

Knotted 3-Valent Graphs, Branched Braids, and Multiplication-Conjugation Relations in a Group

Victoria LEBED

OCAMI, Osaka

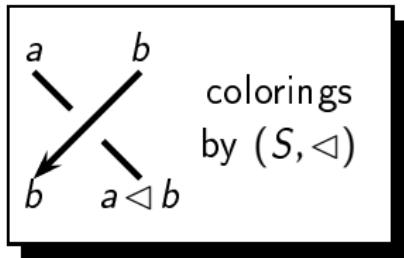
Intelligence of Low-dimensional Topology
Kyoto, May 21-23, 2014

Based on:

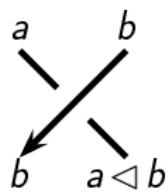
- ﴿ Qualgebras and Knotted 3-Valent Graphs, ArXiv, 2014
- ﴿ (With S. Kamada) Alexander and Markov Theorems for Graph-Braids,
in progress

Part 1: How a Knot Theorist Would Invent Quandles

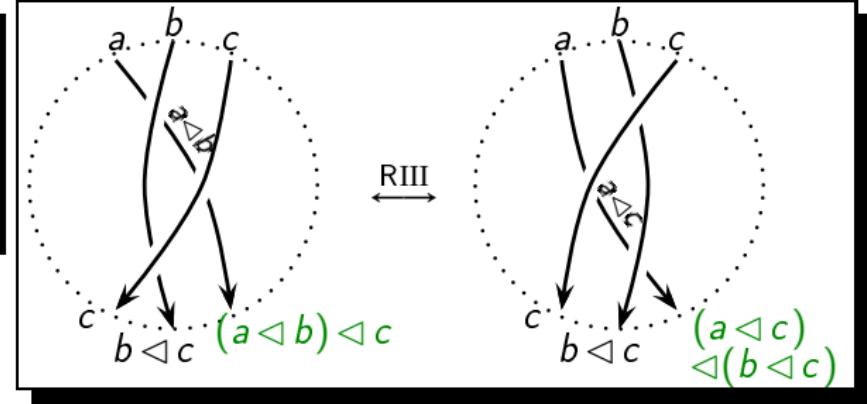
Quandle colorings of knots



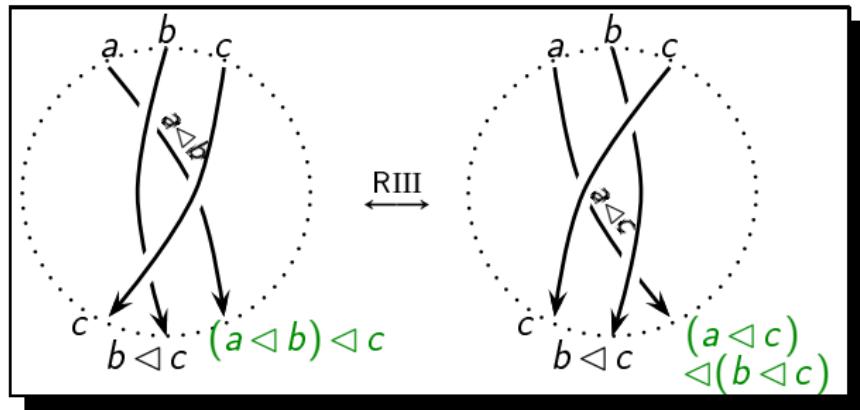
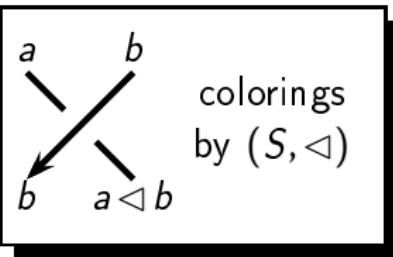
Quandle colorings of knots



colorings
by (S, \triangleleft)

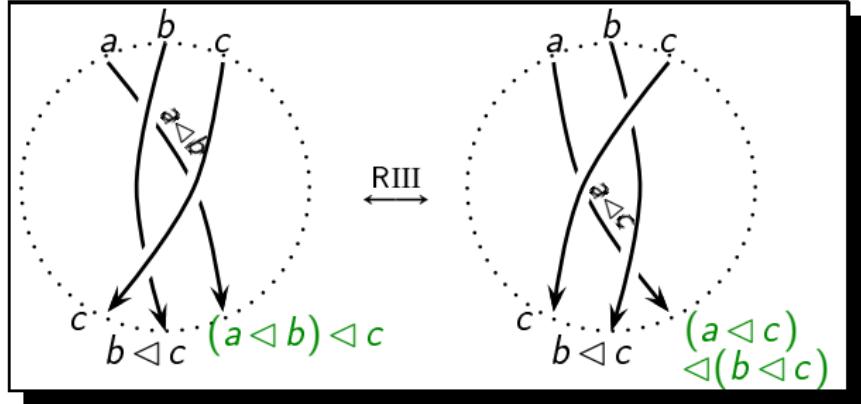
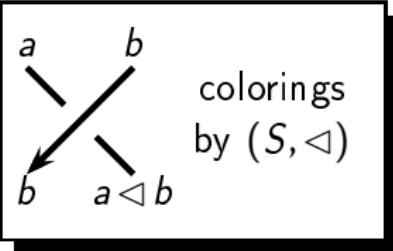


Quandle colorings of knots



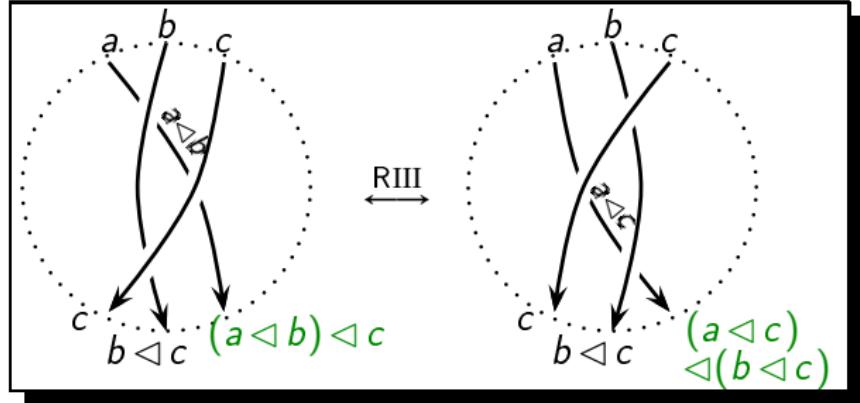
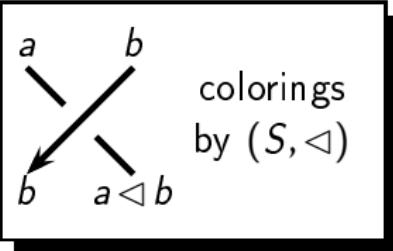
$$\text{RIII} \leftrightarrow (a \triangleleft b) \triangleleft c = (a \triangleleft c) \triangleleft (b \triangleleft c) \quad (\text{SD})$$

