Substitutes and Complements in Network Flows Viewed as Discrete Convexity ¹

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Abstract

A new light is shed on "substitutes and complements" in the maximum weight circulation problem with reference to the concepts of L-convexity and M-convexity in the theory of discrete convex analysis. This provides us with a deeper understanding of the relationship between convexity and submodularity in combinatorial optimization.

Keywords: network flow, submodularity, convexity, combinatorial optimization

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1 Introduction

The relationship between convexity and submodularity has been discussed in the literature of combinatorial optimization (see [1, 2, 3, 10]). In this paper, we address this issue with reference to "substitutes and complements in network flows" discussed by Gale–Politof [5], and show that the concepts of L-convexity and M-convexity due to Murota [11, 12] help us better understand the relationship between convexity and submodularity.

We consider a network flow problem. Let G = (V, A) be a directed graph with vertex set V and arc set A. For each arc $a \in A$, we are given a nonnegative capacity c(a) for flow and a weight w(a) per unit flow. The maximum weight circulation problem is to find a flow $\xi = (\xi(a) \mid a \in A)$ that maximizes the total weight $\sum_{a \in A} w(a)\xi(a)$ subject to the capacity (feasibility) constraint:

$$0 \le \xi(a) \le c(a) \qquad (a \in A)$$

and the conservation constraint:

$$\sum \{\xi(a) \mid a \text{ leaves } v\} - \sum \{\xi(a) \mid a \text{ enters } v\} = 0 \qquad (v \in V).$$

$$(1.1)$$

We denote by F the maximum weight of a feasible circulation, i.e.,

$$F = \max\{w^{\mathrm{T}}\xi \mid N\xi = \mathbf{0}, \ \mathbf{0} \le \xi \le c\},\tag{1.2}$$

where $N\xi = \mathbf{0}$ represents the conservation constraint (1.1).

Our concern here is how the weight F depends on the problem parameters (w, c). Namely, we are interested in the function F = F(w, c) in $w \in \mathbf{R}^A$ and $c \in \mathbf{R}_+^A$. We first look at convexity and concavity.

Proposition 1. F is convex in w and concave in c.

Proof. The function F = F(w, c) given by (1.2) is the maximum of linear functions in w and hence convex in w. By linear programming duality, we obtain an alternative expression $F = \min\{c^{\mathrm{T}}\eta \mid N^{\mathrm{T}}p + \eta \geq w, \ \eta \geq \mathbf{0}\}$, which shows the concavity of F in c.

We next consider submodularity and supermodularity. A function $f: \mathbf{R}^n \to \mathbf{R} \cup \{+\infty\}$ is said to be submodular if

$$f(x) + f(y) \ge f(x \lor y) + f(x \land y) \qquad (x, y \in \mathbf{R}^n),$$

and supermodular if

$$f(x) + f(y) \le f(x \lor y) + f(x \land y)$$
 $(x, y \in \mathbf{R}^n),$

where $x \vee y$ and $x \wedge y$ are defined by

$$(x \lor y)(i) = \max\{x(i), y(i)\}, \qquad (x \land y)(i) = \min\{x(i), y(i)\} \qquad (i = 1, 2, \dots, n).$$

With economic terms of *substitutes* and *complements* we have the following correspondences:

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f is submodular \iff goods are substitutes,

f is supermodular \iff goods are complements,
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where f is interpreted as representing a utility function.

Two arcs are said to be "parallel" if every (undirected) simple cycle containing both of them orients them in the opposite direction, and "series" if every (undirected) simple cycle containing both of them orients them in the same direction. A set of arcs is said to be "parallel" if it consists of pairwise "parallel" arcs, and "series" if it consists of pairwise "series" arcs. With notations $w_P = (w(a) \mid a \in P), c_P = (c(a) \mid a \in P), w_S = (w(a) \mid a \in S),$ and $c_S = (c(a) \mid a \in S),$ the following statements hold true.

Theorem 2 (Gale-Politof [5]). Let P be a "parallel" arc set and S a "series" arc set.

- (i) F is submodular in w_P and in c_P .
- (ii) F is supermodular in w_S and in c_S .

See [6, 7, 8] for some extensions and generalizations of this result.

Combination of Proposition 1 and Theorem 2 yields that

$$F$$
 is convex and submodular in w_P ,
 F is concave and submodular in c_P ,
 F is convex and supermodular in w_S ,
 F is concave and supermodular in c_S .

