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THE NUMBER OF RIBBON SCHUR FUNCTIONS

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Abstract. We present formulas for the number of distinct ribbon Schur functions of a given size, and of a given size and height.

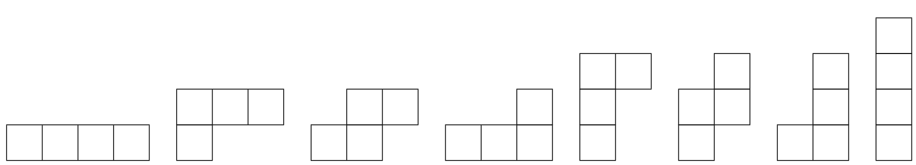
1. Introduction

An important basis for the space of homogeneous symmetric functions of degree n is the set of *Schur functions* s_λ, indexed by partitions λ of n. A larger set of homogeneous symmetric functions of degree n is the set of *skew Schur functions* $s_{\lambda/\mu}$. These are indexed by skew shapes λ/μ of size n, i.e., pairs of partitions (λ, μ) such that

- $\lambda = (\lambda_1 \geq \lambda_2 \geq \cdots \geq \lambda_k > 0)$ is a partition of $n + m$,
- $\mu = (\mu_1 \geq \mu_2 \geq \cdots \geq \mu_\ell > 0)$ is a partition of m,
- k, the number of parts of λ, is strictly larger than ℓ, the number of parts of μ,
- and $\mu_i \leq \lambda_i$ for $i \leq \ell$.

When μ is the empty partition, i.e., $\ell = 0$, $s_{\lambda/\mu} = s_\lambda$. Since the set of Schur functions is a basis, there must be relations between skew Schur functions. Equalities between skew functions have been studied by Stephanie van Willigenburg, Peter McNamara, Vic Reiner and Kristin Shaw [8, 5, 4] and, in the multiplicity-free case, by Christian Gutschwager [2]. However, so far only partial results and a conjecture are available.

The situation is very different for the subset of *ribbon Schur functions*, that are indexed by *ribbons*, also known as rim hooks or border strips. These are skew shapes that satisfy $\lambda_{i+1} = \mu_i + 1$ for $1 \leq i < k$, where we set $\mu_i = 0$ for $i > \ell$. Here are the ribbons of size 4:



It can be shown that the space of homogeneous symmetric functions of degree n is also generated by the set of ribbon Schur functions of size n. For these functions, Louis J. Billera, Hugh Thomas, and Stephanie van Willigenburg [1] give a criterion for deciding when they are equal. In this article we use this criterion to count the number of distinct ribbon Schur functions of a given size, see Theorem 3.4. Furthermore, letting the *height* of a ribbon Schur function be the number of parts of

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λ minus one, we also find the generating function for the number of distinct ribbon Schur functions of a given size by height, see Theorem 4.2.

Note that ribbons λ/μ of size n and height $m-1$ can be identified with compositions α of size n and length m by setting $\alpha_i = \lambda_i - \mu_i$ for all i. Two compositions α and β are called equivalent, denoted $\alpha \sim \beta$, if and only if the corresponding ribbon Schur functions are equal.

In the following section we recall a binary operation on compositions from [1], that makes the set of compositions into a monoid with (almost) unique factorisation. One of the main theorems of [1] shows that equivalence of compositions is easily determined given their factorisations.

In Section 3 we present a relatively appealing formula for the number of distinct ribbon Schur functions of a given size, while in Section 4 we exhibit a (not nearly as beautiful) formula for the number of distinct ribbon Schur functions of a given size and height. For more information on symmetric functions we refer to Chapter 7 of Enumerative Combinatorics 2 [7].

2. Composition of Compositions and equality of ribbon Schur functions

In this section we collect the definitions and results from [1] that are relevant for our approach. As mentioned before, the basic objects we will be working with are compositions:

Definition 2.1. A composition α of a positive integer m, denoted $\alpha \vDash m$, is a list of positive integers $(a_1, a_2, \dots, a_k)$ such that $a_1 + a_2 + \dots + a_k = m$. We refer to each of the a_i as components, and say that α has length $l(\alpha) = k$ and size $|\alpha| = m$.