Quandle colorings of knots



RIII	\leftrightarrow	$(a \triangleleft b) \triangleleft c = (a \triangleleft c) \triangleleft (b \triangleleft c)$	(SD)
RII	\leftrightarrow	$a \mapsto a \triangleleft b$ is invertible	(Inv)
RI	\leftrightarrow	$a \triangleleft a = a$	(Idem)

Quandle colorings of knots



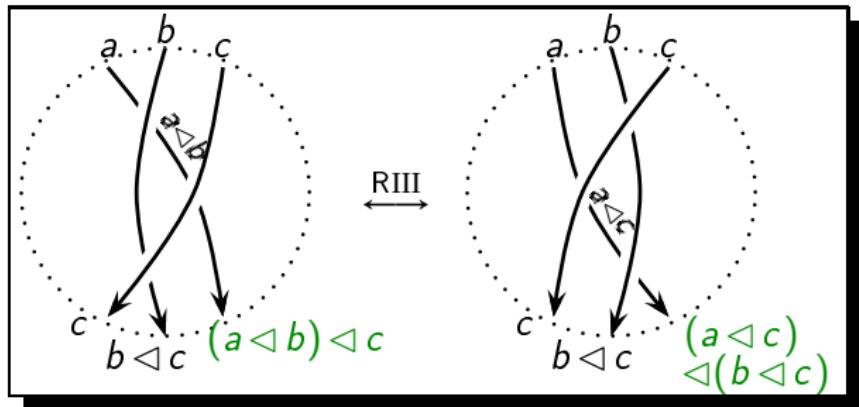
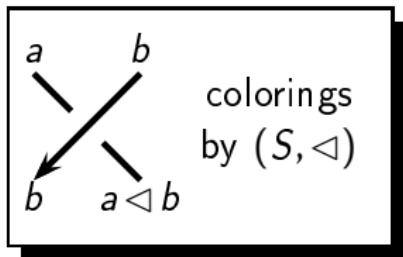
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Quandle
 } (1982 D. Joyce,
 S. Matveev)

Example

Group $G \rightsquigarrow$
 $Conj(G) = (G, g \triangleleft h = h^{-1}gh).$

Quandle colorings of knots



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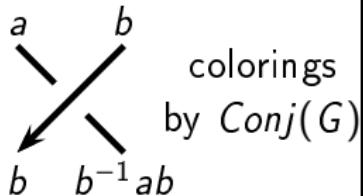
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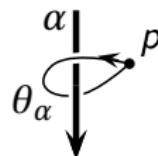
Group $G \rightsquigarrow$
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knot invariants $\overset{\text{colorings}}{\leadsto}$ quandle

Quandle colorings of knots



Wirtinger presentation:



colorings by $\text{Conj}(G)$
 \downarrow
 $\text{Rep}(\pi_1(\mathbb{R}^3 \setminus K), G)$

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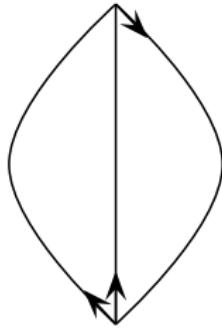
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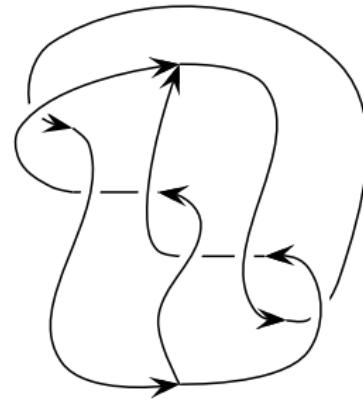
Group $G \rightsquigarrow$
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Knotted 3-valent graphs

Standard and Kinoshita-Terasaka Θ -curves:



Θ_{st}



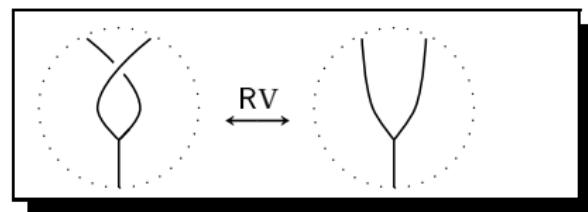
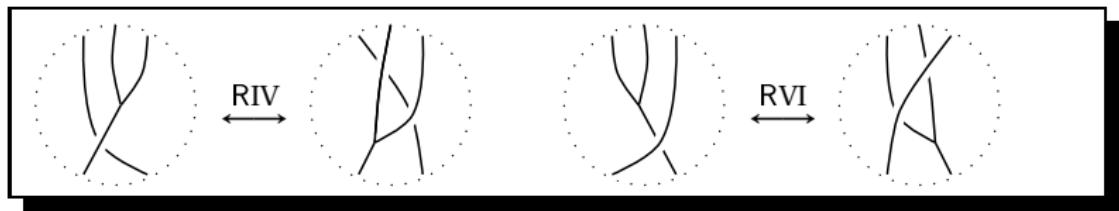
Θ_{KT}

Applications:

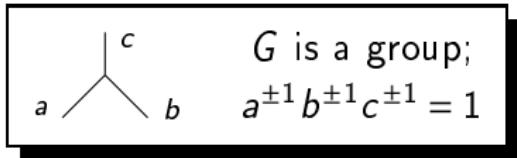
- ❖ handlebody-knots;
- ❖ foams (categorification, 3-manifolds).

Reidemeister moves for 3-graphs

1989 L.H. Kauffman, S. Yamada, D.N. Yetter:



Extending quandle colorings to 3-graphs?

