow m & & \downarrow 1 \otimes \sigma \otimes 1 \\
 \mathcal{H} & \xrightarrow{\Delta} & \mathcal{H} \otimes \mathcal{H} \otimes \mathcal{H} \otimes \mathcal{H} \\
 & & \searrow m \otimes m \\
 & & \mathcal{H} \otimes \mathcal{H}
 \end{array}
 &
 \begin{array}{c}
 \sigma(\psi \otimes \phi) = \phi \otimes \psi \\
 \eta : \mathbb{C} \rightarrow \mathcal{H}
 \end{array}
 \end{array}$$

Diagram illustrating the structure of a Hopf algebra \mathcal{H} using the CD version (counit, pentagon, antipode).

- Top Left:** A commutative diagram showing the counit ϵ . The top row has \mathcal{H} on the left and $\mathcal{H} \otimes \mathcal{H}$ on the right. The bottom row has $\mathcal{H} \otimes \mathcal{H}$ on the left and \mathcal{H} on the right. The top arrow is Δ , the bottom arrow is $1 \otimes \epsilon$, and the right arrow is $\epsilon \otimes 1$. A diagonal arrow from \mathcal{H} to $\mathcal{H} \otimes \mathcal{H}$ is labeled \cong .
- Top Middle:** A commutative diagram showing the pentagon equation. The top row has $\mathcal{H} \otimes \mathcal{H}$ on the left and $\mathcal{H} \otimes \mathcal{H} \otimes \mathcal{H} \otimes \mathcal{H}$ on the right. The bottom row has $\mathcal{H} \otimes \mathcal{H} \otimes \mathcal{H} \otimes \mathcal{H}$ on the left and $\mathcal{H} \otimes \mathcal{H}$ on the right. The top arrow is $\Delta \otimes \Delta$, the bottom arrow is m , and the right arrow is $1 \otimes \sigma \otimes 1$. A diagonal arrow from \mathcal{H} to $\mathcal{H} \otimes \mathcal{H} \otimes \mathcal{H} \otimes \mathcal{H}$ is labeled Δ .
- Bottom Left:** A commutative diagram showing the antipode S . The top row has \mathcal{H} on the left and $\mathcal{H} \otimes \mathcal{H}$ on the right. The bottom row has $\mathcal{H} \otimes \mathcal{H}$ on the left and \mathcal{H} on the right. The top arrow is Δ , the bottom arrow is $1 \otimes S$, and the right arrow is m . A diagonal arrow from \mathcal{H} to $\mathcal{H} \otimes \mathcal{H}$ is labeled ϵ . A diagonal arrow from $\mathcal{H} \otimes \mathcal{H}$ to \mathbb{C} is labeled $S \otimes 1$.
- Bottom Right:** Equations involving the antipode S and the counit ϵ . The top equation is $\sigma(\psi \otimes \phi) = \phi \otimes \psi$. The bottom equation is $\eta : \mathbb{C} \rightarrow \mathcal{H}$.

Symm(X) Hopf algebra (outer HA)



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Symm(X)

- ▶ $\text{Symm}(X)$ is the sub-Hopf algebra $K_0[\text{Rep}(GL(\infty))]$ (i.e. Grothendieck ring) of the repr. category $\text{Rep}(GL(\infty))$ (in formally infinite many variables to avoid syzygies).

Let A, B, C ($A[X], B[X], C[X]$) be in $\text{Symm}(X)$:

- ▶ product: $AB = C; \quad s_\mu s_\nu = \sum c_{\mu,\nu}^\lambda s_\lambda \quad (c_{\mu,\nu}^\lambda \text{ LR-coeff.})$
- ▶ unit: $A1 = A = 1A; \quad s_\mu s_{(0)} = s_\mu = s_{(0)} s_\mu$
- ▶ coproduct: (Sweedler index notation)
$$\Delta(A[X]) = A[X + Y] = A_{(1)}[X]A_{(2)}[Y] \cong A_{(1)} \otimes A_{(2)}$$
$$\Delta s_\lambda = \sum_{(\lambda)} s_{\lambda_{(1)}} \otimes s_{\lambda_{(2)}} = \sum_{\mu,\nu} c_{\mu,\nu}^\lambda s_\mu \otimes s_\nu \quad (c_{\mu,\nu}^\lambda \text{ LR-coeff.})$$
- ▶ counit: $\epsilon(A) = \delta_{A,1}; \quad \epsilon(s_\mu) = \delta_{\mu,(0)}$

Needs of course a proof to be a HA (Geissiger 77, Zelevinsky 81)
 $\text{Symm}(X)$ is a PSH (positive selfadjoint HA)