Definition 2.2. Let $\alpha = (a_1, a_2, \dots, a_k) \vDash m$ and $\beta = (b_1, b_2, \dots, b_\ell) \vDash n$. Then the *concatenation* of α and β is the composition

$$\alpha \cdot \beta = (a_1, \dots, a_k, b_1, \dots, b_\ell) \vDash n + m.$$

Their *near concatenation* is

$$\alpha \odot \beta = (a_1, \dots, a_k + b_1, \dots, b_\ell) \vDash n + m.$$

Writing

$$\alpha^{\odot n} = \underbrace{\alpha \odot \alpha \odot \dots \odot \alpha}_{n}$$

we define the *composition* of α and β as

$$\alpha \circ \beta = \beta^{\odot a_1} \cdot \beta^{\odot a_2} \cdots \beta^{\odot a_k} \vDash mn.$$

The composition $\alpha = (a_1, a_2, \dots, a_k)$ is *symmetric* if it coincides with its *reversal* $\alpha^* = (a_k, a_{k-1}, \dots, a_1)$.

The following theorem shows that composition of compositions is a very well behaved operation indeed:

Theorem 2.3 ([1], Propositions 3.3, 3.7, 3.8 and 3.9). *The set of compositions together with the operation $\circ$ is a monoid, i.e., $\circ$ is associative and has neutral element* (1). *Furthermore, $|\alpha \circ \beta| = |\alpha|\,|\beta|$ and $l(\alpha \circ \beta) = l(\alpha) + |\alpha|\,(l(\beta) - 1)$. Finally, $(\alpha \circ \beta)^* = \alpha^* \circ \beta^*$.*

Note that composition of compositions is not commutative. For example, $(1,1) \circ (2) = (2)^{\odot 1} \cdot (2)^{\odot 1} = (2,2)$, but $(2) \circ (1,1) = (1,1)^{\odot 2} = (1,1) \odot (1,1) = (1,2,1)$.

Definition 2.4 ([1], Definition 3.5). If a composition α is written in the form $\alpha_1 \circ \alpha_2 \circ \cdots \circ \alpha_k$ then we call this a *factorisation* of α. A factorisation $\alpha = \beta \circ \gamma$ is called trivial if any of the following conditions are satisfied:

(1) one of β and γ is the composition 1,
(2) the compositions β and γ both have length 1,
(3) the compositions β and γ both have all components equal to 1.

A factorisation $\alpha_1 \circ \alpha_2 \circ \cdots \circ \alpha_k$ is called *irreducible* if no $\alpha_i \circ \alpha_{i+1}$ is a trivial factorisation, and each α_i admits only trivial factorisations. We call a composition α *irreducible*, if it has not length 1, not all of its components are equal to 1 and it admits only trivial factorisations.

We remark that irreducible compositions are *not* defined in [1]. In particular, the notion does not coincide with the notion of irreducible factors.

Theorem 2.5 ([1], Theorem 3.6). *The irreducible factorisation of any composition is unique.*

It is not surprising that such a theorem is very useful to enumerate the underlying objects. For experimentation it was also of great help to have a relatively efficient test for irreducibility, which is exhibited in Definition 4.11 and Lemma 4.15 of [1].[1]

Finally, equivalence of compositions and therefore equality of ribbon Schur functions is reduced to factorisation by the following theorem. Note that it was well known before that reversal of compositions yields the same ribbon Schur functions, see for example Exercise 7.56 in Enumerative Combinatorics 2 [7], which includes also the natural extension to skew Schur functions.

Theorem 2.6 ([1], Theorem 4.1). *Two compositions β and γ satisfy $\beta \sim \gamma$ if and only if for some k, $\beta = \beta_1 \circ \beta_2 \circ \cdots \circ \beta_k$ and $\gamma = \gamma_1 \circ \gamma_2 \circ \cdots \circ \gamma_k$ where, for each i, either $\gamma_i = \beta_i$ or $\gamma_i = \beta_i^*$.*

3. The number of ribbon Schur functions of a given size

Definition 3.1. We order the set of compositions of a given length lexicographically. Thus, let $\alpha = (a_1, a_2, \ldots, a_k)$ and $\beta = (b_1, b_2, \ldots, b_k)$ be two compositions. Then $\alpha < \beta$ if and only if $a_s < b_s$ for some s, such that $a_r = b_r$ for all $r < s$. The composition α is *lexicographic minimal* if $\alpha \leq \alpha^*$.