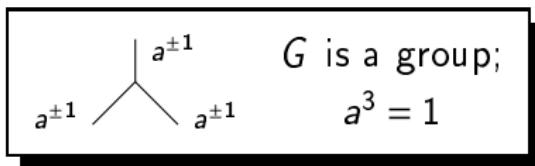
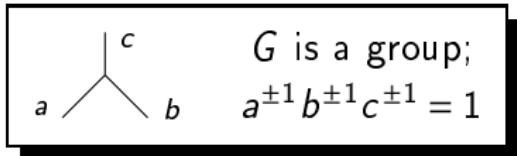


generalizations:

G -family of quandles (2012
Ishii-Iwakiri-Jang-Oshiro),

multiple conjugation quandle
(2013 A. Ishii)

Extending quandle colorings to 3-graphs?



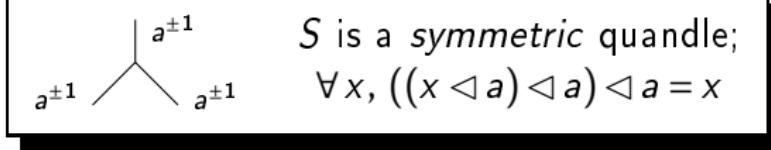
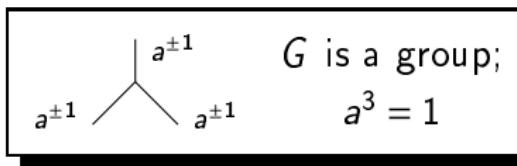
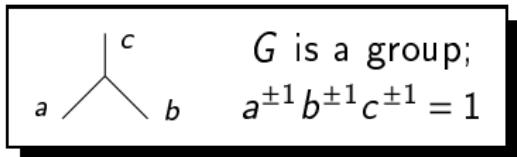
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a *vertex constant version*
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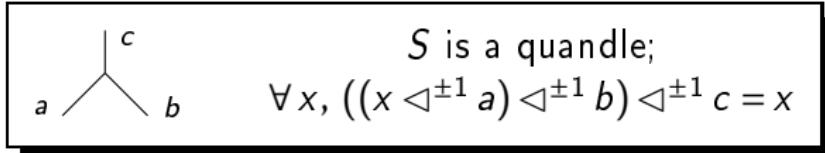
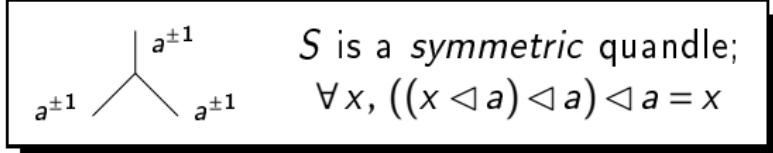
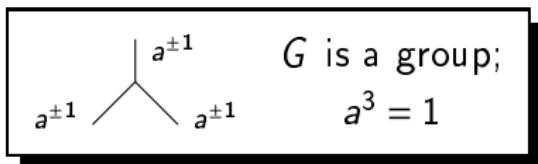
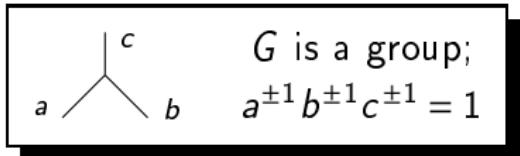
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(2007
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Extending quandle colorings to 3-graphs?



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(2010
M. Niebrzydowski)

Orientation

Well-oriented 3-graphs: only zip



and unzip



vertices.

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Well-oriented 3-graphs: only zip  and unzip  vertices.

Proposition: Every 3-graph admits a well-orientation.

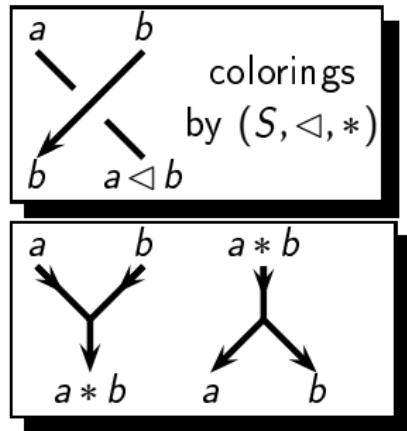
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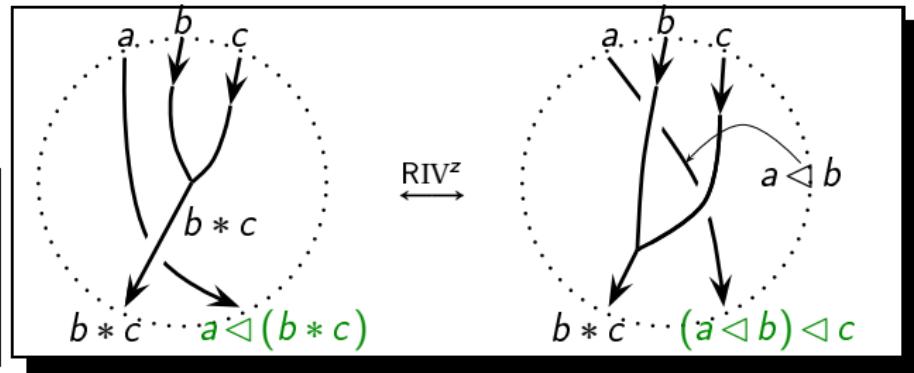
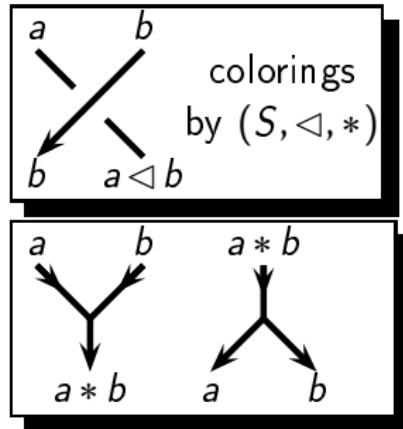
Proposition: Every 3-graph admits a well-orientation.

Unoriented 3-graph $\longrightarrow \{ \text{its well-orientations} \}$.