Thus all combinations of convexity/concavity and submodularity/supermodularity arise in our network flow problem. This demonstrates that convexity and submodularity are mutually independent properties.

Although convexity and submodularity are mutually independent, the combinations of convexity/concavity and submodularity/supermodularity in (1.3) are not accidental phenomena but logical consequences that can be explained in terms of L-convexity and M-convexity.

The concepts of M-convex and L-convex functions are introduced by Murota [11, 12] (see also [13, 14]), aiming to identify the well-behaved structure in (nonlinear) combinatorial optimization. These concepts were originally introduced for functions over the integer lattice; subsequently, their variants called M^{\natural} -convexity and L^{\natural} -convexity were introduced by Murota-Shioura [15] and by Fujishige-Murota [4], respectively. Recently, Murota-Shioura [16] extended these concepts to polyhedral convex functions defined over the real space.

A polyhedral convex function $f: \mathbf{R}^n \to \mathbf{R} \cup \{+\infty\}$ is said to be M-convex if dom $f \neq \emptyset$ and f satisfies (M-EXC):

(M-EXC)
$$\forall x, y \in \text{dom } f, \ \forall i \in \text{supp}^+(x-y), \ \exists j \in \text{supp}^-(x-y), \ \exists \alpha_0 > 0:$$

$$f(x) + f(y) \ge f(x - \alpha(\chi_i - \chi_j)) + f(y + \alpha(\chi_i - \chi_j)) \qquad (\forall \alpha \in [0, \alpha_0]),$$

where

$$\operatorname{dom} f = \{x \in \mathbf{R}^n \mid f(x) < +\infty\},\$$

$$\operatorname{supp}^+(x) = \{i \mid x(i) > 0\}, \quad \operatorname{supp}^-(x) = \{i \mid x(i) < 0\} \qquad (x \in \mathbf{R}^n),\$$

$$\chi_i \in \{0, 1\}^n \text{: the } i\text{-th unit vector } (i = 1, 2, \dots, n),\$$

$$[0, \alpha_0] = \{\alpha \in \mathbf{R} \mid 0 \le \alpha \le \alpha_0\}.$$

A polyhedral convex function $f: \mathbf{R}^n \to \mathbf{R} \cup \{+\infty\}$ is said to be \mathbf{M}^{\natural} -convex if the function $\widehat{f}: \mathbf{R}^n \times \mathbf{R} \to \mathbf{R} \cup \{+\infty\}$ defined by

$$\widehat{f}(x,x_0) = \begin{cases} f(x) & ((x,x_0) \in \mathbf{R}^n \times \mathbf{R}, \ x_0 = -\sum_{i=1}^n x(i)), \\ +\infty & \text{(otherwise)} \end{cases}$$

is M-convex. On the other hand, a polyhedral convex function $g: \mathbf{R}^n \to \mathbf{R} \cup \{+\infty\}$ is said to be L-convex if dom $g \neq \emptyset$ and g satisfies (LF1) and (LF2):

- (LF1) g is submodular,
- (**LF2**) $\exists r \in \mathbf{R} \text{ such that } q(p + \lambda \mathbf{1}) = q(p) + \lambda r \quad (\forall p \in \mathbf{R}^n, \ \lambda \in \mathbf{R});$

it is called L^{\natural}-convex if the function $\widehat{g}: \mathbf{R}^n \times \mathbf{R} \to \mathbf{R} \cup \{+\infty\}$ defined by

$$\widehat{g}(p, p_0) = g(p - p_0 \mathbf{1})$$
 $((p, p_0) \in \mathbf{R}^n \times \mathbf{R})$

is L-convex.

The main aim of this paper is to show that the function F defined by (1.2) is endowed with L^{\natural} -convexity and M^{\natural} -convexity, as follows. The proof is given in Section 2.

Theorem 3. Let P be a "parallel" arc set and S a "series" arc set.

- (i) F is L^{\natural} -convex in w_P and M^{\natural} -concave in c_P .
- (ii) F is M^{\natural} -convex in w_S and L^{\natural} -concave in c_S .

In general, L^{\natural} -convexity implies submodularity by definition, whereas M^{\natural} -convexity implies supermodularity [16, Theorem 4.24].

Theorem 4.

- (i) An L^{\natural} -convex function is submodular.
- (ii) An M^{\dagger}-convex function is supermodular.

