

## HAMMING SPACES AND MAXIMAL SELF DUAL CODES OVER $GF(q)$ , $q = \text{ODD}$

CHAT YIN HO<sup>(\*)</sup>

### Abstract

The group of automorphisms of a Hamming space is determined. Self dual codes over odd characteristic finite field with respect to bilinear forms are treated. Under the subgroup of the monomial group preserving the inner product, we classify the maximal self dual codes over  $GF(5)$  with respect to the inner product of dimension  $\leq 8$ . The Hamming weight distribution and the order of the automorphism of the code are given.

### 0. Introduction

Although coding theory started from an engineering problem in the late 1940, the subject has developed by using more and more mathematical techniques. Besides the application of the error correction, one of its application is using the self dual codes over  $GF(2)$  to study the projective plane of order 10. Generalization of this to self dual codes over odd characteristic finite fields turns out to be very successful in obtaining a bound of the order of a group of automorphisms of a finite geometry, and in proving non-existence of various kinds of results for planes (cf. see [3]). This leads to classification of self dual codes of small dimension. In [5], self dual codes over  $GF(3)$  of dimension less than 12 have been classified.

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In this note, we first determine the group of automorphisms of a Hamming space in 1.1. Self dual codes over odd characteristic finite fields with respect to bilinear forms are treated in general in section 2. The subgroup  $H$  of the monomial group preserving the inner product is shown to be isomorphic to  $\mathbb{Z}_2 \wr S_n$ . As an illustration, in section 3, we classify the maximal self dual codes over  $GF(5)$  with respect to the standard inner product of dimension  $n \leq 8$  under  $H$ . The Hamming weight distribution and the order of the automorphism group of the code are given. Even though  $H$  preserves the Hamming weight, it does not preserve the complete weight which is invariant under the subgroup  $S$  of the symmetric group on  $n$  letters. It is shown that for  $n \leq 4$ , a  $H$ -orbit remains to be a  $S$ -orbit in the maximal self dual codes, and for  $n = 5$  one of the 2  $H$ -orbits becomes the union of 3  $S$ -orbits with different complete weight distributions.

Most of our results are self contained and notations are standard taken from Dembowski [1], Gorenstein [2], Huppert [4], or MacWilliams and Sloane [7].

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## 1. Preliminaries

A discrete metric space is a set  $X$  together with a function  $d$  from  $X \times X$  into the nonnegative integers such that for  $x, y, z \in X$  we have a)  $d(x, y) = 0$  if and only if  $x = y$ ; b)  $d(x, y) = d(y, x)$ ; c)  $d(x, z) \leq d(x, y) + d(y, z)$ . A code is a subset of a discrete metric space.

Let  $(X_1, d_1)$ ,  $(X_2, d_2)$  be two discrete metric spaces. Their direct product is the set  $X_1 \times X_2$  together with the functions  $d = d_1 \times d_2$  defined by

$$d((x_1, x_2), (y_1, y_2)) = d_1(x_1, y_1) + d_2(x_2, y_2).$$

An isomorphism between two discrete metric spaces  $(X_1, d_1)$  and  $(X_2, d_2)$  is a one-to-one and onto map  $\alpha$  between the sets such that  $d_2(\alpha(x), \alpha(y)) = d_1(x, y)$  for all  $x, y \in X_1$ . An automorphism is an isomorphism such that the two sets are equal. The set of all automorphisms of the discrete metric space  $(X, d)$  is denoted by  $\text{Aut}(X, d)$ . It is a group under composition.

Let  $\alpha \in \text{Aut}(X_1, d_1)$  and  $\beta \in \text{Aut}(X_2, d_2)$ . Define  $\alpha \times \beta \in \text{Aut}(X_1 \times X_2, d_1 \times d_2)$  by  $(\alpha \times \beta)(x_1, x_2) = (\alpha(x_1), \beta(x_2))$  for  $x_1 \in X_1$ , and  $x_2 \in X_2$ . This shows that

$$\text{Aut}(X_1, d_1) \times \text{Aut}(X_2, d_2) \subseteq \text{Aut}(X_1 \times X_2, d_1 \times d_2).$$

Let  $n$  be a positive integer and let  $S_n$  be the symmetric group on  $n$  letter. Then  $\text{Aut}(\prod_{i=1}^n X_i, \prod_{i=1}^n d_i)$ , where  $X_i = X$  and  $d_i = d$  for  $i=1, \dots, n$ , contains a subgroup isomorphic to  $\text{Aut}(X, d) \wr S_n$ , the wreath product of  $\text{Aut}(X, d)$  by  $S_n$ .

A discrete metric space  $(X, d)$  is called an equal distance space if  $d(x, y) = 1$  whenever  $x \neq y$  in  $X$ . For simplicity we will just write that  $X$  is an equal distance space and  $\text{Aut}(X)$  for  $\text{Aut}(X, d)$  when no confusion may be caused. Thus  $\text{Aut}(X) = S(X)$ , the symmetric group on  $X$ .

Let  $\Omega$  and  $F$  be two non-empty finite sets and  $V$  the set of all functions from  $\Omega$  into  $F$ . For  $f, g \in V$  define

$$d(f, g) = |\{x \in \Omega \mid f(x) \neq g(x)\}|.$$

Then  $(V, d)$  is called a Hamming space. The functional notation or the  $n$ -tuple notation for  $V$  will be used according to its convenience.

Thus  $V$  is the direct product of  $|\Omega|$  copies of the equal distance space  $F$ . Let  $|\Omega| = n$ . Then  $\text{Aut}(V, d)$  contains the group  $T = \prod_{x \in \Omega} T_x$ , where  $T_x \cong S(F)$  for  $x \in \Omega$ . Define  $S = \{S \in \text{Aut}(V, d) \mid \text{for } f \in V \text{ and } x \in \Omega, \text{ we have } (sf)(x) = f(\sigma_S(x)) \text{ where } \sigma_S \in S(\Omega)\}$ . Clearly  $S$  can be identified as  $S(\Omega)$  and is isomorphic to  $S_n$ . Also  $T$  is normalized by  $S$ .



**1.1 Theorem.** Using the above notation, we have  
 $\text{Aut}(V, S) = TS \cong S(F) \sim S_n$ .

**Proof.** Clearly  $TS \subseteq \text{Aut}(V, d)$ . Let  $\alpha \in \text{Aut}(V, d)$ . We apply induction on  $|\Omega|$  to show  $\alpha \in TS$ .