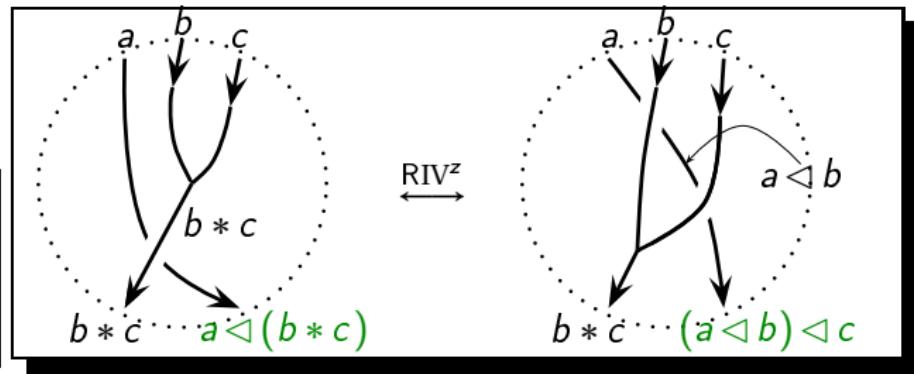
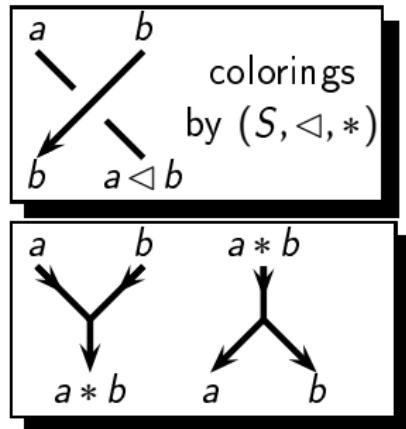
Qualgebra colorings for 3-graphs



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Qualgebra colorings for 3-graphs

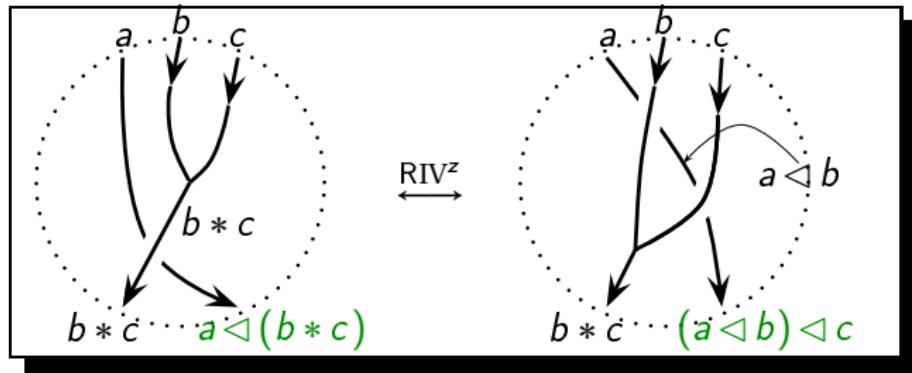
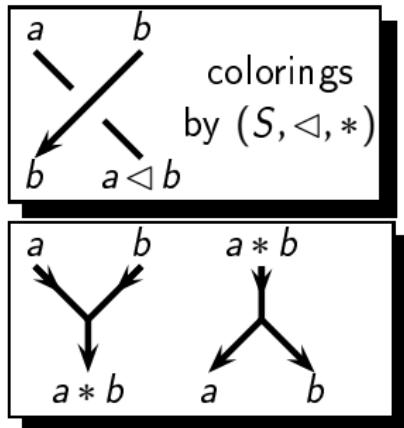


$$\text{RIV} \leftrightarrow a \triangleleft (b * c) = (a \triangleleft b) \triangleleft c$$

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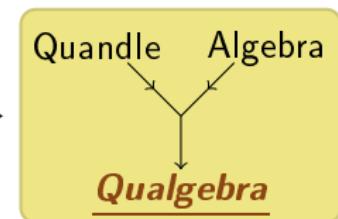
$$\text{RV} \leftrightarrow a * b = b * (a \triangleleft b)$$

Qualgebra colorings for 3-graphs

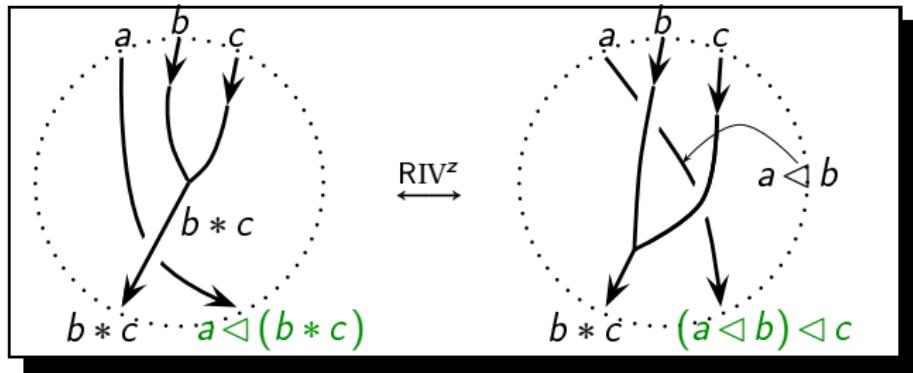
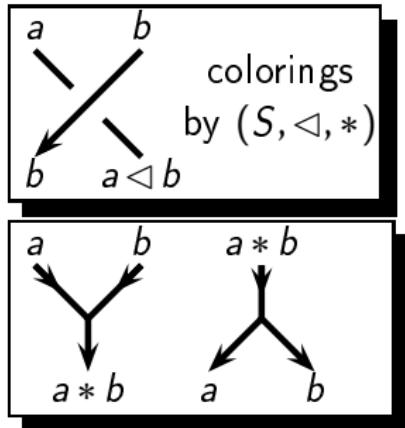


RIV	\leftrightarrow	$a \triangleleft (b * c) = (a \triangleleft b) \triangleleft c$
RVI	\leftrightarrow	$(a * b) \triangleleft c = (a * c) \triangleleft (b * c)$
RV	\leftrightarrow	$a * b = b * (a \triangleleft b)$

& (SD),
 (Inv),
 (Idem)