Accordingly, L^{\natural} -concavity implies supermodularity and M^{\natural} -concavity submodularity. With the aid of these general results on L^{\natural} -convex and M^{\natural} -convex functions, Theorem 3 above provides us with a somewhat deeper understanding of (1.3). Namely, it is understood that

F is L^{\natural} -convex, hence convex and submodular, in w_P , F is M^{\sharp} -concave. hence concave and submodular, in c_P , $F \text{ is } \mathbf{M}^{\sharp}\text{-convex},$ hence convex and supermodular, w_S , F is L^{\natural} -concave, hence concave and supermodular, c_S .

It is left for future research to consider the results of [6, 7, 8] from the viewpoint of discrete convexity.

2 Proofs

This section gives the proof of Theorem 3. We start with basic properties of "parallel" and "series" arc sets that we use in the proof. Let us call $\pi: A \to \{0, \pm 1\}$ a *circuit* if $\partial \pi = \mathbf{0}$ and the set $\sup_{x \to 0} \mathbf{0} = \mathbf{0}$ and the set $\sup_{x \to 0} \mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and the set $\mathbf{0} = \mathbf{0}$ are $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{0} = \mathbf{0}$ are supp- $\mathbf{0} = \mathbf{0}$ and $\mathbf{$

Proposition 5. Let π be a circuit.

- (i) $|\operatorname{supp}^+(\pi) \cap P| \le 1$ and $|\operatorname{supp}^-(\pi) \cap P| \le 1$ for a "parallel" arc set P.
- (ii) $|\operatorname{supp}^+(\pi) \cap S| = 0$ or $|\operatorname{supp}^-(\pi) \cap S| = 0$ for a "series" arc set S.

Proposition 6. Let S be a "series" arc set, and π_1 and π_2 be circuits. If $\operatorname{supp}^+(\pi_1) \cap \operatorname{supp}^+(\pi_2) \cap S \neq \emptyset$, there exists a circuit π such that $\operatorname{supp}^+(\pi) \subseteq \operatorname{supp}^+(\pi_1) \cup \operatorname{supp}^+(\pi_2)$, $\operatorname{supp}^-(\pi) \subseteq \operatorname{supp}^-(\pi_1) \cup \operatorname{supp}^-(\pi_2)$, and $\operatorname{supp}^+(\pi) \cap S = (\operatorname{supp}^+(\pi_1) \cup \operatorname{supp}^+(\pi_2)) \cap S$.

Proof. Suppose $a \in (\operatorname{supp}^+(\pi_2) \setminus \operatorname{supp}^+(\pi_1)) \cap S$. By elementary graph argument we can find a circuit π' such that $\operatorname{supp}^+(\pi') \subseteq \operatorname{supp}^+(\pi_1) \cup \operatorname{supp}^+(\pi_2)$, $\operatorname{supp}^-(\pi') \subseteq \operatorname{supp}^-(\pi_1) \cup \operatorname{supp}^-(\pi_2)$, and $\operatorname{supp}^+(\pi') \cap S \supseteq (\operatorname{supp}^+(\pi_1) \cap S) \cup \{a\}$. Repeating this we can find π .

The main technical tool in the proof is the *conformal decomposition* (see, e.g., [9, 17]) of a circulation ξ , which is a representation of ξ as a positive sum of circuits conformal to ξ , i.e.,

$$\xi = \sum_{i=1}^{m} \beta_i \pi_i,$$

where $\beta_i > 0$ and $\pi_i : A \to \{0, \pm 1\}$ is a circuit with $\operatorname{supp}^+(\pi_i) \subseteq \operatorname{supp}^+(\xi)$ and $\operatorname{supp}^-(\pi_i) \subseteq \operatorname{supp}^-(\xi)$ for $i = 1, 2, \ldots, m$.

2.1 Proof of L^{\dagger}-convexity in w_P

L^{\dagger}-convexity of F in w_P is equivalent to submodularity of $F(w-w_0\chi_P,c)$ in (w_P,w_0) , which in turn is equivalent to

$$F(w + \lambda \chi_a, c) + F(w + \mu \chi_b, c) \ge F(w, c) + F(w + \lambda \chi_a + \mu \chi_b, c), \tag{2.1}$$

$$F(w + \lambda \chi_a, c) + F(w - \mu \chi_P, c) \ge F(w, c) + F(w + \lambda \chi_a - \mu \chi_P, c)$$
(2.2)

