Reconstruction of polytopes and Kalai's conjecture on reconstruction of spheres

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Séminaire Lotharingien de Combinatoire
Strobl
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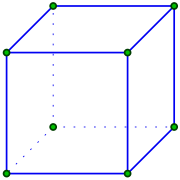
Outline

- Convex polytopes
- Kalai's conjecture on reconstruction of spheres
- Subword complexes
- Manifolds

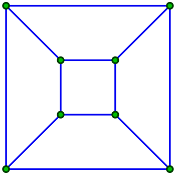
Convex polytopes

(Convex) polytope P:
convex hull of finitely many points in Euclidian space.

The graph $G(P)$:
the graph consisting of the vertices and edges of P.



P



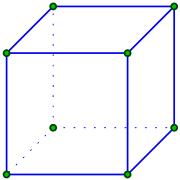
$G(P)$

Simple polytope P:
number of edges incident to each vertex equals the dimension of P.

Reconstruction of polytopes

Theorem (Blind–Mani, 1987)

If P is a simple polytope, then the graph $G(P)$ determines the entire combinatorial structure of P.



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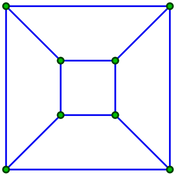

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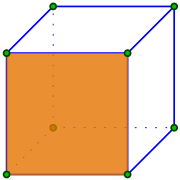

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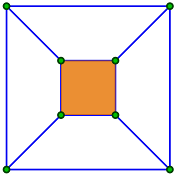

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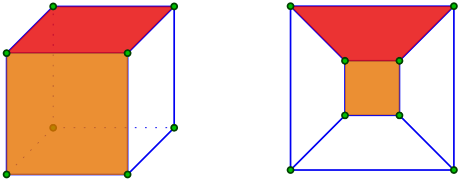

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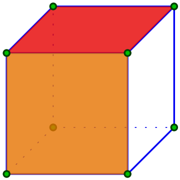

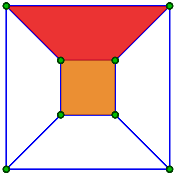

Kalai, 1988: A simple constructive proof.

Reconstruction of polytopes

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This holds for arbitrary polytopes (not only simple) in dimension 3 (Steinitz Theorem), but not in higher dimensions.

Example

Let Δ_m be a m-dimensional simplex. The following are two non isomorphic 6-dimensional polytopes with the same graph (complete graph on 8 vertices)

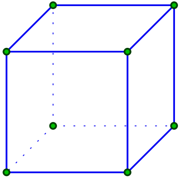
$$(\Delta_2 \times \Delta_4)^* \ncong (\Delta_3 \times \Delta_3)^*$$

Duality of polytopes

Every nonempty d-polytope P in $\mathbb{R}^d$ admits a dual polytope in $\mathbb{R}^d$:

$$P^* = \{y \in \mathbb{R}^d : x^T y \leq 1 \text{ for all } x \in P\}$$

where P is assumed to contain the origin in its interior.



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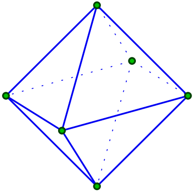

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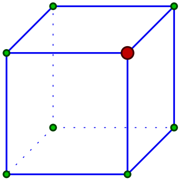

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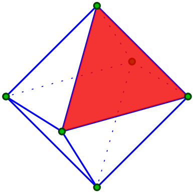

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$$\begin{array}{rcl} P & \longleftrightarrow & P^* \\ \text{vertices} & \longleftrightarrow & \text{facets (maximal dimensional faces)} \end{array}$$

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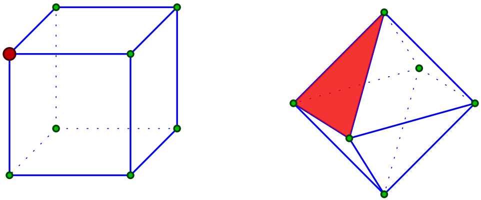