Symm(X) additional structure



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inner product (S_n tensor product of characters)

- ▶ tensor product in $R(S_n)$ yields via Frobenius characteristics a second product on $\text{Symm}(X)$
- ▶ $s_\mu \star s_\nu = \sum_\lambda g_{\mu,\nu}^\lambda s_\lambda$
- ▶ $g_{\mu,\nu}^\lambda$ is symmetric in all three partitions, computable via the Murnaghan-Nakayama rule
- ▶ $\delta(s_\lambda[X]) = s_\lambda[X \otimes Y] = \sum_{[\lambda]} s_{\lambda_{[1]}}[X] s_{\lambda_{[2]}}[Y] \cong \sum_{[\lambda]} s_{\lambda_{[1]}} \otimes s_{\lambda_{[2]}}$

Scalar product (self duality)

$\{s_\lambda\}_{\lambda \vdash n, n \in \mathbb{N}}$ basis of irreducibles \Rightarrow orthonormal (Schur's lemma)

$$\langle \cdot | \cdot \rangle : \Lambda \otimes \Lambda \rightarrow \mathbb{Z} \quad \langle s_\mu | s_\nu \rangle = \delta_{\mu,\nu}$$

Symm(X) additional structure, cont.



Schur functors (on $R(G) = \text{Rep}(G)$)

Let V, W be G -modules of $R(G)$. Consider a the Schur functor
 $\mathbb{S}_\lambda : R(G) \rightarrow R(G); \quad \mathbb{S}_\lambda(V) = W =: V^\lambda \quad (\text{Weyl module})$

$$\mathbb{S}_{(n)}(V) = \text{Sym}(V^{\otimes n}) = V^{(n)}$$

$$\mathbb{S}_{(1^n)}(V) = V^{\wedge n} = V^{(1^n)}$$

$$\mathbb{S}_\mu \circ \mathbb{S}_\nu(V) = \mathbb{S}_\mu(V^\nu) = (V^\nu)^\mu \qquad U \xrightarrow{\mathbb{S}_\nu} V \xrightarrow[\mathbb{S}_\mu \circ \mathbb{S}_\nu]{\mathbb{S}_\mu} W$$

Plethysm (composition on Symm(X) $\cong K_0[\text{Rep}(G)]$)

$$s_\mu[s_\nu] = \sum_{\lambda} a_{\mu,\nu}^\lambda s_\lambda$$

$$s_\nu[X] = Y = y_1[X] + y_2[X] + \dots + y_m[X]$$

with m number of monomials (incl. multiplicities)

$$\mathbb{S}_\nu(V) = V^\nu = W \quad \text{and} \quad \mathbb{S}_\mu(\mathbb{S}_\nu(V)) = \mathbb{S}_\mu(W) = W^\mu = (V^\nu)^\mu$$

$$s_\mu[s_\nu][X] = s_\mu[s_\nu[X]] = s_\mu[Y[X]]$$



Symm(X) : summary of costructures

σ (succ. map), HA (outer HA), BI (bialgebra), HG (Hopf gebra)

$$\begin{aligned}\sigma : \quad \sigma(s_\lambda)[X] &= s_\lambda[x_0 + X] = \sum_{(\eta), \nu} c_{(\eta), \nu}^\lambda s_{(\eta)}[x_0] s_\nu[X] \\ &= \sum_{\lambda \subset \nu} x_0^n s_\nu \text{ with } \nu \setminus \lambda \text{ horizontal strip with } n \text{ boxes}\end{aligned}$$

$$\text{HA: } \Delta(s_\lambda)[X] = s_\lambda[X + Y] = s_{\lambda_{(1)}}[X] s_{\lambda_{(2)}}[Y] = \sum_{\mu, \nu} c_{\mu, \nu}^\lambda s_\mu[X] s_\nu[Y]$$

$$\text{BI: } \delta(s_\lambda)[X] = s_\lambda[XY] = s_{\lambda_{[1]}}[X] s_{\lambda_{[2]}}[Y] = \sum_{\mu, \nu} g_{\mu, \nu}^\lambda s_\mu[X] s_\nu[Y]$$

$$\text{HG: } \hat{\Delta}(s_\lambda)[X] = s_\lambda[Y^X] = s_{\lambda_{\langle 1 \rangle}}[X] s_{\lambda_{\langle 2 \rangle}}[Y] = \sum_{\mu, \nu} p_{\mu, \nu}^\lambda s_\mu[X] s_\nu[Y]$$