for $a, b \in P$ with $a \neq b$ and $\lambda, \mu \in \mathbf{R}_+$, where $\chi_P \in \{0, 1\}^A$ denotes the characteristic vector of $P \subseteq A$. To show (2.1) let ξ and $\overline{\xi}$ be optimal circulations for w and $w + \lambda \chi_a + \mu \chi_b$. We can establish (2.1) by constructing feasible circulations ξ_a and ξ_b such that

$$\xi_a + \xi_b = \xi + \overline{\xi}, \qquad \lambda[\xi_a(a) - \overline{\xi}(a)] + \mu[\xi_b(b) - \overline{\xi}(b)] \ge 0,$$
 (2.3)

since this implies

$$(w + \lambda \chi_a)^{\mathrm{T}} \xi_a + (w + \mu \chi_b)^{\mathrm{T}} \xi_b \ge w^{\mathrm{T}} \xi + (w + \lambda \chi_a + \mu \chi_b)^{\mathrm{T}} \overline{\xi}$$

of which the left-hand side is bounded by $F(w + \lambda \chi_a, c) + F(w + \mu \chi_b, c)$ and the right-hand side is equal to $F(w, c) + F(w + \lambda \chi_a + \mu \chi_b, c)$. If $\overline{\xi}(a) \leq \xi(a)$, we can take $\xi_a = \xi$ and $\xi_b = \overline{\xi}$ to meet (2.3).

If $\overline{\xi}(b) \leq \xi(b)$, we can take $\xi_a = \overline{\xi}$ and $\xi_b = \xi$ to meet (2.3). Otherwise, we make use of the conformal decomposition $\overline{\xi} - \xi = \sum_{i=1}^m \beta_i \pi_i$. Since $a \in \operatorname{supp}^+(\overline{\xi} - \xi)$ we may assume $\pi_i(a) = 1$ for $i = 1, 2, \ldots, \ell$ and $\pi_i(a) = 0$ for $i = \ell + 1, \ell + 2, \ldots, m$. We have $\pi_i(b) = 0$ for $i = 1, 2, \ldots, \ell$ by Proposition 5 (i), since P is "parallel" and $\{a, b\} \subseteq \operatorname{supp}^+(\overline{\xi} - \xi)$. Then $\xi_a = \xi + \sum_{i=1}^\ell \beta_i \pi_i$ and $\xi_b = \xi + \sum_{i=\ell+1}^m \beta_i \pi_i$ are feasible circulations that satisfy (2.3).

To show (2.2) let ξ and $\overline{\xi}$ be optimal circulations for w and $w + \lambda \chi_a - \mu \chi_P$. We can establish (2.2) by constructing feasible circulations ξ_a and ξ_P such that

$$\xi_a + \xi_P = \xi + \overline{\xi}, \qquad \lambda[\xi_a(a) - \overline{\xi}(a)] + \mu[\sum_{a' \in P} \overline{\xi}(a') - \sum_{a' \in P} \xi_P(a')] \ge 0$$
 (2.4)

since this implies

$$(w + \lambda \chi_a)^{\mathrm{T}} \xi_a + (w - \mu \chi_P)^{\mathrm{T}} \xi_P \ge w^{\mathrm{T}} \xi + (w + \lambda \chi_a - \mu \chi_P)^{\mathrm{T}} \overline{\xi}.$$

If $\overline{\xi}(a) \leq \xi(a)$, we can take $\xi_a = \xi$ and $\xi_P = \overline{\xi}$ to meet (2.4). Otherwise we use the conformal decomposition $\overline{\xi} - \xi = \sum_{i=1}^m \beta_i \pi_i$, in which we assume $\pi_i(a) = 1$ for $i = 1, 2, ..., \ell$ and $\pi_i(a) = 0$ for $i = \ell + 1, \ell + 2, ..., m$. Since P is "parallel" we have $|\sup^-(\pi_i) \cap P| \leq 1$ by Proposition 5 (i), and hence $\sum_{a' \in P} \pi_i(a') \geq 0$ for $i = 1, 2, ..., \ell$. Therefore, $\xi_a = \xi + \sum_{i=1}^{\ell} \beta_i \pi_i$ and $\xi_P = \xi + \sum_{i=\ell+1}^{m} \beta_i \pi_i$ are feasible circulations with the properties in (2.4).