Fix  $P \in \Omega$ . Let  $V_2 = \{\text{all functions from } \Omega - \{P\} \text{ to } F\}$  and  $V_1 = \{\text{all functions from } P \text{ to } F\}$ . Let  $d_1$  and  $d_2$  be the Hamming distances of  $V_1$  and  $V_2$  respectively. Then  $V = V_1 \times V_2$ . Since  $\{\alpha f(P) \mid f \in V\} = F$ , there exists  $t_P \in T_P$  such that  $t_P(\alpha f(P)) = f(P)$  for all  $f \in V$ . Let  $t = t_P \times 1$ . Then  $t \in T$  and  $(t\alpha f)(P) = f(P)$  for all  $f \in V$ . For  $f \in V$ , define  $f_1 \in V_1$  and  $f_2 \in V_2$  by  $f_1(P) = f(P)$  and  $f_2(x) = f(x)$  for  $x \neq P$ . These give onto mappings from  $V$  to  $V_1$  and  $V$  to  $V_2$ , and  $d_2(g_2, h_2) = d(g, h) - d_1(g_1, h_1)$  for  $f, h \in V$ . Since  $(t\alpha f)_1 = f_1$  for all  $f \in V$ , we have  $d_2((t\alpha g)_2, (t\alpha h)_2) = d(t\alpha g, t\alpha h) - d_1((t\alpha g)_1, (t\alpha h)_1) = d(g, h) - d_1(g_1, h_1) = d_2(g_2, h_2)$ . Therefore the function  $\beta$  from  $V_2$  to  $V_2$  defined by  $\beta f_2 = (t\alpha f)_2$  for  $f \in V$  is well defined and  $\beta \in \text{Aut}(V_2, d_2)$ . By induction we have  $\beta \in \left( \prod_{\substack{x \in \Omega \\ x \neq P}} T_x \right) S(\Omega - \{P\}) := G$ . Since  $1 \times G \subseteq TS$  we have  $1 \times \beta \in TS$ . Since  $(t\alpha f)_1 = f_1$  for all  $f \in V$ ,  $1 \times \beta = t\alpha$ . Therefore  $t\alpha \in TS$  as desired.

Suppose  $F$  is a field. Then we can identify  $V$  with  $F^n$ . A linear code of  $V$  is a subspace of  $V$ . An affine subspace of  $V$  is a set  $v+U$ , where  $U$  is a subspace of  $V$ .

**1.2. Corollary.** Suppose  $|F| = 2$ . Two linear codes of  $V$  are in the same orbit of  $\text{Aut}(V, d)$  on the affine subspaces of  $V$  if and only if they are in the same orbit of  $S$  on the subspaces of  $V$ .

**Proof.** Since  $|F| = 2$ ,  $T$  is the group of all translations of  $V$ . By 1.1,  $\text{Aut}(V, d) = TS$ , which is a subgroup of the affine transformation group of  $V$ . This implies 1.2.

For  $v \in V$ , define the Hamming weight of  $v$  to be  $w(v) = d(v, 0)$ . For  $C \subseteq V$  and  $0 \leq i \leq n$ , let  $w_i(C) = |\{v \in C \mid w(v) = i\}|$ . The Hamming weight distribution of  $C$  is  $HW(C) = (w_0(C), w_1(C), w_2(C), \dots, w_n(C))$ .

## 2. Self dual codes over $\text{GF}(q)$ , $q$ odd

In this section let  $F = \text{GF}(q)$ , the finite field of  $q$  elements, where  $q$  is odd, and let  $\{\epsilon_1, \dots, \epsilon_n\}$  be the standard basis for  $V = F^n$ , and  $F^x = F \setminus \{0\}$ .

The monomial group of  $V$  is the set of all non-singular matrices such that each row and each column has exactly one non-zero element. Let  $S$  be the subgroup of the monomial group whose non-zero entries are 1. Then  $S$  is isomorphic to the symmetric group on  $n$  letters, and the monomial group is the product of  $S$  with the subgroup consisting of all non-singular diagonal matrices. The monomial group of  $V$  preserves the Hamming distance of  $V$ .

Let  $(,)$  be a non-degenerate symmetric bilinear form of  $V$ . By [Huppert, p. 238 Satz 10.9] we see that  $V = H_1 \perp \dots \perp H_m \perp V'$ , where  $H_i$  is a hyperbolic plane for  $i = 1, \dots, m$  and  $\dim V' = \delta \leq 2$ . If  $V' = \langle v \rangle$ , then  $(v, v) \neq 0$ . If  $\delta = 2$ , then  $V'$  has an orthogonal basis  $\{v_1, v_2\}$  such that  $(v_1, v_1) = -k(v_2, v_2)$ , where  $k$  is not a square in  $F^x$ .

The maximal totally isotropic subspaces of  $V$  have the same dimension  $m$  which is called the index of  $V$ . From a result of Serge [Dembowski, p. 46 Article 45] we have the following.

**2.1 Lemma.** The number of maximal totally isotropic subspaces of  $V$  is  $\prod_{i=1}^m (q^{i+\delta-1} + 1)$ .

For  $x = (x_1, \dots, x_n)$  and  $y = (y_1, \dots, y_n) \in V$ , let  $B_1(x, y) = x_1 y_1 + \dots + x_n y_n$  and  $B_2(x, y) = x_1 y_1 + \dots + x_n y_n - x_n y_n$ .



Let the discriminant of  $B_i(.,.)$ ,  $i=1,2$  be denoted by  $\Delta_i$ ,  $i=1,2$ . Then  $\Delta_1 \equiv 1$  and  $\Delta_2 \equiv -1 \pmod{(F^x)^2}$ , where  $F^x = F \setminus \{0\}$ .

**2.2. Lemma.** Concerning  $\delta$  we have the following:

- a)  $\delta = 1$  iff  $n$  is odd.
- b) If  $q \equiv 1 \pmod{4}$ , then  $\delta \leq 1$  for  $(.,.) = B_1(.,.)$  or  $B_2(.,.)$ .
- c) Suppose  $n \equiv 0 \pmod{4}$ . If  $(.,.) = B_1(.,.)$ , then  $\delta = 0$ . If  $(.,.) = B_2(.,.)$  and  $q \equiv 3 \pmod{4}$ , then  $\delta = 2$ .
- d) Suppose  $n \equiv 2 \pmod{4}$ . If  $(.,.) = B_2(.,.)$ , then  $\delta = 0$ . If  $(.,.) = B_1(.,.)$  and  $q \equiv 3 \pmod{4}$ , then  $\delta = 2$ .