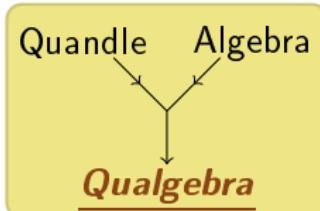


Qualgebra colorings for 3-graphs



$$\begin{array}{l} \text{RIV} \leftrightarrow a \triangleleft (b * c) = (a \triangleleft b) \triangleleft c \\ \text{RVI} \leftrightarrow (a * b) \triangleleft c = (a * c) \triangleleft (b * c) \\ \text{RV} \leftrightarrow a * b = b * (a \triangleleft b) \end{array}$$

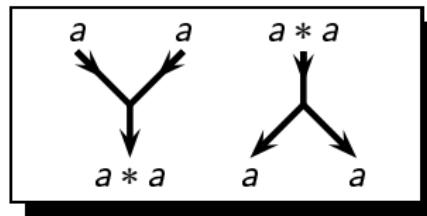
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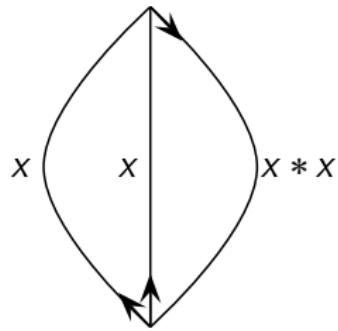
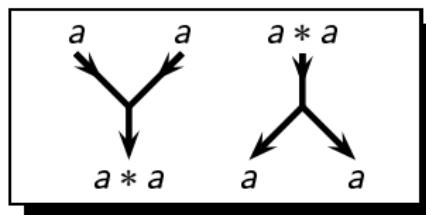
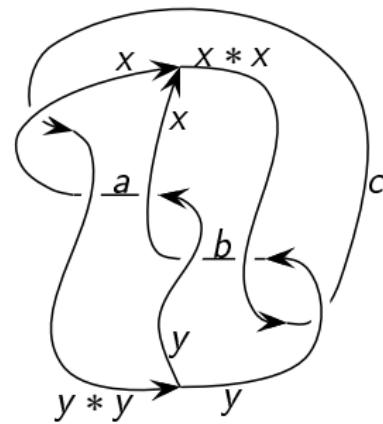
3-graph invariants $\stackrel{\text{colorings}}{\leadsto}$ qualgebra

Computation example

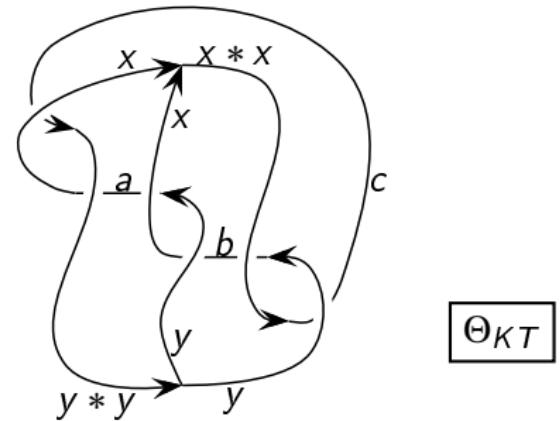
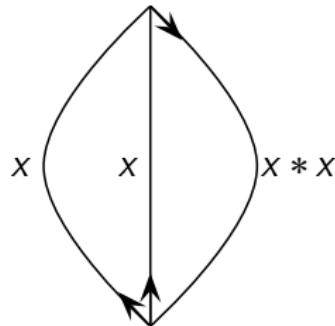
Isosceles colorings:



Computation example

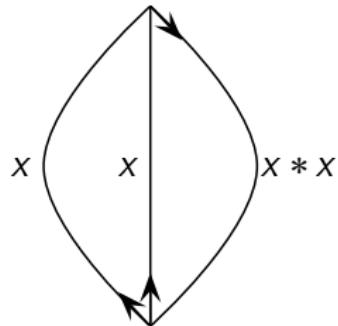
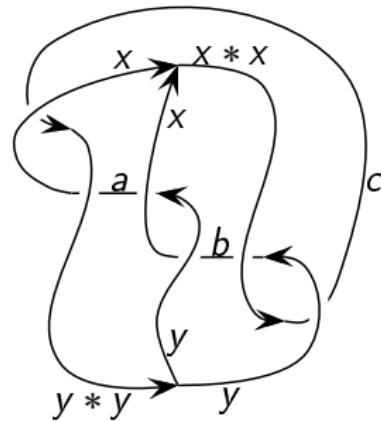
Isosceles colorings: Θ_{st}  Θ_{KT}

Computation example



$$(\star) \left\{ \begin{array}{lcl} a & = & x \triangleleft (y * y) = y \triangleleft x, \\ b & = & x \tilde{\triangleleft} y = y \tilde{\triangleleft} (x * x), \\ c & = & (y * y) \triangleleft x = (x * x) \tilde{\triangleleft} y. \end{array} \right.$$

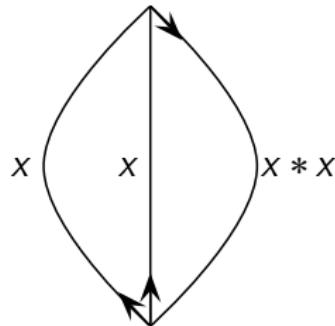
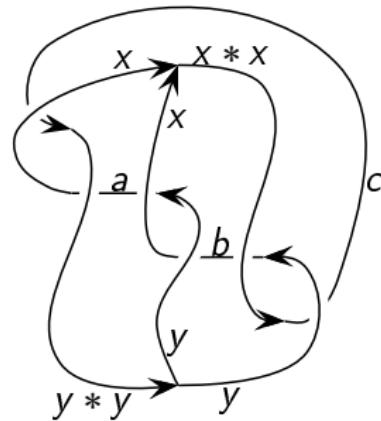
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Qualgebra $(S_4, g \triangleleft h = h^{-1}gh, g * h = gh) \rightsquigarrow (\star) \Leftrightarrow xyx = yxy.$

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 Θ_{st}  Θ_{KT}

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Qualgebra $(S_4, g \triangleleft h = h^{-1}gh, g * h = gh)$ $\leadsto (\star) \Leftrightarrow xyx = yxy$.

Solutions: $x = y$ and $x = (123), y = (432)$ and...