In view of Theorem 2.5 and Theorem 2.6, we call a composition *normalised*, if all factors in its irreducible factorisation are lexicographic minimal.

[1] An implementation can be obtained from the author of the present article.

Thus, to determine the number of distinct ribbon Schur functions, it is sufficient to count normalised compositions. This is not hard to achieve using a suitable combinatorial decomposition. The validity of our decomposition hinges on the following lemma:

Lemma 3.2. *Consider a composition α with irreducible factorisation $\alpha_1 \circ \alpha_2 \circ \cdots \circ \alpha_k$. Then α is symmetric if and only if all α_i are symmetric for $i \in \{1, \dots, k\}$.*

If α is asymmetric, then there is an $\ell \in \{1, \dots, k\}$ such that α_ℓ is asymmetric, and α_i is symmetric for all $i > \ell$. In this situation, $\alpha < \alpha^*$ if and only if $\alpha_\ell < \alpha_\ell^*$.

Proof. By the last statement of Theorem 2.3, an irreducible factorisation of the reversal of α is $\alpha^* = \alpha_1^* \circ \alpha_2^* \circ \cdots \circ \alpha_k^*$. Thus, by Theorem 2.5, if $\alpha = \alpha^*$, all the factors α_i are symmetric. The reverse direction follows immediately from the last statement of Theorem 2.3.

Suppose now that α is asymmetric. We have to show that $\alpha < \alpha^*$ if and only if $\alpha_\ell < \alpha_\ell^*$, where ℓ is maximal such that α_ℓ is asymmetric. Let us first prove for compositions β, γ and δ:

$$\beta \circ \delta < \gamma \circ \delta \text{ if and only if } \beta < \gamma, \quad \text{whenever } l(\beta) = l(\gamma). \tag{1}$$

Indeed, if $\beta = (b_1, \dots, b_r) < \gamma = (g_1, \dots, g_r)$, then there is an index j such that $b_j < g_j$ and $b_i = g_i$ for all $i < j$. Since $\beta \circ \delta = \delta^{\odot b_1} \cdots \delta^{\odot b_r}$ and $\gamma \circ \delta = \delta^{\odot g_1} \cdots \delta^{\odot g_r}$, it suffices to compare $\delta^{\odot b_j}$ and $\delta^{\odot g_j}$.

If $\delta = (d_1)$ has length 1, we have $\delta^{\odot b_j} = (b_j d_1) < (g_j d_1) = \delta^{\odot g_j}$. Let us now consider the case that $\delta = (d_1, \dots, d_s)$ with $s \geq 2$. Then the component with index $1 + b_j\,(l(\delta) - 1)$ of $\delta^{\odot b_j}$, i.e., its last component, equals d_s. However, since $b_j < g_j$, the component of $\delta^{\odot g_j}$ with the same index is $d_s + d_1$, which is strictly greater than d_s. Hence $\beta \circ \delta < \gamma \circ \delta$. The converse follows by symmetry: given $\beta \circ \delta < \gamma \circ \delta$, assuming that $\beta > \gamma$ leads to a contradiction by what we have just shown.

Next, we prove for compositions β, γ, δ and ϵ:

$$\beta \circ \delta < \gamma \circ \epsilon \text{ if and only if } \delta < \epsilon, \qquad \text{whenever } l(\beta) = l(\gamma),\ l(\delta) = l(\epsilon),\ |\delta| = |\epsilon| \text{ and } \delta \neq \epsilon. \tag{2}$$

Suppose that $\delta = (d_1, d_2, \dots, d_r) < \epsilon = (e_1, e_2, \dots, e_r)$. It suffices to compare the first $r-1$ components of $\beta \circ \delta$ and $\gamma \circ \epsilon$, which are $d_1, d_2, \dots, d_{r-1}$ and $e_1, e_2, \dots, e_{r-1}$, respectively. Let j be minimal such that $d_j < e_j$. Since $|\delta| = |\epsilon|$, the two compositions cannot differ only in the last component, so $j \leq r-1$, which implies $\beta \circ \delta < \gamma \circ \epsilon$. Again, the converse follows by symmetry.