**Proof.** a) is obvious.

Let the discriminate of  $(.,.)$  be denoted by  $\Delta$ . If  $\delta = 2$ , the decomposition of  $V$  into directed sum of hyperbolic planes and  $V'$  indicated above implies that  $\Delta \equiv (-1)^{m+1}k \pmod{(F^x)^2}$ .

b) Since  $q \equiv 1 \pmod{4}$ ,  $-1 \equiv 1 \pmod{(F^x)^2}$ . If  $\delta = 2$ , then  $k \equiv (-1)^{m+1}k \equiv \Delta \equiv 1 \equiv -1 \pmod{(F^x)^2}$ . However this implies that  $k$  is a square, a contradiction. Therefore  $\delta \leq 1$  as asserted.

In proving c) and d) we have  $n = 2m + \delta$ ,  $\delta$  is even.

c) Suppose  $n \equiv 0 \pmod{4}$ . First suppose  $\delta = 0$ . Then  $m = \frac{n}{2}$  is even. Hence  $\Delta \equiv (-1)^m \equiv 1 \pmod{(F^x)^2}$ . If  $(.,.) = B_2(.,.)$ , then  $\Delta = \Delta_2 \equiv -1 \pmod{(F^x)^2}$ . This implies that  $-1$  is a square in  $(F^x)$ , which forces  $q \equiv 1 \pmod{4}$ .

Now suppose  $\delta = 2$ , then  $m = \frac{n}{2} + 1$  is odd. Hence  $\Delta \equiv (-1)^{m+1}k \equiv k \pmod{(F^x)^2}$ . If  $(.,.) = B_1(.,.)$ , then  $1 \equiv \Delta_1 \equiv \Delta \equiv k \pmod{(F^x)^2}$ . This implies that  $k$  is square, a contradiction. Therefore  $\delta = 0$  when  $(.,.) = B_1(.,.)$ .

Assume  $(.,.) = B_2(.,.)$ . Then  $-1 \equiv \Delta_2 \equiv \Delta \equiv (-1)^{m+1}k \equiv k \pmod{(F^x)^2}$ . Since  $k$  is not a square,  $-1$  is not a square. Therefore  $q \equiv 3 \pmod{4}$  in this case.

d) Similar argument as in c).

In the rest of this section we assume that  $(.,.) = B_1(.,.)$ . Suppose  $\alpha = \text{diag}[1, \dots, \alpha, \dots, 1]$ , where  $\alpha \neq 0$  appears at the  $i$ -th row. If  $\alpha$  preserves  $(.,.)$ , then from  $(\epsilon_i, \epsilon_i) = (\alpha\epsilon_i, \alpha\epsilon_i) = \alpha^2(\epsilon_i, \epsilon_i)$  we have  $\alpha^2 = 1$ . This shows that the subgroup  $H$  of the monomial group preserving  $(.,.)$  is isomorphic to  $\{\pm 1\} \sim S_n$  and is a subgroup of  $\{F^x\} \sim S_{q-1} \sim S$ . For  $W \subseteq V$ , let  $W^\perp = \{v \in V \mid (w, v) = 0 \text{ for all } w \in W\}$ . For any linear code  $C$ , let  $\text{Aut}(C)$  be the subgroup of  $H$  leaving  $C$  invariant. If  $C \subseteq C^\perp$ , then  $C$  is a self dual code. Thus  $\{\text{maximal self dual codes of } V\} = \{\text{maximal totally isotropic subspaces of } V\}$ , and we denote this set by  $M_n$ . The  $H$ -orbits of  $M$  gives a classification of the members of  $M_n$ .

**2.3. Lemma.** Let  $q \equiv 1 \pmod{4}$ , and  $n = 2m$ . Let  $\Gamma = \{W \subseteq V \mid \dim W = m-1\}$ . We have the following:

$$a) |\Gamma| = \frac{q^m - 1}{q - 1} \prod_{i=2}^m (q^{i-1} + 1).$$

b) For  $Y \in \Gamma$ , there are exactly two maximal isotropic subspaces containing  $Y$  and they lie in the same orbit of  $H$ .

**Proof.** a) This is because

$$\begin{aligned} |\Gamma| &= \prod_{i=2}^m (q^{i-1} - 1)(q^{i-1} + 1) q^{\frac{(m-1)(m-2)}{2}} / |GL(m-1, q)| \\ &= \frac{q^m - 1}{q - 1} \prod_{i=2}^m (q^{i-1} + 1). \end{aligned}$$

b) For  $X = X^\perp$  let  $m(X) = |\{W \in \Gamma \mid W \subseteq X\}|$ . For  $W \in \Gamma$  let  $n(W) = |\{X \mid X = X^\perp \text{ and } W \subseteq X\}|$ . By Witt's theorem  $m(X)$ ,  $n(X)$  are constants. We now count  $\Lambda = \{(X, W) \mid X = X^\perp, W \subseteq X \text{ and } W \in \Gamma\}$  in two ways. Since  $q \equiv 1 \pmod{4}$  and  $n = 2m$ , 2.2 implies  $\delta = 0$ . By 2.1, there are  $\prod_{i=1}^m (q^{i-1} + 1)$  possibilities for  $X$ . Since

$$m(X) = \frac{q^n - 1}{q - 1},$$







isotropic subspace of dimension  $\frac{1}{2}(\dim U - 1) = \frac{1}{2}(n-2) = \text{index of } V$ . Hence  $U$  contains a  $x$ -subspace of  $V$ . So # equals the number of maximal totally isotropic subspaces of  $U$ , which is  $\frac{n-2}{2} \prod_{i=1}^{\frac{n-2}{2}} (q^{i+1})$  by 2.1.

d) By the proof of b) we may assume that  $n$  is even. Hence  $\dim U = n-1$  is odd, and the index of  $U$  is  $\frac{1}{2}(n-2)$ . If  $n \equiv 0 \pmod{4}$  2.2. c) implies that the index of  $V$  is  $\frac{1}{2}n$ . Therefore # = 0 in this case. If  $n \equiv 2 \pmod{4}$  but  $q \equiv 1 \pmod{4}$ , then the index of  $V$  is  $\frac{1}{2}n$  by 2.2 b) and again # = 0 as asserted.