2.2 Proof of M^{\natural} -concavity in c_P

 M^{\natural} -convexity of a function $f: \mathbf{R}^n \to \mathbf{R} \cup \{+\infty\}$ is characterized by the following property [16, Theorem 4.21]:

(M^{\(\pi\)}-EXC)
$$\forall x, y \in \text{dom } f, \ \forall i \in \text{supp}^+(x-y), \ \exists j \in \text{supp}^-(x-y) \cup \{0\}, \ \exists \alpha_0 > 0:$$

$$f(x) + f(y) \ge f(x - \alpha(\chi_i - \chi_j)) + f(y + \alpha(\chi_i - \chi_j)) \qquad (\forall \alpha \in [0, \alpha_0]),$$

where $\chi_0 = \mathbf{0}$ by convention. We prove the M^{\dagger}-concavity of F in c_P by establishing (M^{\dagger}-EXC) for -F as a function in c_P . In our notation this reads as follows:

Let $c_1, c_2 \in \mathbf{R}_+^A$ be capacities with $c_1(a') = c_2(a')$ for all $a' \in A \setminus P$. For each $a \in \text{supp}^+(c_1 - c_2)$, there exist $b \in \text{supp}^-(c_1 - c_2) \cup \{0\}$ and a positive number α_0 such that

$$F(w, c_1) + F(w, c_2) \le F(w, c_1 - \alpha(\chi_a - \chi_b)) + F(w, c_2 + \alpha(\chi_a - \chi_b))$$
 $(\forall \alpha \in [0, \alpha_0]).$

Let ξ_1 and ξ_2 be optimal circulations for c_1 and c_2 , respectively. We shall find $\alpha_0 > 0$ and $b \in \text{supp}^-(c_1 - c_2) \cup \{0\}$ such that, for any $\alpha \in [0, \alpha_0]$, there exist circulations ξ'_1 and ξ'_2 such that

$$\xi_1' + \xi_2' = \xi_1 + \xi_2, \quad \mathbf{0} \le \xi_1' \le c_1 - \alpha(\chi_a - \chi_b), \quad \mathbf{0} \le \xi_2' \le c_2 + \alpha(\chi_a - \chi_b).$$
 (2.5)

If $\xi_1(a) < c_1(a)$, we can take $\alpha_0 = c_1(a) - \xi_1(a)$, b = 0, $\xi'_1 = \xi_1$ and $\xi'_2 = \xi_2$ to meet (2.5). Suppose $\xi_1(a) = c_1(a)$. We have $\xi_1(a) = c_1(a) > c_2(a) \ge \xi_2(a)$. Let π be a circuit such that $a \in \operatorname{supp}^+(\pi) \subseteq \operatorname{supp}^+(\xi_1 - \xi_2)$ and $\operatorname{supp}^-(\pi) \subseteq \operatorname{supp}^-(\xi_1 - \xi_2)$. Since P is "parallel" and $a \in \operatorname{supp}^+(\pi)$, we have $\operatorname{supp}^+(\pi) \cap P = \{a\}$ and $|\operatorname{supp}^-(\pi) \cap P| \le 1$ by Proposition 5 (i). If $|\operatorname{supp}^-(\pi) \cap P| = 1$, define b by $\{b\} = \operatorname{supp}^-(\pi) \cap P$; otherwise put b = 0. We can take $\alpha_0 > 0$ such that $\alpha_0 \le |\xi_1(a') - \xi_2(a')|$ for all $a' \in \operatorname{supp}^+(\pi) \cup \operatorname{supp}^-(\pi)$. Then $\xi'_1 = \xi_1 - \alpha\pi$ and $\xi'_2 = \xi_2 + \alpha\pi$ satisfy (2.5) if $\alpha \in [0, \alpha_0]$.

2.3 Proof of M^{\dagger}-convexity in w_S

We prove the M^{\dagger} -convexity of F in w_S by establishing (M^{\dagger} -EXC). In our notation this reads as follows:

Let $w_1, w_2 \in \mathbf{R}^A$ be weight vectors with $w_1(a') = w_2(a')$ for all $a' \in A \setminus S$. For each $a \in \operatorname{supp}^+(w_1 - w_2)$, there exist $b \in \operatorname{supp}^-(w_1 - w_2) \cup \{0\}$ and a positive number α_0 such that

$$F(w_1, c) + F(w_2, c) > F(w_1 - \alpha(\chi_a - \chi_b), c) + F(w_2 + \alpha(\chi_a - \chi_b), c)$$
 $(\forall \alpha \in [0, \alpha_0]).$

Let ξ_1 and ξ_2 be optimal circulations for w_1 and w_2 , respectively, with $\xi_1(a)$ minimum and $\xi_2(a)$ maximum.

