# Insertion algorithms for shifted domino tableaux

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#### Séminaire Lotharingien de Combinatoire

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Insertion algorithms

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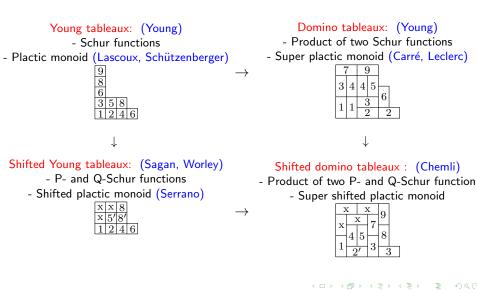








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A partition  $\lambda$  of *n* is a non-increasing sequence  $(\lambda_1, \lambda_2, \ldots, \lambda_k)$  such that  $\lambda_1 + \lambda_2 + \cdots + \lambda_k = n$ . We represent a partition by its Ferrers diagram.

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Figure: The Ferrers diagram of  $\lambda = (5,4,3,3,1)$ 

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Figure: A Young tableau of shape  $\lambda = (5,4,3,3,1)$ 

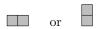
A Young tableau is a filling of a Ferrers diagram with positive integers such that rows are non-decreasing and columns are strictly increasing.

Two adjacent boxes form a domino:

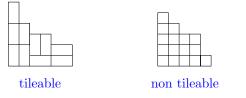


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A diagram is tileable if we can tile it by non intersecting dominos.



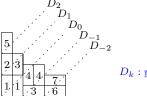
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$$\begin{array}{c} D_{2} \\ D_{1} \\ D_{0} \\ D_{-1} \\ D_{-2} \\ \hline \\ 2 \\ 3 \\ 1 \\ 1 \\ 1 \\ 3 \\ 6 \\ \end{array}$$

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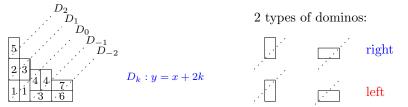


 $D_k: y = x + 2k$ 

2 types of dominos:

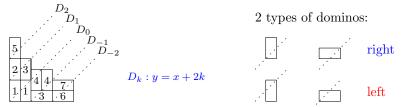


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A tilling is acceptable iff there is no vertical domino d on  $D_0$  such that the only domino adjacent to d on the left is strictly above  $D_0$ .





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Given an acceptable tilling, a shifted domino tableau is:

- a filling of dominos strictly above  $D_0$  by x
- a filling of other dominos with integers in  $\{1' < 1 < 2' < 2 < \cdots\}$
- columns and rows are non decreasing
- an integer without ' appears at most once in every column
- an integer with ' appears at most once in every row





### Theorem (Chemli, P. (2016))

There is a bijective algorithm f, with a bicolored word as input and a pair (P, Q) of shifted domino tableaux as output such that:

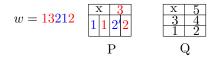
• P and Q have same shape

$$w = 13212 \qquad \boxed{\begin{array}{c} x & 3 \\ 1 & 2 & 2 \\ \end{array}} \qquad \boxed{\begin{array}{c} x & 5 \\ 3 & 4 \\ 1 & 2 \\ \end{array}} \qquad P \qquad Q$$

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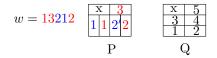
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There is a bijective algorithm f, with a bicolored word as input and a pair (P, Q) of shifted domino tableaux as output such that:

- P and Q have same shape
- P is without ' on  $D_0$
- Q is standard without '



Let  $w_1$  be a word in  $\mathbb{N} \times \{L\}$  with P-tableau of shape  $\mu$ , and  $w_2$  be a word in  $\mathbb{N} \times \{R\}$  with P-tableau of shape  $\nu$ . Let  $\lambda$  be the shape of the P-tableau of the word  $w_1w_2$ . We have:

$$\sum_{T, sh(T)=\lambda} x^T = P_{\mu} P_{\nu}$$

, where  $P_{\mu}$  is a P-Schur function.

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Two words belong to the same class of the super shifted plactic monoid iff they have the same P-tableau.

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## Theorem (Chemli, P. (2016))

There is an algorithm g with a bicolored standard word as input and a pair (P, Q) of shifted domino tableaux as output such that :

- P and Q have the same shape
- P is standard without '
- Q is standard without ' in D<sub>0</sub>

### Conjecture 1

If  $\sigma$  is a signed permutation (that we identify with a bicolored standart word) then

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#### Conjecture 2

Algorithm *f* commutes with standardization and truncation.

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- Enumerative consequences
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- Hook formula for shifted domino tableaux

Thank you for your attention!

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