So, $\#\mathcal{C}_{S_4}^{iso}(\Theta_{KT}) > \#S_4 = \#\mathcal{C}_{S_4}^{iso}(\Theta_{st})$.

Part 2:

How an Algebraist
Would Invent Qualgebras

Group qualgebras

Example 1

Group $G \rightsquigarrow \underline{\text{group qualgebra}}$ $QA(G) = (G, g \triangleleft h = h^{-1}gh, g * h = gh)$.

Group qualgebras

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$$QA(G)\text{-colorings} \xleftarrow[\text{presentation}]{\text{Wirtinger}} Rep(\pi_1(\mathbb{R}^3 \setminus \Gamma), G)$$

Group quandles

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Group $G \rightsquigarrow \text{group qualgebra } QA(G) = (G, g \triangleleft h = h^{-1}gh, g * h = gh).$

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abstract level	quandle axioms	specific qualgebra axioms
topology	moves RI-RIII	moves RIV-RVI
groups	conjugation	conjugation-multiplication interaction

Group qualgebras

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abstract level	quandle axioms	specific qualgebra axioms
topology	moves RI-RIII	moves RIV-RVI
groups	conjugation	conjugation-multiplication interaction

Quandle axioms \Rightarrow all properties of conjugation.

⚠ Qualgebra axioms $\not\Rightarrow$ all properties of conjugation/multiplication interaction.

$$(b \triangleleft a) * (a \triangleleft b) = ((a \widetilde{\triangleleft} b) \triangleleft a) * b \\ (= a^{-1}bab^{-1}ab)$$

Other qualgebra examples

Example 1

Group $G \rightsquigarrow$ group qualgebra $QA(G) = (G, g \triangleleft h = h^{-1}gh, g * h = gh)$.

Example 1'

Group G & $X \subset G \rightsquigarrow$ the sub-qualgebra of $QA(G)$ generated by X .

Other qualgebra examples

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Example 0

Trivial qualgebra $(S, a \triangleleft b = a, a * b)$, where $*$ is commutative.

Other qualgebra examples

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Group $G \rightsquigarrow$ group qualgebra $QA(G) = (G, g \triangleleft h = h^{-1}gh, g * h = gh)$.

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Group G & $X \subset G \rightsquigarrow$ the sub-qualgebra of $QA(G)$ generated by X .

Example 0

Trivial qualgebra $(S, a \triangleleft b = a, a * b)$, where $*$ is commutative.

\rightsquigarrow abstract graph invariants

“Qualgebrability”

Existence

Dihedral quandle $(\mathbb{Z}/n\mathbb{Z}, a \triangleleft b = 2b - a)$ is not “qualgebrizable”:

$$\begin{aligned}(a \triangleleft b) \triangleleft c &= 2c - 2b + a \\ a \triangleleft (b * c) &= 2(b * c) - a\end{aligned}$$

“Qualgebrability”

Existence

Dihedral quandle $(\mathbb{Z}/n\mathbb{Z}, a \triangleleft b = 2b - a)$ is not “qualgebrizable”:

$$\begin{aligned}(a \triangleleft b) \triangleleft c &= 2c - 2b + a \\ a \triangleleft (b * c) &= 2(b * c) - a\end{aligned}$$

Uniqueness

$QA(S_3)$ is the unique “qualgebraization” of $\text{Conj}(S_3)$.

Small examples

Example 4

Non-trivial qualgebra structures on $Q = \{p, q, r, s\}$?

Description: put $\overline{p} = q, \overline{q} = p, \overline{r} = r, \overline{s} = s$;
 $x \triangleleft r = \overline{x}, \quad x \triangleleft y = x$ for other y ;
* is commutative,
 $r * x = r$ for $x \neq r, \quad r * r = s * s = s,$
 $p * q = s, \quad p * p = \overline{q * q} \in \{p, q, s\},$
 $p * s = \overline{q * s} \in \{p, q, s\}.$

Small examples

Example 4

Non-trivial qualgebra structures on $Q = \{p, q, r, s\}$?

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* is commutative,

$\leadsto 9$ structures.

$r * x = r$ for $x \neq r, \quad r * r = s * s = s$,

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Small examples

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$p * s = \overline{q * s} \in \{p, q, s\}$.

✿ Not cancellative \Rightarrow do not come from groups.

Small examples

Example 4

Non-trivial qualgebra structures on $Q = \{p, q, r, s\}$?

Description: put $\overline{p} = q, \overline{q} = p, \overline{r} = r, \overline{s} = s$;

$x \triangleleft r = \overline{x}, \quad x \triangleleft y = x$ for other y ;

* is commutative,

$\leadsto 9$ structures.

$r * x = r$ for $x \neq r, \quad r * r = s * s = s$,

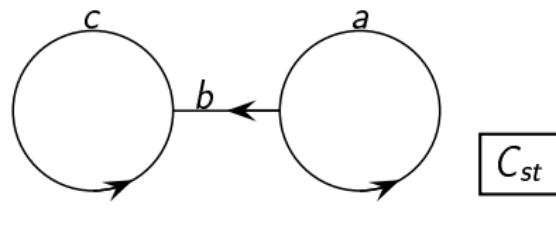
$p * q = s, \quad p * p = \overline{q * q} \in \{p, q, s\}$,

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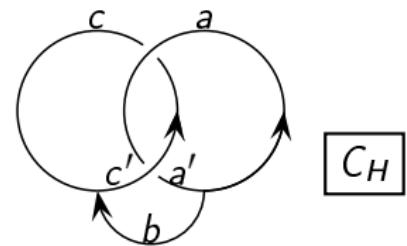
- ❖ Not cancellative \Rightarrow do not come from groups.
- ❖ Two are associative, three have neutral elements, none are unital associative.

Computation example

Standard and Hopf cuff graphs:



C_{st}



C_H

$$\#\mathcal{C}_Q(C_{st}) = \#\{(a, b, c) \in Q \mid b * a = a, b * c = c\} = 18,$$

$$\#\mathcal{C}_Q(C_H) = \#\{(a, b, c) \in Q \mid b * a = a \triangleleft c, b * c = c \triangleleft a\} = 14.$$

Part 3:
Variations on
Qualgebra Ideas

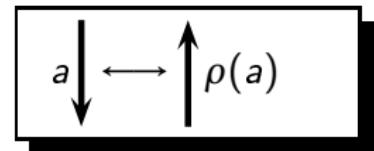
Symmetric qualgebras and orientation independence

Problem: Qualgebra invariants depend on orientations.