Proposition 7. There exists $\alpha_0 > 0$ such that ξ_1 is optimal for $w_1 - \alpha \chi_a$ and ξ_2 is optimal for $w_2 + \alpha \chi_a$ for all $\alpha \in [0, \alpha_0]$.

Proof. For any circuit π such that $\pi(a) = -1$ and $\mathbf{0} \leq \xi_1 + \beta \pi \leq c$ for some $\beta > 0$, we have $w_1^{\mathrm{T}}(\xi_1 + \beta \pi) < w_1^{\mathrm{T}}\xi_1$ by the choice of ξ_1 . Let $\alpha_1 > 0$ be the minimum of $-w_1^{\mathrm{T}}\pi$ over all such circuits π . Then ξ_1 is optimal for $w_1 - \alpha \chi_a$ for all $\alpha \in [0, \alpha_1]$, since $(w_1 - \alpha \chi_a)^{\mathrm{T}}(\xi_1 + \beta \pi) \leq (w_1 - \alpha \chi_a)^{\mathrm{T}}\xi_1$ for any $\beta > 0$ and circuit π such that $\mathbf{0} \leq \xi_1 + \beta \pi \leq c$. Similarly, let $\alpha_2 > 0$ be the minimum of $-w_2^{\mathrm{T}}\pi$ over all circuits π such that $\pi(a) = 1$ and $\mathbf{0} \leq \xi_2 + \beta \pi \leq c$ for some $\beta > 0$. Then ξ_2 is optimal for $w_2 + \alpha \chi_a$ for all $\alpha \in [0, \alpha_2]$. Put $\alpha_0 = \min(\alpha_1, \alpha_2)$.

If $\xi_1(a) \geq \xi_2(a)$, we can take b = 0 in (M^{\delta}-EXC), since

$$F(w_1, c) + F(w_2, c) = w_1^{\mathrm{T}} \xi_1 + w_2^{\mathrm{T}} \xi_2$$

$$\geq (w_1 - \alpha \chi_a)^{\mathrm{T}} \xi_1 + (w_2 + \alpha \chi_a)^{\mathrm{T}} \xi_2 = F(w_1 - \alpha \chi_a, c) + F(w_2 + \alpha \chi_a, c),$$

where the last equality is by Proposition 7. In what follows we assume $\xi_1(a) < \xi_2(a)$.

By Proposition 5 (ii), we can impose further conditions on ξ_1 and ξ_2 that, for each $b \in S \setminus \{a\}$, $\xi_1(b)$ is maximum among all optimal ξ_1 for w_1 with $\xi_1(a)$ minimum, and $\xi_2(b)$ is minimum among all optimal ξ_2 for w_2 with $\xi_2(a)$ maximum.

Proposition 8. There exists $\alpha_0 > 0$ such that ξ_1 is optimal for $w_1 - \alpha(\chi_a - \chi_b)$ and ξ_2 is optimal for $w_2 + \alpha(\chi_a - \chi_b)$ for all $b \in S \setminus \{a\}$ and for all $\alpha \in [0, \alpha_0]$.

Proof. For any circuit π such that $\pi(a) - \pi(b) = -1$ for some $b \in S \setminus \{a\}$ and $\mathbf{0} \leq \xi_1 + \beta \pi \leq c$ for some $\beta > 0$, we have $w_1^{\mathrm{T}}(\xi_1 + \beta \pi) < w_1^{\mathrm{T}}\xi_1$ by the choice of ξ_1 . Let $\alpha_1 > 0$ be the minimum of $-w_1^{\mathrm{T}}\pi$ over all such circuits π . Then ξ_1 is optimal for $w_1 - \alpha(\chi_a - \chi_b)$ for all $\alpha \in [0, \alpha_1]$. Similarly, let $\alpha_2 > 0$ be the minimum of $-w_2^{\mathrm{T}}\pi$ over all circuits π such that $\pi(a) - \pi(b) = 1$ for some $b \in S \setminus \{a\}$ and $\mathbf{0} \leq \xi_2 + \beta \pi \leq c$ for some $\beta > 0$. Then ξ_2 is optimal for $w_2 + \alpha(\chi_a - \chi_b)$ for all $\alpha \in [0, \alpha_2]$. Put $\alpha_0 = \min(\alpha_1, \alpha_2)$.