To conclude the proof, we write $\alpha = \beta \circ \alpha_\ell \circ \gamma$, where ℓ is maximal such that α_ℓ is asymmetric – if $\ell = 1$ then $\beta = (1)$, if $\ell = k$ then $\gamma = (1)$. Thus, by the last statement of Theorem 2.3, $\alpha^* = \beta^* \circ \alpha_\ell^* \circ \gamma$. By Condition (1) $\alpha < \alpha^*$ if and only if $\beta \circ \alpha_\ell < \beta^* \circ \alpha_\ell^*$. By Condition (2), this is the case if and only if $\alpha_\ell < \alpha_\ell^*$, as desired. ☐

In the following lemma we collect the facts we need about Dirichlet generating functions:

Lemma 3.3. *Let A and B be sets of compositions, let $A \uplus B$ be their disjoint union and define $A \circ B := \{\alpha \circ \beta : \alpha \in A, \beta \in B\}$. For any set of compositions A, let $A(s) = \sum_{\alpha \in A} |\alpha|^{-s}$ the associated Dirichlet generating function. Then*

$$(A \uplus B)(s) = A(s) + B(s) \quad \textit{and}$$
$$(A \circ B)(s) = A(s)B(s).$$

The latter equality is equivalent to the statement, that the coefficient of n^{-s} in $(A \circ B)(s)$ is $a_n * b_n$, where a_n and b_n are the coefficients of n^{-s} in $A(s)$ and $B(s)$, respectively, and $a_n * b_n$ denotes the Dirichlet convolution $\sum_{d|n} a_d b_{n/d}$.

Remark. A full-fledged combinatorial theory of Dirichlet series within the theory of combinatorial species was developed by Manuel Maia and Miguel Méndez [3]. Although the proofs below are written in the spirit of that theory, they are quite elementary.

Theorem 3.4. *The number of normalised compositions of size n is*

$$2 \cdot 2^{n-1} * 2^{\lfloor \frac{n}{2} \rfloor} * \left(2^{n-1} + 2^{\lfloor \frac{n}{2} \rfloor}\right)^{-1},$$

where $a_n * b_n$ denotes the Dirichlet convolution, and the reciprocal is the inverse with respect to Dirichlet convolution.

Remark. Thus, the numbers of ribbon Schur functions of size 1 to 33 turn out to be:

$$1, 2, 3, 6, 10, 20, 36, 72, 135, 272, 528, 1052, 2080, 4160, 8244, 16508, 32896, 65770,$$
$$131328, 262632, 524744, 1049600, 2098176, 4196200, 8390620, 16781312, 33558291,$$
$$67116944, 134225920, 268451240, 536887296, 1073774376, 2147515424.$$

This is sequence `http://oeis.org/A120421` in the on-line encyclopedia of integer sequences [6].

It may be interesting to compare the number of ribbon Schur functions with the number of lexicographic minimal compositions. Since $|\alpha \circ \beta| = |\alpha| \cdot |\beta|$, it is clear that the numbers coincide when n is prime. For $n = 9$, there are 136 lexicographic minimal compositions, but two of them are equivalent. Here are the differences and their positions up to $n = 33$:

n :	9	12	15	16	18	20	21	24	25	27	28	30	32	33
difference :	1	4	12	4	22	24	56	152	36	237	112	600	216	992

Proof. Let R be the set of normalised compositions. Let S be the set of symmetric compositions, $P^\times$ be the set of (normalised) asymmetric irreducible compositions and

$$R^1 = P^\times \circ S, \tag{3}$$

i.e., the set of (normalised) compositions whose first factor in the irreducible factorisation is asymmetric, and all remaining factors (if any) are symmetric. We can then decompose the set of normalised compositions recursively as

$$R = S \uplus (R \circ R^1), \tag{4}$$

since a normalised composition is either symmetric, or can be written in a unique way as a product of a normalised composition, an asymmetric irreducible factor and a symmetric composition.

The set R^1 can be described in terms of the set of all compositions C and the set of asymmetric lexicographic minimal compositions $L^\times$ by Lemma 3.2. Namely,

$$L^\times = C \circ R^1, \tag{5}$$

since an asymmetric composition is lexicographic minimal, if and only if the last asymmetric factor in its irreducible factorisation is lexicographic minimal.

Finally, we have (again by Lemma 3.2)

$$2L^\times = C \setminus S,$$

where $2L^\times$ is interpreted as the set of asymmetric compositions whose last asymmetric factor is either lexicographic minimal or lexicographic maximal.