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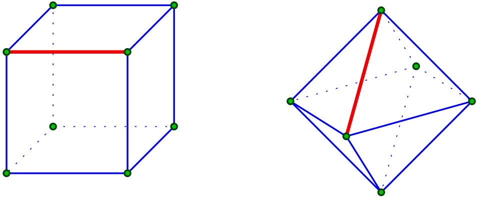

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Under this duality:

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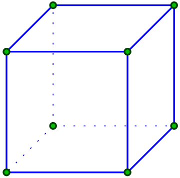

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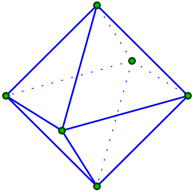

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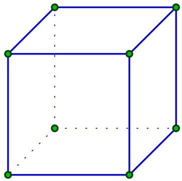
Simple vs simplicial

Simplicial polytope P:
all faces are simplices.

The facet-ridge graph $G_{FR}(P)$:
the graph whose vertices are facets of P
two facets are connected by an edge if they intersect in a ridge.

$$\begin{array}{ccc} P \text{ is simple} & \longleftrightarrow & P^* \text{ is simplicial} \\ G(P) & = & G_{FR}(P^*) \end{array}$$



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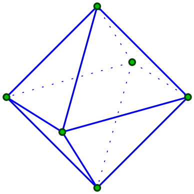

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Reconstruction of polytopes and spheres

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Simplicial polytopes are completely determined by their facet-ridge graphs.

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Simplicial spheres are completely determined by their facet-ridge graphs.

A simplicial sphere is a simplicial complex which is homeomorphic to a sphere.

Most spheres are not polytopal

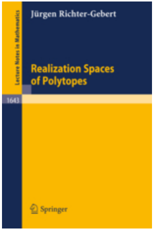
For $d \geq 3$, most d-spheres are not polytopal.

- Goodman–Pollack, 1986
- Kalai, 1988
- Pfeifle–Ziegler, 2004

Deciding polytopality of spheres is a difficult problem

Mnëv's Universality theorem: Realization spaces of polytopes can take arbitrary (semi-algebraic) shapes and thus can exhibit all kinds of pathologies.

The realizability problem for 4-polytopes is NP-hard.



Goal

Our initial goal was:
Look for a counterexample to Kalai's Conjecture among a special family of simplicial spheres which are conjectured to be polytopal.
(kill two conjectures at once)

Instead:
We proved the conjecture for this family.
(spherical subword complexes)

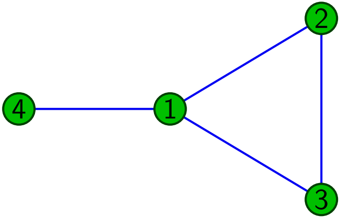
Rest of the talk:
Introduce subword complexes and state our main result.

Subword complexes preliminaries

Simplicial Complex Δ: A collection of subsets of a ground set E which is closed under containment:

$$\sigma \in \Delta \text{ and } \tau \subseteq \sigma \longrightarrow \tau \in \Delta$$

faces: subsets in Δ
vertices: singleton sets
facets: maximal sets
ridges: facets missing a single element



$$\Delta = \{\emptyset, \{1\}, \{2\}, \{3\}, \{4\}, \{1,2\}, \{1,3\}, \{1,4\}, \{2,3\}\}$$

Subword complexes preliminaries

Symmetric group $\mathbb{S}_{n+1}$:
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length of w: smallest r such that $w = s_{i_1} \ldots s_{i_r}$
longest element: permutation $[n+1,\ \ldots\ , 1]$
reduced expression for w: expression for w of minimal length

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In this talk: **finite Coxeter groups**
(very similar to the symmetric group)

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W finite Coxeter group with generating set S

$Q = (q_1, \ldots, q_m)$ a word in S

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Definition (Knutson–Miller, 2004)

The subword complex $\Delta(Q, \pi)$ is the simplicial complex whose

faces $\longleftrightarrow$ subwords P of Q such that $Q \setminus P$ contains a reduced expression of π