Proposition 8 implies that for all $b \in S \setminus \{a\}$ we have

$$F(w_1, c) + F(w_2, c) - F(w_1 - \alpha(\chi_a - \chi_b), c) - F(w_2 + \alpha(\chi_a - \chi_b), c)$$

$$= w_1^{\mathrm{T}} \xi_1 + w_2^{\mathrm{T}} \xi_2 - (w_1 - \alpha(\chi_a - \chi_b))^{\mathrm{T}} \xi_1 - (w_2 + \alpha(\chi_a - \chi_b))^{\mathrm{T}} \xi_2$$

$$= \alpha[(\xi_2(b) - \xi_1(b)) - (\xi_2(a) - \xi_1(a))]. \tag{2.6}$$

We want to find $b \in \text{supp}^-(w_1 - w_2)$ for which (2.6) is nonnegative.

We make use of the conformal decomposition $\xi_2 - \xi_1 = \sum_{i=1}^m \beta_i \pi_i$. Since S is "series" we may assume, by Proposition 6, that

$$a \in \operatorname{supp}^+(\pi_1) \cap S \subseteq \operatorname{supp}^+(\pi_2) \cap S \subseteq \cdots \subseteq \operatorname{supp}^+(\pi_\ell) \cap S$$

and $\pi_i(a) = 0$ for $i = \ell + 1, \ell + 2, \dots, m$; then $\operatorname{supp}^-(\pi_i) \cap S = \emptyset$ for $i = 1, 2, \dots, \ell$.

Proposition 9. There exists $b \in (\operatorname{supp}^+(\pi_1) \cap S) \cap \operatorname{supp}^-(w_1 - w_2)$.

Proof. We have $w_1^T \pi_1 \leq 0$, since ξ_1 is optimal for w_1 and $\mathbf{0} \leq \xi_1 + \beta_1 \pi_1 \leq c$. Similarly, we have $-w_2^T \pi_1 \leq 0$. Hence

$$0 \ge (w_1 - w_2)^{\mathrm{T}} \pi_1 = \sum_{b \in S} (w_1(b) - w_2(b)) \pi_1(b) = \sum_{b \in \mathrm{supp}^+(\pi_1) \cap S} (w_1(b) - w_2(b)).$$

Since $w_1(a) - w_2(a) > 0$ in this summation, we must have $w_1(b) - w_2(b) < 0$ for some $b \in \operatorname{supp}^+(\pi_1) \cap S$.

For $b \in (\text{supp}^+(\pi_1) \cap S) \cap \text{supp}^-(w_1 - w_2)$ in Proposition 9, we have

$$\xi_2(b) - \xi_1(b) = \sum_{i=1}^{\ell} \beta_i + \sum_{i=\ell+1}^{m} \beta_i \pi_i(b) \ge \sum_{i=1}^{\ell} \beta_i = \xi_2(a) - \xi_1(a),$$

which shows the nonnegativity of (2.6).

2.4 Proof of L $^{\natural}$ -concavity in c_S

L^{\dagger}-concavity of F in c_S is equivalent to supermodularity of $F(w, c - c_0 \chi_S)$ in (c_S, c_0) , which in turn is equivalent to

$$F(w, c + \lambda \chi_a) + F(w, c + \mu \chi_b) \le F(w, c) + F(w, c + \lambda \chi_a + \mu \chi_b), \tag{2.7}$$

$$F(w, c + \lambda \chi_a) + F(w, c - \mu \chi_S) \le F(w, c) + F(w, c + \lambda \chi_a - \mu \chi_S)$$
(2.8)

for $a, b \in S$ with $a \neq b$ and $\lambda, \mu \in \mathbf{R}_+$, where $\chi_S \in \{0, 1\}^A$ denotes the characteristic vector of $S \subseteq A$. To show (2.7) let ξ_a and ξ_b be optimal circulations for $c + \lambda \chi_a$ and $c + \mu \chi_b$. We can establish (2.7) by constructing circulations ξ and $\overline{\xi}$ such that

$$\xi + \overline{\xi} = \xi_a + \xi_b, \quad \mathbf{0} \le \xi \le c, \quad \mathbf{0} \le \overline{\xi} \le c + \lambda \chi_a + \mu \chi_b.$$
 (2.9)