Symmetric qualgebras and orientation independence

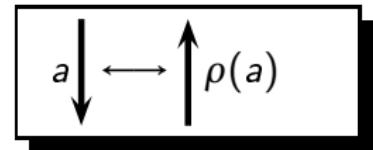
Problem: Qualgebra invariants depend on orientations.

Solution: a *good involution* $\rho : S \rightarrow S$.



Symmetric quandles and orientation independence

Problem: Quandle invariants depend on orientations.



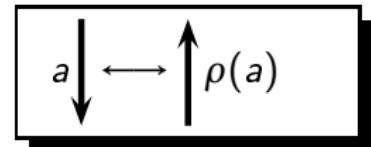
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$$\begin{aligned}\rho(\rho(a)) &= a \\ \rho(a) \triangleleft b &= \rho(a \triangleleft b) \\ a \triangleleft \rho(b) &= a \tilde{\triangleleft} b \\ (a * b) * \rho(b) &= \rho(b) * (b * a) = a\end{aligned}$$

$$\left. \begin{array}{c} \text{Symmetric} \\ \text{quandle} \\ (1996) \\ \text{S. Kamada} \end{array} \right\} \quad \left. \begin{array}{c} \text{Symmetric} \\ \text{qualgebra} \end{array} \right\}$$

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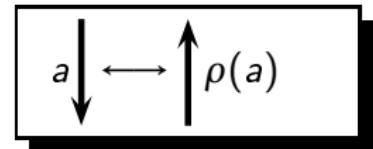
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Example

Group $G \rightsquigarrow QA^*(G) = (G, g \triangleleft h = h^{-1}gh, g * h = h, \rho(h) = h^{-1})$.

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abstract level	specific symmetric qualgebra axioms
topology	unoriented 3-graphs
groups	conjugation- and multiplication-inversion interaction

Symmetric qualgebras: examples

$$\rho(\rho(a)) = a$$

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Properties

- ✿ $\forall b$, maps $a \mapsto a * b$ and $a \mapsto b * a$ are bijections. \leadsto pseudo-sudoku
- ✿ Symmetric qualgebras with associative $*$ \longleftrightarrow groups.

Symmetric qualgebras: examples

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Example 3

$S = \{x, y, z\}$, $a \triangleleft b = a$, $*$ is commutative.

*	x	y	z
x	x	y	z
y	y	z	x
z	z	x	y
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Symmetric qualgebras: examples

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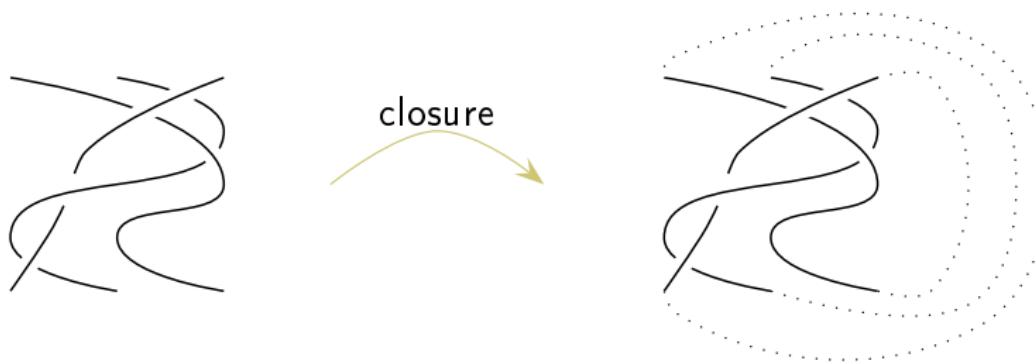
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Example 4

- ❖ Non-trivial qualgebras of size 4 are not symmetric.
- ❖ Trivial qualgebras: $QA^*(\mathbb{Z}/4\mathbb{Z})$, $QA^*(\mathbb{Z}/2\mathbb{Z} \times \mathbb{Z}/2\mathbb{Z})$, and two non-associative ones.

Alexander-Markov theorem



Theorem (Alexander, 1923; Markov, 1935)

- ❖ Surjectivity.
- ❖ Kernel: moves M1 and M2.

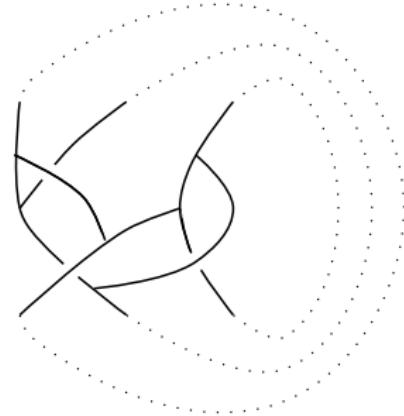


Branched Alexander-Markov theorem

Branched
braids



closure
→

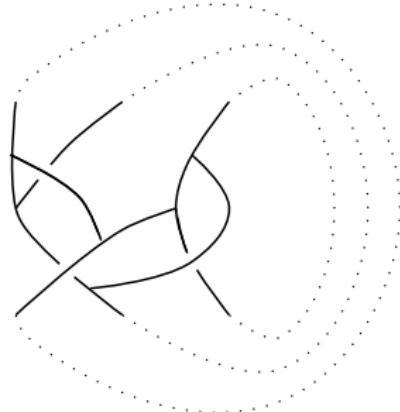


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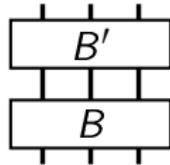


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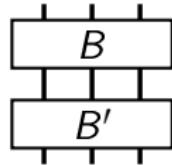


Theorem (S. Kamada - L., 2014)

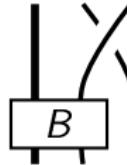
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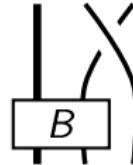
M1



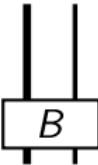
M2



M2



M2



Branched Alexander-Markov theorem

branched braids $\xrightarrow{\text{closure}}$ 3-graphs

Theorem (S. Kamada - L., 2014)

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- ❖ Kernel: moves M1 and M2.