We say that a linear code  $C$  is decomposable if

$$\{\epsilon_1, \dots, \epsilon_n\} = \bigcup_{j=1}^k B_j \text{ is a partition with } k \geq 2 \text{ and}$$

$$C = \bigoplus_{j=1}^k (C \cap \langle B_j \rangle). \text{ Otherwise } C \text{ is called indecomposable.}$$

### 3. Maximal self dual codes over GF(5)

In this section, under the group  $H$  we classify the maximal self dual codes over  $F = GF(5)$  with respect to  $(, ) = B(, )$  for  $\dim V = n \leq 8$ .

Let  $X \in M_n$ . The orbit containing  $X$  is  $X^H$  and  $|X^H| = [H : \text{Aut}(X)]$ . For any subset  $L$  of  $\text{Aut}(X)$  and  $v \in X$ , set  $C_L(v) = \{\ell \in L \mid v^\ell = v\}$  and  $L_{\langle v \rangle} = \{\ell \in L \mid \langle v \rangle^\ell = \langle v \rangle\}$ .

For  $a, b, c \in F^x$ , we have  $a^2 + b^2 + c^2 \neq 0$ . as a consequence we have the following.

**3.1. Lemma.** If  $X \in M_n$ , then  $w_3(X) = 0$ .

**3.2. Theorem.** The following holds, where \* means that the code is indecomposable. Also the code is generated by the row vector of the matrix when this matrix is given in the description of the code.

$n$	$ M_n $	$ H $	R(Representative from each orbit of $M_n$ under $H$ )	$ \text{Aut}(R) $	$ R^H $	$HW(R)$
1	0	1				
2	2	8	$X_2 = \langle (1, 2) \rangle$	4	2	(1, 0, 4)
3	6	$2^4 \cdot 3$	$X_3 = X_2 \oplus \langle 0 \rangle = \{(x, 0) \mid x \in X_2\}$	8	6	(1, 0, 4, 0)
4	12	$2^7 \cdot 3$	$X_4 = X_2 \oplus X_2$	$2^5$	6	(1, 0, 8, 0, 16)
5	$2^2 \cdot 3 \cdot 13$	$2^8 \cdot 3 \cdot 5$	$X_5 = X_2 \oplus X_2 \oplus \langle 0 \rangle$	$2^6$	$2^2 \cdot 3 \cdot 5$	(1, 0, 8, 0, 16, 0)
			$U_5 = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 0 & 1 & 2 & 3 & 4 \end{pmatrix}^*$	$2^3 \cdot 5$	$2^5 \cdot 3$	(1, 0, 0, 0, 20, 4)
6	$2^3 \cdot 3 \cdot 13$	$2^6 \cdot 6!$	$X_6 = X_2 \oplus X_2 \oplus X_2$	$2^7 \cdot 3$	$2^3 \cdot 3 \cdot 5$	(1, 12, 0, 48, 0, 64)
			$U_6 = \begin{pmatrix} 1 & 0 & 0 & 1 & 2 & 2 \\ 0 & 1 & 0 & 2 & 1 & 3 \\ 0 & 0 & 1 & 2 & 3 & 1 \end{pmatrix}^*$	$2^4 \cdot 3 \cdot 5$	$2^6 \cdot 3$	(1, 0, 0, 60, 24, 40)
7	$2^3 \cdot 3^3 \cdot 7 \cdot 13$	$2^7 \cdot 7!$	$X_7 = X_6 \oplus \langle 0 \rangle$	$2^8 \cdot 3$	$2^3 \cdot 3 \cdot 5 \cdot 7$	(1, 12, 0, 48, 0, 64, 0)
			$C_2 = U_6 \oplus \langle 0 \rangle$	$2^5 \cdot 3 \cdot 5$	$2^6 \cdot 3 \cdot 7$	(1, 0, 0, 62, 24, 40, 0)
			$C_3 = U_5 \oplus X_2$	$2^5 \cdot 5$	$2^6 \cdot 3^2 \cdot 7$	(1, 4, 0, 20, 4, 80, 16)
			$U_7 = \begin{pmatrix} 1 & 0 & 0 & 1 & 0 & 2 & 2 \\ 0 & 1 & 0 & 1 & 2 & 0 & 2 \\ 0 & 0 & 1 & 1 & 2 & 2 & 0 \end{pmatrix}^*$	$2^4 \cdot 3$	$2^7 \cdot 3 \cdot 5 \cdot 7$	(1, 0, 0, 24, 12, 72, 16)
8	$2^4 \cdot 3^3 \cdot 7 \cdot 13$	$2^8 \cdot 8!$	$X_8 = X_2 \oplus X_2 \oplus X_2 \oplus X_2$	$2^{11} \cdot 3$	$2^4 \cdot 3 \cdot 5 \cdot 7$	(1, 16, 0, 96, 0, 256, 0, 256)
			$T_2 = U_6 \oplus X_2$	$2^6 \cdot 3 \cdot 5$	$2^9 \cdot 3 \cdot 7$	(1, 4, 0, 60, 24, 280, 96, 160)
			$U_8 = \begin{pmatrix} 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 0 & 2 & 2 \\ 0 & 0 & 1 & 0 & 1 & 2 & 0 & 2 \\ 0 & 0 & 0 & 1 & 1 & 2 & 2 & 0 \end{pmatrix}^*$	$2^7 \cdot 3$	$2^8 \cdot 3 \cdot 5 \cdot 7$	(1, 0, 0, 48, 32, 288, 128, 128)



**Proof.** If  $n = 1$ , then  $|M_1| = 0$ . For  $n > 1$ , by 2.1 and

2.2 we see that  $|M_n| = \prod_{i=1}^{\frac{n}{2}} (5^{i-1} + 1)$  when  $n$  is even, and

$$|M_n| = \prod_{i=1}^{\frac{n-1}{2}} (5^{i-1} + 1) \text{ when } n \text{ is odd.}$$

Suppose  $n = 2$ . Then  $|M_2| = 2$  and  $|H| = 2^2 \cdot 2 = 8$ . Let  $X_2 = \langle (1, 2) \rangle$ . Then  $\text{Aut}(X_2) = \langle \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix} \rangle$  has order 4. Hence  $M_2 = X_2^H$ .