Knutson–Miller. Gröbner geometry of Schubert polynomials. Ann. Math., 161(3), '05
Knutson–Miller. Subword complexes in Coxeter groups. Adv. Math., 184(1), '04

Subword complexes - Example 1

In type A_2:

$W = \mathbb{S}_3$, $S = \{s_1, s_2\} = \{(1\ 2), (2\ 3)\}$

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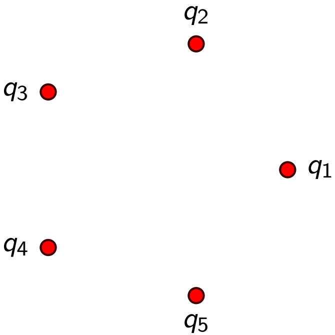

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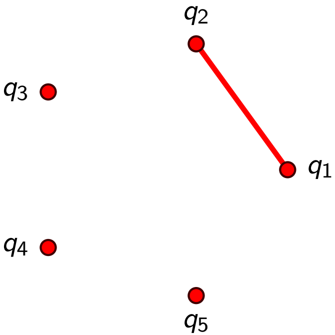

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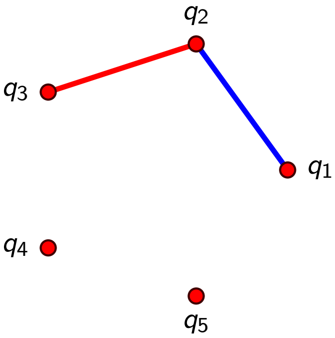

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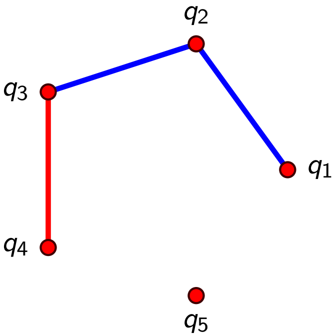

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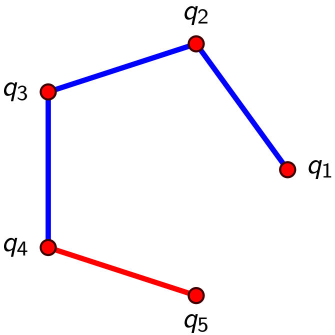

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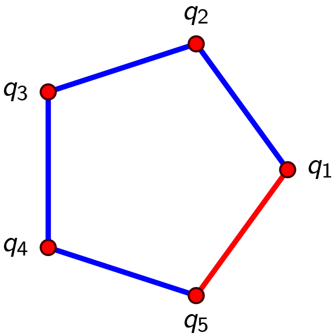

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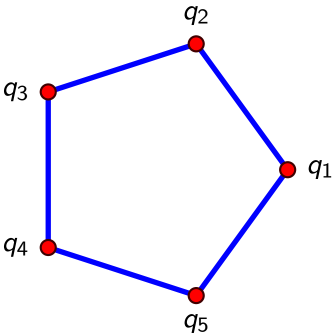

Subword complexes - Example 2

In type A_3:

$W = \mathbb{S}_4$, $S = \{s_1, s_2, s_3\} = \{(1\ 2), (2\ 3), (3\ 4)\}$

$$Q = \begin{matrix} (\ s_1, s_2, s_1, s_2, s_1, {\color{red}s_3}\ ) \\ q_1, q_2, q_3, q_4, q_5, {\color{red}q_6} \end{matrix} \quad \text{and } \pi = [3\ 2\ 1] = s_1 s_2 s_1 = s_2 s_1 s_2$$

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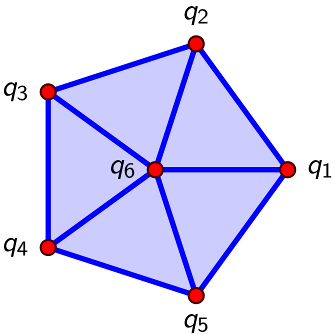

Subword complexes

Theorem (Knutson-Miller, 2004)

Subword complexes are vertex decomposable spheres or balls.