If $\xi_a(a) \leq c(a)$, we can take $\xi = \xi_a$ and $\overline{\xi} = \xi_b$ to meet (2.9). If $\xi_b(b) \leq c(b)$, we can take $\xi = \xi_b$ and $\overline{\xi} = \xi_a$ to meet (2.9). Otherwise, we have $\xi_a(a) > c(a) \geq \xi_b(a)$ and $\xi_a(b) \leq c(b) < \xi_b(b)$, and therefore $a \in \operatorname{supp}^+(\xi_a - \xi_b)$ and $b \in \operatorname{supp}^-(\xi_a - \xi_b)$. We make use of the conformal decomposition $\xi_a - \xi_b = \sum_{i=1}^m \beta_i \pi_i$, where we assume $\pi_i(a) = 1$ for $i = 1, 2, \ldots, \ell$ and $\pi_i(a) = 0$ for $i = \ell + 1, \ell + 2, \ldots, m$. We have $\pi_i(b) = 0$ for $i = 1, 2, \ldots, \ell$ by Proposition 5 (ii), since S is "series" and $a \in \operatorname{supp}^+(\xi_a - \xi_b)$ and $b \in \operatorname{supp}^-(\xi_a - \xi_b)$. Then $\xi = \xi_a - \sum_{i=1}^{\ell} \beta_i \pi_i$ and $\overline{\xi} = \xi_b + \sum_{i=1}^{\ell} \beta_i \pi_i$ satisfy (2.9).

To show (2.8) let ξ_a and ξ_S be optimal circulations for $c + \lambda \chi_a$ and $c - \mu \chi_S$. We can establish (2.8) by constructing circulations ξ and $\overline{\xi}$ such that

$$\xi + \overline{\xi} = \xi_a + \xi_S, \quad \mathbf{0} \le \xi \le c, \quad \mathbf{0} \le \overline{\xi} \le c + \lambda \chi_a - \mu \chi_S.$$
 (2.10)

If $\xi_a(a) \leq c(a)$, we can take $\xi = \xi_a$ and $\overline{\xi} = \xi_S$ to meet (2.10). Otherwise, we have $\xi_a(a) > c(a) \geq \xi_S(a)$, and therefore $a \in \text{supp}^+(\xi_a - \xi_S)$. We use the conformal decomposition $\xi_a - \xi_S = \sum_{i=1}^m \beta_i \pi_i$. Since S is "series" we may assume by Proposition 6 that

$$a \in \operatorname{supp}^+(\pi_1) \cap S \subseteq \operatorname{supp}^+(\pi_2) \cap S \subseteq \cdots \subseteq \operatorname{supp}^+(\pi_\ell) \cap S$$

and $\pi_i(a) = 0$ for $i = \ell + 1, \ell + 2, ..., m$; then $\sup_{i=1}^{\infty} (\pi_i) \cap S = \emptyset$ for $i = 1, 2, ..., \ell$. Noting $\sum_{i=1}^{\ell} \beta_i = \xi_a(a) - \xi_S(a) \ge \xi_a(a) - c(a)$, let k be the smallest integer with $\sum_{i=1}^{k} \beta_i \ge \xi_a(a) - c(a)$ and define $\beta' = [\xi_a(a) - c(a)] - \sum_{i=1}^{k-1} \beta_i$. Then $\xi = \xi_a - \sum_{i=1}^{k-1} \beta_i \pi_i - \beta' \pi_k$ and $\overline{\xi} = \xi_S + \sum_{i=1}^{k-1} \beta_i \pi_i + \beta' \pi_k$ satisfy (2.10), since $\xi(a) = \xi_a(a) - \sum_{i=1}^{k-1} \beta_i - \beta' = c(a)$, $\overline{\xi}(a) = \xi_S(a) + \sum_{i=1}^{k-1} \beta_i + \beta' = \xi_S(a) + \xi_a(a) - c(a) \le c(a) + \lambda - \mu$, and, for any $b \in \sup_{i=1}^{k} (\pi_k) \cap S$, we have

$$\overline{\xi}(b) = \xi_S(b) + \sum_{i=1}^{k-1} \beta_i \pi_i(b) + \beta' = \xi_S(b) + \sum_{i=1}^{k-1} \beta_i \pi_i(b) + [\sum_{i=k}^{\ell} \beta_i + \xi_S(a) - c(a)]$$

$$= [\xi_S(b) + \sum_{i=1}^{\ell} \beta_i \pi_i(b)] + \xi_S(a) - c(a) \le \xi_a(b) + \xi_S(a) - c(a) \le c(b) - \mu.$$

This completes the proof of Theorem 3.

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