Suppose  $n = 3$ . Then  $|M_3| = 6$  and  $|H| = |Z_2 \vee S_3| = 2^4 \cdot 3$ . Let  $X_3 = X_2 \oplus \langle 0 \rangle = \{(x, 0) | x \in X_2\}$ . Then  $|\text{Aut}(X_3)| = 8$  and  $M_3 = X_3^H$ .

Suppose  $n = 4$ . Then  $|M_4| = 12$  and  $|H| = 2^4 \cdot 4! = 2^7 \cdot 3$ . Let  $X_4 = X_2 \oplus X_2$ . Thus  $\text{Aut}(X_4) \cong \text{Aut}(X_2) \vee S_2$  has order  $2^5$  and  $M_4 = X_4^H$ .

Suppose  $n = 5$ . Then  $|M_5| = 2^2 \cdot 3 \cdot 13$  and  $|H| = 2^8 \cdot 3 \cdot 5$ . Let  $X_5 = X_2 \oplus X_2 \oplus \langle 0 \rangle$ . Thus  $|\text{Aut}(X_5)| = |(\text{Aut}(X_2) \vee S_2) \times Z_2| = 2^6$ . So  $|X_5^H| = 2^2 \cdot 3 \cdot 5$ . Let  $U_5 = \langle u_1, u_2 | u_1 = (1, 1, 1, 1, 1) \text{ and } u_2 = (0, 1, 2, 3, 4) \rangle$ . Then  $U_5 \in M$ . Let  $k \in \text{Aut}(U_5) = K$  be such that  $u_2^k = u_2$ . The corresponding permutation on the coordinates induced by  $k$  must be the identity and the corresponding scalars not 1 can only appear possibly in the first coordinate. Thus

$|C_K(u_2)| \leq 2$ . Since there are 4 nonzero vectors in  $\langle u_2 \rangle$ ,  $[K_{\langle u_2 \rangle} : C_K(u_2)] \leq 4$ . Since there are 6 1-dimensional subspaces in  $U_5$  and  $w(u_1) \neq w(u_2)$ ,  $[K : K_{\langle u_2 \rangle}] \leq 5$ . Therefore  $|K| \leq 5 \cdot 4 \cdot 2$  and so  $|U_5^H| = \frac{|H|}{|K|} \geq \frac{2^8 \cdot 3 \cdot 5}{2^3 \cdot 5} = 2^5 \cdot 3$ . Hence  $M_5 = X_5^H \cup U_5^H$  and  $|U_5^H| = 2^5 \cdot 3$ .

Suppose  $n = 6$ . Then  $|M_6| = 2^3 \cdot 3 \cdot 13$  and  $|H| = 2^6 \cdot 6!$

Let  $X_6 = X_2 \oplus X_2 \oplus X_2$ . Thus  $\text{Aut}(X_6) = \text{Aut}(X_2) \vee S_3$  has order  $2^7 \cdot 3$  and  $|X_6^H| = 2^3 \cdot 3 \cdot 5$ .

Let  $U_6$  be the subspace generated by the  $u_1 = (1, 0, 0, 1, 2, 2)$ ,  $u_2 = (0, 1, 0, 2, 1, 3)$  and  $u_3 = (0, 0, 1, 2, 3, 1)$ . Then  $U_6 \in M$  and is indecomposable. So  $U_6 \notin X_6^H$ . Note that  $U_6 = \{(a, b, c, a + 2b + 2c, 2a + b - 2c, 2a - 2b + c) | a, b, c \in F\}$ . Let  $K = \text{Aut}(U_6)$ . Then  $K$  induces a permutation group  $\bar{K}$  on the 6 coordinate places  $\bar{1}, \bar{2}, \dots, \bar{6}$ . Since  $\{\pm 2\} \cap \{\pm 1\} = \emptyset$  in  $F$ ,  $\bar{K}_{\langle u \rangle}$  leaves invariant  $\{\bar{2}, \bar{3}\}$ ,  $\{\bar{1}, \bar{4}\}$ ,  $\{\bar{5}, \bar{6}\}$ . Let  $(y_1, \dots, y_6)$  be a vector in  $V$ . If  $k \in C_K(u_1)$ , then  $k$  permutes  $y_2, y_3$  with possible multiplication by a factor  $\pm 1$ , and permutes  $y_1, y_4$ , and permutes  $y_5, y_6$ . Suppose  $\bar{2}^{\bar{k}} = \bar{2}$ . Then  $\bar{3}^{\bar{k}} = \bar{3}$ . If  $y_1^k = y_1$ , then  $y_4^k = y_4$ . This implies  $u_2^k = u_2$  which forces  $y_2^k = y_2$ . Also  $u_3^k = u_3$  under the present assumption on  $k$ . Hence  $y_t^k = y_t$ , for  $t = 3, 5, 6$ . Therefore  $k$  is the identity linear transformation in this case. If  $y_1^k = y_4$ , then  $y_4^k = y_1$ . Thus  $u_2^k = (2, \pm 1, 0, 0, *, *) \in U_6$ , which implies  $u_2^k = (2, -1, 0, 0, 3, 1)$  by the description of the vectors in  $U_6$ . Hence  $k$  interchanges  $y_5$  and  $y_6$ , and  $y_2^k = -y_2$ . Since  $u_3^k = (2, 0, \pm 1, 0, 1, -2) \in U_6$ ,  $u_3^k = (2, 0, -1, 0, 1, -2)$ . Hence  $y_3^k = -y_3$ . Therefore if  $\bar{2}^{\bar{k}} = \bar{2}$ , then either  $k$  is the identity linear transformation or  $y_1^k = y_4$ ,  $y_4^k = y_1$ ,  $y_2^k = -y_2$ ,  $y_3^k = -y_3$ ,  $y_5^k = y_6$ , and  $y_6^k = y_5$ . Next suppose  $\bar{2}^{\bar{k}} = \bar{3}$ . If  $y_1^k = y_4$ , then  $y_4^k = y_1$  and so  $u_2^k = (2, 0, \pm 1, 0, **, *) \in U_6$ . This implies  $u_2^k = (2, 0, -1, 0, 1, 3)$  which shows  $y_5^k = y_5$  and  $y_6^k = y_6$ . Hence  $u_3^k = (2, \pm 1, 0, 0, -2, 1)$  which forces  $y_3^k = -y_3$ . If  $y_1^k = y_1$ , then  $y_4^k = y_4$ . Thus  $u_2^k = (0, 0, \pm 1, 2, *, *)$  which implies  $u_2^k = u_3$ . Hence  $y_2^k = y_3$ ,  $y_5^k = y_6$  and  $y_6^k = y_5$ . So