I call this **Arithmetic Complexity**



Plethysm properties

Let $a, b, c, \dots \in R$ and $A, B, C, \dots \in \text{Symm}[X]$

$\varepsilon^n(s_\lambda(x_1, x_2, \dots)) = s_\lambda(1, \dots, 1, 0, \dots)$ with n ones

defining rules :

$$(aA + bB)[C] = aA[C] + bB[C] \quad \text{left distributivity}$$

$$(AB)[C] = A[C]B[C] \quad \text{left multiplicativity}$$

$$A[B + C] = A_{(1)}[B]A_{(2)}[C] \quad \text{right additivity}$$

('+' convolution on $\text{Symm}(X)$)

$$A[-B] = (\text{S}(A))[B] \quad \text{antipode (negative alphabets)}$$

$$A[BC] = A_{[1]}[B]A_{[2]}[C] \quad \text{right multiplicativity}$$

('.' convolution on $\text{Symm}(X)$)

$$A[B * C] = A[\sum_{i=1}^N g_{BC}^{D_i} D_i]$$

$$= \prod_{i=1}^N \varepsilon^{g_{BC}^{D_i}} (A_{(i)[1]}) A_{(i)[2]}[D_i] \quad \text{right inner product expansion}$$

$$(A[B])[C] = A[B[C]] \quad \text{associativity}$$

Relating outer and inner coproducts '+'



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We give two ways to compute the plethysm $s_\mu[n s_\nu]$:

- ▶ Lemma [FJK] : $(A \otimes B)[C] = A[C] \otimes B[C]$
- ▶ Lemma : $A[B + B] = A_{(1)}[B]A_{(2)}[B] = (m^1 \circ \Delta^1)(A))[B]$
 $: A[n B] = \prod_{i=1}^n A_{(i)}[B] = (m^{(n-1)} \circ \Delta^{(n-1)})(A))[B]$

Proof: additive version

$$\begin{aligned}s_\mu[n s_\nu] &= s_\mu[s_\nu + (n-1)s_\nu] && n > 1 \\&= s_{\mu(1)}[s_\nu] s_{\mu(2)}[(n-1)s_\nu] \\&= \dots \\&= \prod_{i=1}^n s_{\mu(i)}[s_\nu] \\&= \left(m^{(n-1)} \circ \Delta^{(n-1)}\right)(s_\mu)[s_\nu]\end{aligned}$$

To stop the induction process we have to require: $n \in \mathbb{N}$



Relating outer and inner coproducts !!

- ▶ $\varepsilon^n(s_{(1)}) = 1 + \dots + 1 = \dim_{GL(n)}(s_{(1)})$ is a Hom
- ▶ Lemma [FJKW] : $\Delta(A[B]) = A[\Delta(B)]$
- ▶ with $m(A \otimes B) = AB$ and the HA axioms we get further
Lemma [FJK] : $A[B \otimes C] = A_{[1]}[B] \otimes A_{[2]}[C]$

2nd proof: multiplicative version

$$\begin{aligned}s_\mu[n s_\nu] &= m(s_\mu[n \otimes s_\nu]) \\&= m(s_\mu[\varepsilon^n(s_{(1)}) \otimes s_\nu]) \\&= m \circ (\varepsilon^n \otimes \text{Id})(s_\mu[s_{(1)} \otimes s_\nu]) \\&= m(\varepsilon^n(s_{\mu[1]}) \otimes s_{\mu[2]}[s_\nu]) \\&= \varepsilon^n(s_{\mu[1]}) s_{\mu[2]}[s_\nu] \\&= \dim_{GL(n)}(s_{\mu[1]}) s_{\mu[2]}[s_\nu]\end{aligned}$$

Relating outer and inner coproducts Δ, δ



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- ▶ $A = A[s_{(1)}]$ for all A
- ▶ $[1 + n](A) = (m^n \Delta^n)(A) = A[(1 + n)s_{(1)}]$
 $[1 + n](A[B]) = (m^n \Delta^n)(A[B]) = A[(1 + n)B]$
 $[1 + n]$ is a scaling operator
- ▶ $[1 + n](A[B]) = A[(1 + n)B] = \dim_{GL(n+1)}(A_{[1]})A_{[2]}[B]$
 $[1 + n] = (m^n \Delta^n) = (\varepsilon^{(n+1)} \otimes \text{Id})\delta$

Result : t -generalization

Since we have that ε^n is a binomial expression in n it can be generalized to arbitrary $t \in \mathbb{C}$, that is $\varepsilon^n \Rightarrow \varepsilon^t$:

$$[1 + t] := (\varepsilon^{1+t} \otimes \text{Id})\delta \quad (= (m^t \Delta^t) \quad t \in \mathbb{N})$$
$$[1 + t](A[B]) = A[(1 + t)B] = \dim_{GL(t+1)}(A_{[1]})A_{[2]}[B]$$

For $t \notin \mathbb{N}$ generically non-recursive.