We can now apply Lemma 3.3 to obtain the Dirichlet generating function for the set of normalised compositions. We have

$$\begin{aligned} L^\times(s) &= 1/2\,(C(s) - S(s)) \\ R^1(s) &= L^\times(s)/C(s) \end{aligned}$$

and therefore

$$\begin{aligned} R(s) &= \frac{S(s)}{1 - R^1(s)} \\ &= \frac{2C(s)S(s)}{2C(s) - (C(s) - S(s))} \\ &= \frac{2C(s)S(s)}{C(s) + S(s)}. \end{aligned}$$

Since $C(s) = \sum_{n\geq 1} 2^{n-1} n^{-s}$ and $S(s) = \sum_{n\geq 1} 2^{\lfloor \frac{n}{2} \rfloor} n^{-s}$, the claim follows. □

Remark. It is not difficult to obtain more information using the preceding theorem and the decompositions in its proof. In particular, we can easily refine the count of normalised compositions by taking into account the number of asymmetric irreducible factors. Denoting the number of asymmetric irreducible factors of a composition ρ by $\alpha(\rho)$ and defining $R(s, z) = \sum_{\rho\in R} |\rho|^{-s} z^{\alpha(s)}$, we find

$$R(s, z) = \frac{S(s)}{1 - zR^1(s)} = \frac{2C(s)S(s)}{2C(s) - z\,(C(s) - S(s))}.$$

Perhaps more interesting, we can determine the generating function for irreducible compositions by size using the following proposition:

Proposition 3.5. *Let $P(s)$ be the Dirichlet generating function for (normalised) irreducible compositions, $P^*(s)$ the Dirichlet generating function for symmetric irreducible compositions and $R(s)$ the Dirichlet generating function for normalised compositions by size.*

Furthermore, let $S(s) = \sum_{n\geq 1} 2^{\lfloor \frac{n}{2} \rfloor} n^{-s}$ *be the Dirichlet generating function of symmetric compositions, and* $\zeta(s) = \sum_{n\geq 1} n^{-s}$ *the Riemann zeta function. We then have*

$$P(s) = 2\zeta^{-1}(s) - 1 - R^{-1}(s) \tag{6}$$

$$P^*(s) = 2\zeta^{-1}(s) - 1 - S^{-1}(s) \tag{7}$$

and

$$P^\times(s) = S^{-1}(s) - R^{-1}(s). \tag{8}$$

Remark. Thus, the numbers of (normalised) irreducible compositions of size 1 to 33 are:

$$0, 0, 1, 2, 8, 10, 34, 56, 126, 234, 526, 972, 2078, 4018, 8186, 16240, 32894, 65164,$$
$$131326, 261544, 524530, 1047490, 2098174, 4191680, 8390520, 16772994, 33557508,$$
$$67100304, 134225918, 268416590, 536887294, 1073708400, 2147512258.$$

Note that, whenever n is prime, there are precisely two normalised compositions (or, equivalently, lexicographic minimal compositions) that are not irreducible, namely the composition with all components equal to 1 and the composition (n).

For $n = 4$, the irreducible normalised compositions are $(1, 3)$ and $(1, 1, 2)$. For $n = 6$, they are $(1, 5)$, $(1, 1, 4)$, $(1, 4, 1)$, $(1, 2, 3)$, $(2, 1, 3)$, $(1, 1, 1, 3)$, $(1, 1, 2, 2)$, $(1, 1, 3, 1)$, $(2, 1, 1, 2)$, $(1, 1, 1, 1, 2)$.

Proof. Let E be the set of compositions with at least two components, all equal to 1. Let K be the set of compositions different from 1 with only one component. Let R be the set of all normalised compositions. Let R_E be the set of normalised compositions with no factors in the irreducible factorisation having only one component, i.e., all factors being irreducible compositions or having at least two components, all equal to 1. Finally, let P be the set of (normalised) irreducible compositions.