$u_3^k = (0, \pm 1, 0, 2, 1, -2)$  which implies  $Y_3^k = Y_2$ . This shows that  $C_K(u_1) = \langle \pi, \delta \rangle$  is an elementary abelian 2-group of order 4, where  $\pi$  interchanges  $Y_1$  and  $Y_4$ ,  $Y_5$  and  $Y_6$ , also  $Y_2^\pi = -Y_2$  and  $Y_3^\pi = -Y_3$ ; and  $\delta$  fixes  $Y_1, Y_4$ , interchanges  $Y_2$  and  $Y_3, Y_5$  and  $Y_6$ . Define  $\alpha$  in the following way: it interchanges  $Y_1$  and  $Y_2, Y_4$  and  $Y_5$ , also  $Y_3^\alpha = -Y_3, Y_6^\alpha = -Y_6$ . Since  $u_1^\alpha = u_2, u_2^\alpha = u_1$  and  $u_3^\alpha = -u_3$ , we have  $\alpha \in K$ . Let  $\beta = \alpha\delta$ . Then  $\beta^3 = -I$  and  $\beta$  has order 6. Let  $g \in K$  such that  $g$  induces identity on  $U_6$ . Then  $g \in C_K(u_1)$ . From  $u_2^g = u_2$  and  $u_3^g = u_3$ , we see that  $g = 1$ . Therefore  $K$  is isomorphic to a subgroup of  $GL(U_6) \cong GL(3, 5)$ . Hence  $|K|$  divides  $5^3 \cdot 3! \cdot 3 \cdot 2^7$ . Since  $|K| \mid |H| = 2^6 \cdot 6!$ ,  $|K|$  divides  $5 \cdot 3 \cdot 2^7$ . Let  $\Gamma = \{1\text{-dimensional subspaces of } U_6\}$ . Then  $|\Gamma| = 31$  and  $K$  acts on  $\Gamma$ . Let  $e_1 = u_1 + u_2 + u_3 = (1, 1, 1, 0, 1, 1)$ ,  $e_2 = e_1^\beta = (1, -1, 1, 1, -1, 0)$ ,  $e_3 = e_2^\beta = (-1, -1, 1, -1, 0, 1)$ , and  $e_4 = e_1^\pi = (0, -1, -1, 1, 1, 1)$ . Since  $\beta \notin K_{\langle e_4 \rangle}$ ,  $|\langle e_1 \rangle^K| \geq 6$ . Let  $f_1 = u_1 - u_2 - u_3 = (1, -1, -1, -3, 3, 3)$ . Let  $f_2 = f_1^\pi = (-3, 1, 1, 1, 3, 3)$ . Then  $(f_2)^{\beta^2} = (-1, 3, 1, -3, -1, 3) \notin \langle f_1 \rangle, \langle f_2 \rangle$ . It is easy to see that  $|\langle f_1 \rangle^K| \geq 6$ . From  $w(u_1) = 4, w(e_1) = 5, w(f_1) = 6$ , we get  $\langle u_1 \rangle^K \neq \langle e_1 \rangle^K \neq \langle f_1 \rangle^K \neq \langle u_1 \rangle^K$ . Thus  $[K : K_{\langle u_1 \rangle}] = |\langle u_1 \rangle^K| \leq 31 - (6+6) = 19$ . Combining the facts that 3 divides  $|K|$  to the first power,  $|K| \mid 5 \cdot 3 \cdot 2^7$ , and  $K_{\langle u_1 \rangle}$  is a 2-group, we see that  $15 \mid [K : K_{\langle u_1 \rangle}]$ . Therefore  $[K : K_{\langle u_1 \rangle}] = 15$ . Since  $K_{\langle u_1 \rangle}$  acts on the 4 non-zero vectors of  $\langle u_1 \rangle$ ,  $[K_{\langle u_1 \rangle} : C_K(u_1)] \leq 4$ . Hence  $|K| = 15 \cdot |K_{\langle u_1 \rangle}| \leq 15 \cdot 2^4$ . Thus  $|U_6^H| = \frac{|H|}{|K|} \geq \frac{2^6 \cdot 6!}{15 \cdot 2^4} = 3 \cdot 2^6 = |M \setminus X_6^H|$ . Therefore  $M_6 = X_6^H \cup U_6^H$  and  $|K| = 2^4 \cdot 3 \cdot 5$ . Define  $\theta$  by  $Y_1^\theta = -Y_6, Y_6^\theta = Y_1, Y_4^\theta = -Y_5,$

$Y_5^\theta = Y_4, Y_2^\theta = -Y_3, Y_3^\theta = Y_2$ . Then  $\theta \in K$  and  $\theta^2 = -I$ . Since  $\langle \theta \rangle$  acts transitively on the 4 non-zero vectors of  $\langle u_1 \rangle$ ,  $K_{\langle u_1 \rangle} = \langle \theta \rangle \cdot C_K(u_1)$ . This shows that  $\bar{K} = K / \langle -I \rangle$ . Since  $\bar{K}$  acts transitively on  $\{\bar{1}, \dots, \bar{6}\}$  and  $|\bar{K}| = 120$  and  $\bar{\theta} = (16)(23)(45)$ , we have  $\bar{K} \cong S_5$ . This completes the analysis for  $n = 6$ .