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Spherical subword complexes are polytopal.

Special cases include:

- Cyclic polytopes
- Duals of associahedra
- Cluster complexes of cluster algebras of finite type
- Duals of pointed-pseudotriangulation polytopes
- Simplicial multi-associahedra (conjectured)

Woo, Pilaud–Pocchiola, Serrano–Stump, Stump, C.-Labbé–Stump, Rote–Santos–Streinu, Jonsson, ...

Our main theorem

Theorem (C.–Doolittle)

Spherical subword complexes of finite type are completely determined by their facet-ridge graph. In other words, they satisfy Kalai's Conjecture.

Our current proof is <u>not</u> constructive.
It is based on the topological tools developed by Blind and Mani.

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Open Problem

Find a combinatorial constructive proof.

Manifolds

Question

Are simplicial manifolds completely determined by their facet-ridge graphs?

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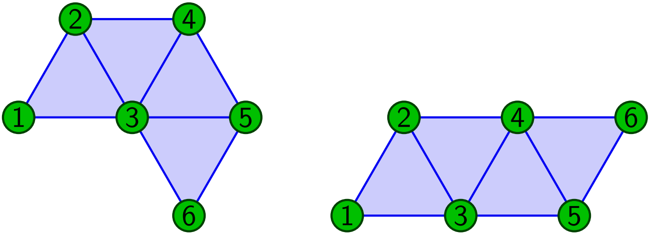
Answer: No



Manifolds

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Are simplicial manifolds without boundary completely determined by their facet-ridge graphs?

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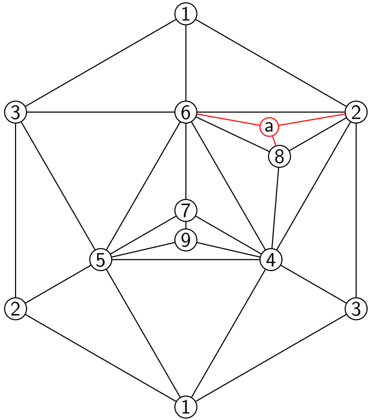
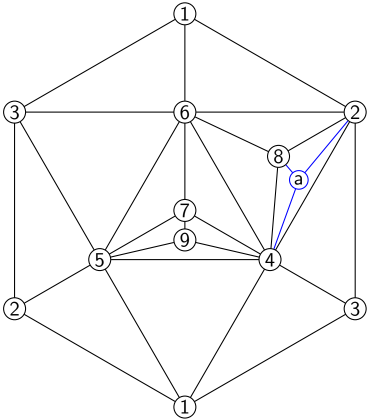
Are simplicial manifolds without boundary completely determined by their facet-ridge graphs?

Answer: No. Counterexamples can be explicitly found using ideas of Blind–Mani and Adiprasito.

Manifolds are not reconstructible

Example 1: Two non-isomorphic triangulations of the projective plane with isomorphic facet-ridge graphs:





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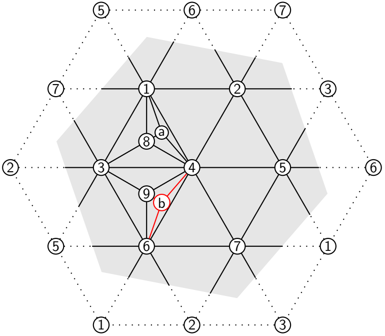
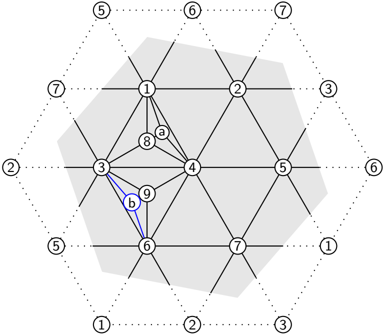
Example 2: Two non-isomorphic triangulations of the torus with isomorphic facet-ridge graphs:





Thank you!