By Theorem 2.5, R_E is the disjoint union of the sets P, $P \circ R_E$, E, $E \circ P$ and $E \circ P \circ R_E$. Passing to (Dirichlet) generating functions, we obtain

$$R_E(s) = E(s) + \big(1 + E(s)\big)P(s)\big(1 + R_E(s)\big). \tag{9}$$

Similarly, R is the disjoint union of the composition (1) and the sets R_E, K, $K \circ R_E$ and $R_E \circ K \circ R$. Informally, we decompose the set of normalised compositions into subsets depending on the position of the first factor with only one component in the irreducible factorisation, if such a factor occurs. Hence

$$R(s) = \big(1 + K(s)\big)\big(1 + R_E(s)\big) + R_E(s)K(s)R(s). \tag{10}$$

We can now extract $R_E(s)$ from Equation (10) and plug it into Equation (9) to obtain $P(s)$ in terms of $E(s)$, $K(s)$ and $R(s)$. Observing that $E(s) = K(s) = \zeta(s) - 1$ we obtain Equation (6). Equation (8) can be derived by combining Equations (3) and (4). Equation (7) then follows from Equations (6) and (8). $\square$

4. The number of ribbon Schur functions of a given size and length

Apart from the size of a composition, the most natural statistic that comes to mind is its length. In this section we derive an expression for the number of normalised compositions with a given size and a given length.

By Theorem 2.3, it is possible to determine the length of a composition of compositions, knowing the size and the length of the factors. However, since the length of a composition of compositions is neither multiplicative or additive, we cannot expect a result as appealing as in Theorem 3.4.

Let us first collect some elementary results:

Proposition 4.1. *Let* $C_n(x) = \sum_{\alpha\in C, |\alpha|=n} x^{l(\alpha)}$ *be the ordinary generating function of all compositions of size* n*, where* x *marks length. Similarly, let* $S_n(x) = \sum_{\alpha\in S, |\alpha|=n} x^{l(\alpha)}$ *the generating function of symmetric compositions, and* $L^\times_n(x) = \sum_{\alpha\in L^\times, |\alpha|=n} x^{l(\alpha)}$ *the generating function of asymmetric lexicographic minimal compositions. Then*

(http://oeis.org/A007318) $$C_n(x) = x(1+x)^{n-1},$$

(http://oeis.org/A051159) $$S_n(x) = \begin{cases} x(1+x)(1+x^2)^{(n-2)/2} & n \text{ even} \\ x(1+x^2)^{(n-1)/2} & n \text{ odd}, \end{cases}$$

(http://oeis.org/A034852) $$L^\times_n(x) = 1/2\left(C_n(x) - S_n(x)\right).$$

Theorem 4.2. *Let* $R_n(x) = \sum_{\rho\in R, |\rho|=n} x^{l(\rho)}$ *be the ordinary generating function of normalised compositions of size* n*, where* x *marks length. Similarly, let* $R^1_n(x) = \sum_{\rho\in R^1, |\rho|=n} x^{l(\rho)}$ *be the ordinary generating function of (normalised) compositions whose first factor in the irreducible factorisation is asymmetric, and all remaining factors (if any) are symmetric. Then we have*

$$R^1_n(x) = \sum_{\substack{k\geq 0 \\ 1=d_0|d_1|\dots|d_k|n \\ d_i\neq d_{i+1} \text{ for } i\in\{0,\dots,k-1\}}} (-1)^k L^\times_{n/d_k}(x^{d_k}) \prod_{i=0}^{k-1} C_{d_{i+1}/d_i}(x^{d_i})/x^{d_i} \tag{11}$$

and

$$R_n(x) = \sum_{\substack{k\geq 0 \\ d_1|d_2|\dots|d_{k+1}=n \\ d_i\neq d_{i+1} \text{ for } i\in\{1,\dots,k\}}} S_{d_1}(x) \prod_{i=1}^{k} R^1_{d_{i+1}/d_i}(x^{d_i})/x^{d_i}. \tag{12}$$

Proof. We reuse the decompositions from the proof of Theorem 3.4. From Equation (5), we obtain the equality of sets (subscripts denoting the size of the compositions we are restricting our attention to)

$$L^\times_n = \dot{\bigcup}_{d|n} C_d \circ R^1_{n/d}.$$

Since $l(\alpha \circ \beta) = l(\alpha) - |\alpha| + |\alpha|\, l(\beta)$, it follows that

$$L^{\times}_n(x) = \sum_{d|n} C_d(x) x^{-d} R^1_{n/d}(x^d). \tag{13}$$

Equation (11) then follows from Equation (15) in Lemma 4.3 below, with $A_n(x) = L^{\times}_n(x)$, $B_n(x) = C_n(x)/x^n$ and $D_n(x) = R^1_n(x)$.