Suppose  $n = 7$ . Then  $|M_7| = 2^3 \cdot 3^3 \cdot 7 \cdot 13$  and  $|H| = 2^7 \cdot 7!$ . Let  $X_7 = X_6 \oplus \langle 0 \rangle, C_2 = U_6 \oplus \langle 0 \rangle, C_3 = U_5 \oplus X_2$ . Then  $|\text{Aut}(X_7)| = 2^8 \cdot 3, |\text{Aut}(C_2)| = 2^5 \cdot 3 \cdot 5, |\text{Aut}(C_3)| = 2^5 \cdot 5$  and so  $|X_7^H| = 2^3 \cdot 3 \cdot 5 \cdot 7, |C_2^H| = 2^6 \cdot 3 \cdot 7, |C_3^H| = 2^6 \cdot 3^2 \cdot 7$ . Let  $U_7$  be the linear code generated by  $t_1 = (1, 0, 0, 1, 0, 2, 2), t_2 = (0, 1, 0, 1, 2, 0, 2)$  and  $t_3 = (0, 0, 1, 1, 2, 2, 0)$ . Then  $U_7 \in M$ . Let  $G = \text{Aut}(U_7)$  and  $\bar{G}$  the permutation group induced by  $G$  on the 7 coordinate places  $\{\bar{1}, \dots, \bar{7}\}$ . Let  $g \in G_{\langle t_1 \rangle}$ . Since  $g$  permutes the 3 zero coordinates of  $v_1, \{\bar{2}, \bar{3}, \bar{5}\}$  is  $\bar{g}$  invariant. If  $g \in G_{\langle t_2 \rangle} \cap G_{\langle t_3 \rangle}$ , then  $\{\bar{1}, \bar{3}, \bar{6}\}$  and  $\{\bar{1}, \bar{2}, \bar{7}\}$  are  $\bar{g}$  invariant by the same argument. Hence  $\bar{g}$  fixes  $\bar{1} = \{\bar{1}, \bar{2}, \bar{7}\} \cap \{\bar{1}, \bar{3}, \bar{6}\}, \bar{2} = \{\bar{1}, \bar{2}, \bar{7}\} \cap \{\bar{2}, \bar{3}, \bar{5}\}$ , and  $\bar{3} = \{\bar{2}, \bar{3}, \bar{5}\} \cap \{\bar{1}, \bar{3}, \bar{6}\}$ . This implies that  $\bar{g}$  also fixes  $\bar{7}, \bar{6}, \bar{5}$  which forces  $\bar{g}$  to be the identity permutation. Easy calculation shows that  $g = \pm I_7$ . Note that  $U_7$  is indecomposable. A typical vector in  $U_7$  has the form  $(a, b, c, a+b+c, 2(b+c), 2(a+c), 2(a+b))$ . In finding  $\langle v \rangle$  with  $w(v) = 7$ , we may assume the first coordinate of  $v$  to be 1. A direct calculation shows that there are 4 possibilities:  $v_1 = (1, 1, 1, 3, 4, 4, 4), v_2 = (1, 1, 2, 4, 1, 1, 4), v_3 = (1, 2, 1, 4, 1, 4, 1)$  and  $v_4 = (1, 3, 3, 2, 2, 3, 3)$ . Therefore  $w_7(U_7) = 16$  and  $G$  acts on  $\{\langle v_i \rangle \mid i=1, 2, 3, 4\}$ . If  $g \in G$  and  $\langle v_i \rangle^g = \langle v_i \rangle$  for  $i=1, 2, 3$ , then  $\langle t_i \rangle^g = \langle t_i \rangle$  for  $i=1, 2, 3$ . This implies  $g = \pm I_7$ .



and so  $|G||S_4| \cdot 2 = 2^4 \cdot 3$ . Hence  $|U_7^H| \geq 2^7 \cdot 7! / 2^4 \cdot 3 = 2^7 \cdot 3 \cdot 5 \cdot 7 = |M_7 \setminus \{X_7^H \cup C_2^H \cup C_3^H\}|$ . Therefore  $M_7 = X_7^H \cup C_2^H \cup C_3^H \cup U_7^H$  and  $|\text{Aut}(U_7)| = 2^4 \cdot 3$ .

Suppose  $n = 8$ . Then  $|M_8| = 2^4 \cdot 3^3 \cdot 7 \cdot 13$ , and  $|H| = 2^8 \cdot 8!$ . Let  $X \in M_8$ . In finding  $X^H$  we may assume that  $X$  is generated by  $(1, 0, 0, 0, \alpha_1)$ ,  $(0, 1, 0, 0, \alpha_2)$ ,  $(0, 0, 1, 0, \alpha_3)$ ,  $(0, 0, 0, 1, \alpha_4)$  where  $\alpha_i \in F^4$ , for  $1 \leq i \leq 4$ . Let  $Y = \langle (1, 0, 0, \alpha_2), (0, 1, 0, \alpha_3), (0, 0, 1, \alpha_4) \rangle$ . Then  $Y \in M_7$ . Without loss of generality, we may assume that  $Y = X_7$ ,  $C_2$ ,  $C_3$  or  $U_7$ . By 2.3, there are at most 4 orbits of  $H$  in  $M_7$ . If  $Y = X_7$ , then  $X \in X_8^H$ , where  $X_8 = X_2 \oplus X_2 \oplus X_2 \oplus X_2$ . Since  $\text{Aut}(X_8) = 2^8 \cdot 2^3 \cdot 3$ ,  $|X_8^H| = 2^4 \cdot 3 \cdot 5 \cdot 7$ . Let  $T_2 = U_6 \oplus X_2$ . Since  $w_2(U_6) = 0$ ,  $|\text{Aut}(T_2)| = (2^4 \cdot 3 \cdot 5) \cdot 4$ . Hence  $|T_2^H| = 2(2^6 \cdot 3 \cdot 7(1+3))$ . Clearly  $C_2$  can be embedded in  $T_2$ . Now  $U_6$  contains  $z_1 = (1, 1, 1, 0, 1, 1)$  and  $z_2 = (0, 1, 4, 0, 3, 2)$ . As  $\langle z_1, 0, 0 \rangle, \langle z_2, 0, 0 \rangle \cong U_5$  and  $\langle (0, 0, 0, 0, 0, 1, 2) \rangle \cong X_2$ ,  $C_3 = U_5 \oplus X_2$  can be embedded in  $T_2$ . Hence  $M_8$  has at most 3 orbits of  $H$ . Let  $U_8 = \langle (1, 0, 0, 0, 1, 1, 1, 1), (0, 1, 0, 0, 1, 0, 2, 2), (0, 0, 1, 0, 1, 2, 0, 2), (0, 0, 0, 1, 1, 2, 2, 0) \rangle$ . Then  $U_8 \in M_8$  and  $U_7$  can be embedded in  $U_8$ . Also  $U_8$  is indecomposable. Hence  $M_8 = X_8^H \cup T_2^H \cup U_8^H$ , and  $|U_8^H| = 2^8 \cdot 3 \cdot 5 \cdot 7$  and  $|\text{Aut}(U_8)| = 2^7 \cdot 3$ .