Other Hom's and specializations



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Specializations (eg. q, t -case : no time ... → [sketch])

- ▶ Let $\Omega[X] = M[X] = \sum h_n(X)$ complete symmetric functions
 $[1+t]\Omega[X] = \Omega[(1+t)X]$ well known plethystic formula
- ▶ $\langle A | B \rangle = \varepsilon^1(A * B)$: counts 1-dim. GL irreps
 $\varepsilon^1(s_\lambda) = \delta_{\lambda,(n)}$: shows that the series $M = \Omega$ is defined by ε^1

duality : $\Lambda \xrightarrow{\varepsilon^1} \mathbb{Z}$ $\Lambda \xrightarrow{\varepsilon^{1+t}} \mathbb{Z}$

$$\begin{array}{ccc} & \xrightarrow{\varepsilon^1} & \\ \Lambda & \xleftarrow{M} & \mathbb{Z} \\ & \xleftarrow{\varepsilon^{1+t}} & \\ & & \mathbb{Z} \\ & \xleftarrow{M^{1+t}} & \end{array}$$

- ▶ $\delta(M[X]) = M[XY] = \sum_\lambda s_\lambda[X] s_\lambda[Y]$: Cauchy kernel
 $\delta(M[-X]) = L[XY] = \sum_\lambda (-1)^{|\lambda|} s_\lambda[X] s_{\lambda'}[Y]$: Cauchy-Binet
- ▶ There are abundantly many further specializations and homomorphisms available to generalize the above given result:
 $M[X] \rightarrow M[\frac{1}{1-q}X] = M[\sum_{n \geq 0} q^n X] = \prod_{n \geq 0} M[q^n X] =$
 $\prod_{n \geq 0} (\sum_\lambda \dim_{GL(q^n)}(s_\lambda) s_\lambda)$
 $A[B] \rightarrow A[\frac{1}{1-q}B] = \prod_{i \geq 1} (\dim_{GL(q^n)}(A_{(i)[1]}) A_{(i)[2]}[B])$
for a general A and a general “alphabet” B .



Applications:

Computational benefits

Timing of the iterated and directly evaluated plethysms
 $s_{(3)}[n s_{(1,1)}]$:

multiplicity	recursive	direct
n=1	0.01	0.02
n=10	0.08	0.02
n=100	0.89	0.01
n=1000	7.32	0.01
n=10000	—	0.01

SchurFkt Maple package [BF.,Rafal Ablamowicz]: Figures obscured by Maple's garbage collection.

SCHUR cannot do the $n = 100, 1000$ cases in reasonable time
(Maple algorithm uses remember hash tables)

Applications:



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Jack symmetric function (eg. [Brenti98])

Define the scalar product

$$\langle A | B \rangle_{\alpha\beta} := \langle A[\alpha s_{(1)}] | B[\beta s_{(1)}] \rangle = \varepsilon^{\alpha\beta} (A * B) = (A * B)[1, \dots, 1, 0, \dots]$$

Schur functions are no longer an orthonormal basis for $\langle \cdot | \cdot \rangle_a$

$$\left(\begin{array}{c|cc} \langle \cdot | \cdot \rangle_{q^2} & s_{(2)} & s_{(11)} \\ \hline s_{(2)} & \frac{1}{2}q^2(q^2 + 1) & \frac{1}{2}q^2(q^2 - 1) \\ s_{(11)} & \frac{1}{2}q^2(q^2 - 1) & \frac{1}{2}q^2(q^2 + 1) \end{array} \right)$$

Jack symmetric functions can be defined as:

$$J_\lambda(x; \alpha) = \prod_{i=1}^{\ell(\lambda)} J_{\lambda_i}(x; \alpha) \quad J_{\lambda_i}(x; \alpha) = \sum_{\mu \vdash n} \alpha^{n-\ell(\mu)} n! \frac{p_\mu}{z_\mu}$$

$$\text{which results in: } h_\lambda\left[\frac{1}{\alpha}x\right] = \frac{1}{\alpha^{|\lambda|}\lambda!} J_\lambda(x; \alpha)$$

Applications:



Powers of Schur function series

let G be a group like Schur function series (M, L, \dots) i.e.
 $\Delta G = G \otimes G$. For such series we have:

$$G[n s_{(1)}] = G^n = GG \cdots G$$

and of course with $G^+ := G - s_{(0)}$

$$\log(1 + G^+) = - \sum_{n \geq 1} \frac{(-1)^n}{n} (G^+)^n = - \sum_{n \geq 1} \sum_{k=0}^n \frac{(-1)^k}{n} \binom{n}{k} (G^+)^k$$

$$\log(1 + M^+) = p_{(1)} + \frac{1}{2}p_2 + \frac{1}{3}p_3 + \frac{1}{4}p_4 + \dots$$

which can be generalized

Applications:



Powers of Schur function series (cont.)