Similarly, from Equation (4), we obtain the equality of sets

$$R_n = S_n \cup \dot{\bigcup}_{d|n, d\neq n} R_d \circ R^1_{n/d},$$

and therefore

$$R_n(x) = S_n(x) + \sum_{d|n, d\neq n} R_d(x) x^{-d} R^1_{n/d}(x^d). \tag{14}$$

Equation (12) then follows from Equation (16) in Lemma 4.3 below, with $A_n(x) = R_n(x)$, $B_n(x) = S_n(x)$ and $D_n(x) = R^1_n(x)/x$. □

Remark. Note that for actually computing the generating function for normalised compositions using a computer, Equations (13) and (14) may be easier to implement than the 'explicit' expressions given in the statement of the theorem.

Again, we can refine the count be marking the number of asymmetric irreducible factors with an additional variable z: every summand in Equation (12) has to be multiplied by z^k, since every composition in R^1_n contains exactly one asymmetric irreducible factor.

Lemma 4.3. *Suppose that $B_1(x) = 1$ and*

$$A_n(x) = \sum_{d|n} B_d(x) D_{n/d}(x^d).$$

Then we have

$$D_n(x) = \sum_{\substack{k\geq 0\\ 1=d_0|d_1|\dots|d_k|n\\ d_i \neq d_{i+1} \text{ for } i\in\{0,\dots,k-1\}}} (-1)^k A_{n/d_k}(x^{d_k}) \prod_{i=0}^{k-1} B_{d_{i+1}/d_i}(x^{d_i}). \tag{15}$$

Given

$$A_n(x) = B_n(x) + \sum_{d|n, d\neq n} A_d(x) D_{n/d}(x^d),$$

we have

$$A_n(x) = \sum_{\substack{k\geq 0\\ d_1|d_2|\dots|d_{k+1}=n\\ d_i \neq d_{i+1} \text{ for } i\in\{1,\dots,k\}}} B_{d_1}(x) \prod_{i=1}^{k} D_{d_{i+1}/d_i}(x^{d_i}). \tag{16}$$

Proof. We prove the statements by induction on n. For $n = 1$, the hypothesis is $A_1(x) = B_1(x)D_1(x) = D_1(x)$, and the right hand side of Equation (15) indeed evaluates to $A_1(x)$.

Now suppose that Equation (15) holds for $n < N$. Then

$$D_N(x) = A_N(x) - \sum_{1<d|N} B_d(x) D_{N/d}(x^d)$$
$$= A_N(x) - \sum_{1<d|N} B_d(x) \sum_{\substack{k\geq 0\\ 1=d_0|d_1|\dots|d_k|N/d\\ d_i\neq d_{i+1} \text{ for } i\in\{0,\dots,k-1\}}} (-1)^k A_{N/(d_k d)}(x^{d_k d}) \prod_{i=0}^{k-1} B_{d_{i+1}/d_i}(x^{d_i d}).$$

Substituting $d'_{i+1} = d_i d$ we obtain

$$D_N(x) = A_N(x) - \sum_{1<d|N} B_d(x) \sum_{\substack{k\geq 0\\ d=d'_1|d'_2|\dots|d'_{k+1}|N\\ d'_i\neq d'_{i+1} \text{ for } i\in\{1,\dots,k\}}} (-1)^k A_{N/(d'_{k+1})}(x^{d'_{k+1}}) \prod_{i=1}^{k} B_{d'_{i+1}/d'_i}(x^{d'_i})$$
$$= A_N(x) - \sum_{\substack{k\geq 0\\ 1=d'_0|d'_1|\dots|d'_{k+1}|N\\ d'_i\neq d'_{i+1} \text{ for } i\in\{0,\dots,k\}}} (-1)^k A_{N/(d'_{k+1})}(x^{d'_{k+1}}) \prod_{i=0}^{k} B_{d'_{i+1}/d'_i}(x^{d'_i}).$$

The final expression is equivalent to the claimed Equation (15), since $A_N(x)$ is precisely the summand corresponding to the chain $1 = d'_0|N$.

Equation (16) can be shown using the same strategy; the calculations are actually a bit easier. □

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