A direct computation using computer's help in the case  $n = 7, 8$  yields the Hamming weight distribution of the representative of each orbit of  $M_n$  under  $H$ . This completes the proof of 3.2.

**3.3. Remark.** Even though  $H$  preserves the Hamming weight, it does not preserve the complete weight of a vector  $\text{comp}(v) = (\xi_0, \xi_1, \xi_2, \xi_3, \xi_4)$  for  $v = (v_1, \dots, v_n)$ , where  $\xi_i = |\{j | v_j = i\}|$

for  $0 \leq i \leq 4$  [See p. 142, 6]. The symmetric subgroup  $S$  of  $H$  does preserve the complete weight. Note that  $S \cong S_n$ .

It is easy to see that for  $2 \leq n \leq 4$ ,  $M_n$  is still an orbit of  $S$ . Also  $X_5^H = X_5^S$ .

Let  $R = S_{U_5}$ . Then  $C_R(u_2) = 1$ , and so  $|R| = [R : R_{\langle u_2 \rangle}] [R_{\langle u_2 \rangle}] \leq 5 \cdot 4$ . Therefore  $|U_5^S| \geq \frac{5!}{5 \cdot 4} = 6$ .

Let  $W = \langle \alpha_1, \alpha_2 \rangle$ , where  $\alpha_1 = (1, -1, 1, 1, 1)$  and  $\alpha_2 = (0, 4, 2, 3, 4)$ . Then  $W \in U_5^H$ . Any image of  $\alpha_2$  under  $S$  will have exactly 2 coordinates equal to 4. Since  $U_5 = \{(x, x+y, x+2y, x+3y, x+4y) | x, y \in F\}$ ,  $\alpha_2^S \notin U_5$ . Hence  $W \notin U_5^S$ . Let  $A = S_W$ . If  $\sigma \in A$ , then  $(\alpha_1, \alpha_1^\sigma) = 0$ , which implies that  $\bar{2}^\sigma = \bar{2}$  (the second coordinate place). Suppose  $\sigma \in A_{\langle \alpha_2 \rangle}$ . Then from  $\bar{2}^\sigma = \bar{2}$  we get  $\sigma \in C_A(\alpha_2)$  which implies that  $\sigma_1 = I$ . Hence  $|A_{\langle \alpha_2 \rangle}| = 1$ . Since  $\bar{2}^A = \bar{2}$ ,  $5 \nmid |A|$ . As  $A$  acts on the 6 1-dimensional subspaces of  $W$  and  $w(\alpha_2) \neq w(\alpha_1)$ ,  $|A| = |A : A_{\langle \alpha_2 \rangle}| = |\langle \alpha_2 \rangle^A| \leq 4$ , as  $5 \nmid |A|$ . Therefore  $|W^S| \geq \frac{5!}{4} = 2 \cdot 3 \cdot 5$ .

Let  $Z = \langle \beta_1, \beta_2 \rangle$ , where  $\beta_1 = (1, -1, -1, 1, 1)$  and  $\beta_2 = (0, 4, 3, 3, 4)$ . Then  $Z \in U_5^H$ . As in the case of  $\alpha_1$ , we see that  $\beta_2^S \notin U_5$ . Hence  $Z \notin U_5^S$ . Since  $W_2 = \{(x, -x+4y, x+2y, x+3y, x+4y) | x, y \in F\}$ ,  $\beta_2^S \notin W_2$  as  $\beta_2$  has 2 distinct pairs of equal coordinates. Hence  $Z \notin W_2^S$ . Let  $B = S_Z$ . Since  $(\beta_1, \beta_1) = 0$ ,  $B$  fixes  $\{\bar{2}, \bar{3}\}$ . Suppose  $b \in B_{\langle \beta_2 \rangle}$ . Then  $\bar{1}^b = \bar{1}$ . Assume that  $\bar{2}^b = \bar{3}$ . So  $\bar{3}^b = \bar{2}$ . Then  $\beta_2^b = (0, 3, 4, *, *) = \lambda \beta_2$  implies that  $3 = \lambda 4$  or  $\lambda = 4 \cdot 3 = 2$ . Hence  $\beta_2^b = (0, 3, 1, 1, 3)$  which is impossible as no coordinate of  $\beta_2$  equals to 1. Hence  $\bar{2}^b = \bar{2}$  and  $\bar{3}^b = \bar{3}$ . Thus  $b \in C_B(\beta_2)$ , and this implies that  $b = I$ . Therefore



$|B_{\langle \beta_2 \rangle}| = 1$ . Now  $Z = \{(x, -x-y, -x-2y, x+3y, x+4y) \mid x, y \in F\}$  contains  $(1, 0, 1, 3, 2) = z_2$ ,  $(1, 2, 0, 2, 4) = z_3$  and  $(1, 3, 2, 4, 0) = z_5$ . By inspection  $\langle \beta_2 \rangle^B \notin \langle \beta_1 \rangle \cup \bigcup_{i=2}^5 \langle z_i \rangle$ . Hence  $|B| = [B : B_{\langle \beta_2 \rangle}] = |\langle \beta_2 \rangle^B| \leq 2$ .

Therefore  $|Z^S| \geq \frac{5!}{2} = 2^2 \cdot 3 \cdot 5$ . Since  $|U_5^S \cup W^S \cup Z^S| \geq 6 + 2 \cdot 3 \cdot 5 + 2^2 \cdot 3 \cdot 5 = 2^5 \cdot 3 = |U_5^H|$ , we have  $U_5^H = U_5^S \cup W^S \cup Z^S$ ,  $|U_5^S| = 6$ ,  $|W^S| = 2 \cdot 3 \cdot 5$ ,  $|Z^S| = 2^2 \cdot 3 \cdot 5$ ,  $|S_{U_5}| = 20$ ,  $|S_W| = 4$  and  $|S_Z| = 2$ . Note also that the complete weight distribution of  $U_5$  and  $W$  are different.

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Chat Yin Ho  
Department of Mathematics  
University of Toronto  
Scarborough College  
Toronto, Ontario  
Canada