We look at $D = M[s_{(2)}]$, $C = M[-s_{(2)}]$, $M_3 = M[s_{(3)}]$ etc.

$$\log(1 + D^+) = \frac{1}{2}(p_2 + p_{(11)}) + \frac{1}{4}(p_4 + p_{(22)}) \dots + \frac{1}{2n}(p_{2n} + p_{(nn)}) \dots$$

$$\log(1 + C^+) = -\frac{1}{2}(p_2 + p_{(11)}) - \frac{1}{4}(p_4 + p_{(22)}) \dots - \frac{1}{2n}(p_{2n} + p_{(nn)}) \dots$$

$$\log(1 + M_3^+) = \frac{1}{3}(p_3 + p_{(111)}) + \frac{1}{6}(p_6 + p_{(222)}) \dots + \frac{1}{3n}(p_{3n} + p_{(nnn)}) \dots$$

Of course : $\log(1 + C^+) + \log(1 + D^+) = 0 \Leftrightarrow CD = 1$

Furthermore we have with $\Theta := \sum_{\lambda} s_{\lambda} s_{\lambda}$ and $\Theta \tilde{\Theta} = 1$ that

$$D^n = D[n s_{(1)}] \tilde{\Theta}^{n-1}$$

$$[1 + n]D = (m^n \Delta^n)D = D[(1 + n) s_{(1)}] \Theta^n$$

a relation important for string functions and QFT.

Thank you!

Thanks to the organizers !



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Peano axioms (NNO : natural number objects)



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Mathematical induction (adding '+1')

$\sigma : \mathbb{N} \rightarrow \mathbb{N}$ successor map, R target ring of series coeff,

$a : \mathbb{N} \rightarrow R$ series (actual a ring, dualizing gives co-structures):

$$\begin{array}{ccc} \mathbb{N} & \xrightarrow{\sigma} & \mathbb{N} \\ \downarrow a \quad \swarrow 0 & & \downarrow a \\ R & \xrightarrow{\sigma} & R \end{array} \qquad \begin{array}{ccc} n & \xrightarrow{\sigma} & 1 + n \\ \downarrow a \quad \swarrow 0 & & \downarrow a \\ a_n & \xrightarrow{\sigma} & a_{n+1} \end{array}$$

Arithmetic complexity (of ordinary arithmetics)

- ▶ $\sigma(n) = 1 + n$ successor map
- ▶ $\sigma^m(n) = 1 + (1 + \dots (1 + n) \dots) =: \text{add}(m, n) = m + n$
iteration of σ : addition
- ▶ $\text{add}(m, 0)^n = (\sigma^m(0))^n = \sigma^{mn}(0) =: \text{mul}(m, n) = mn$
iteration of $\text{add}(n,)$: multiplication

Peano axioms (NNO : natural number objects)



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Arithmetic complexity (of ordinary arithmetics)

$$\begin{aligned} \blacktriangleright \text{ mul}(m, 1)^n &= \text{mul}(m, \sigma(0))^n = (\text{add}(m, 0)^{\sigma(0)})^n \\ &= \prod_{i=1}^n \text{add}(m, 0)^{\sigma(0)} = \prod_{i=1}^n (\sigma^m(0))^{\sigma(0)} \\ &= \prod_{i=1}^n (\sigma^m(0)) = \sigma^{m^n}(0) \end{aligned}$$

Peano axioms (NNO : natural number objects)



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- ▶ $\sigma(n) = 1 + n$ successor map
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iteration of σ : addition
- ▶ $\text{add}(m, 0)^n = (\sigma^m(0))^n = \sigma^{mn}(0) =: \text{mul}(m, n) = mn$
iteration of $\text{add}(n,)$: multiplication
- ▶ $\text{mul}(m, 1)^n = \text{mul}(m, 1)^n =: \text{exp}(m, n) = m^n$
iteration of $\text{mul}(n,)$: exponentiation