Generalizations

- ❖ Graph-braids (vertices of arbitrary valence).
- ❖ Virtual and welded versions.

Branched Alexander-Markov theorem

branched braids $\xrightarrow{\text{closure}}$ 3-graphs

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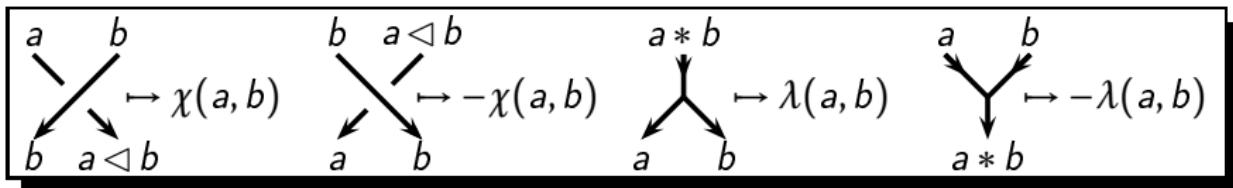
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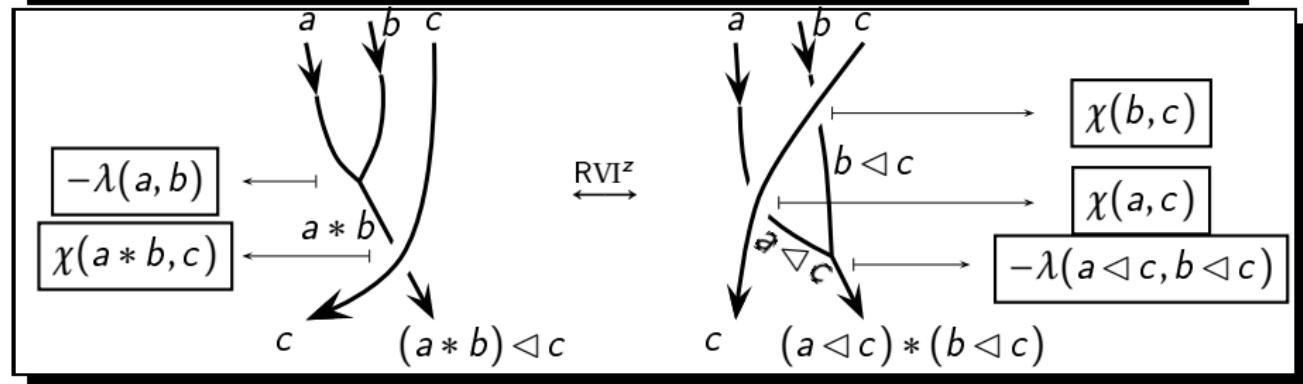
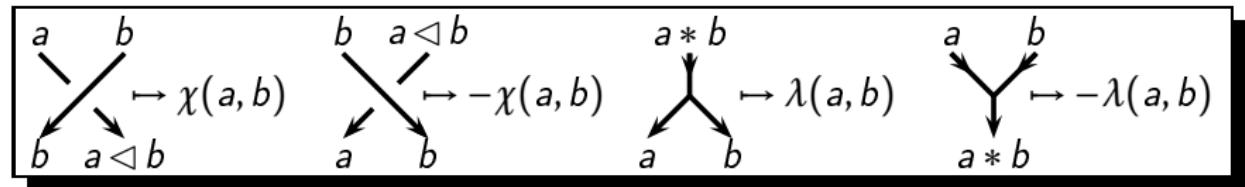
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branched braid invariants $\stackrel{\text{colorings}}{\leadsto}$ weak quandle
 $(\text{omit } a \triangleleft a = a)$

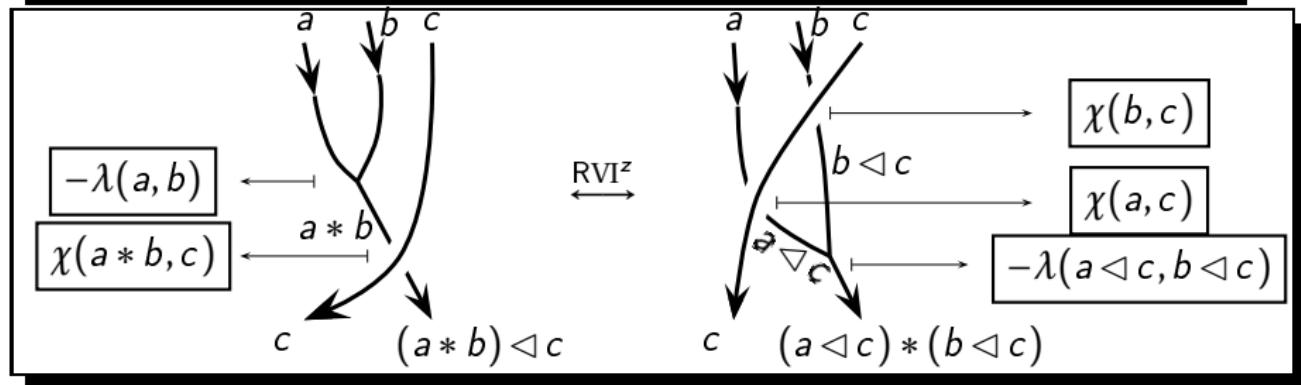
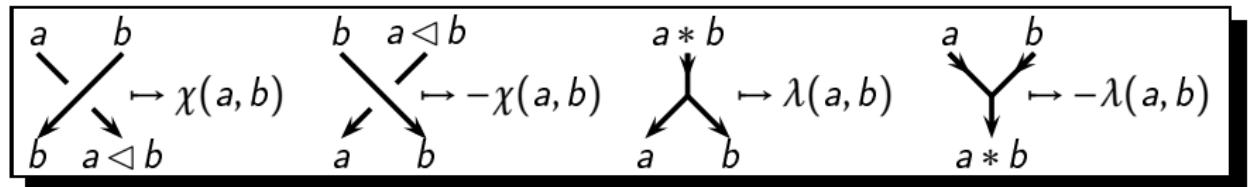
Qualgebra cocycle invariants for 3-graphs



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Qualgebra cocycle invariants for 3-graphs

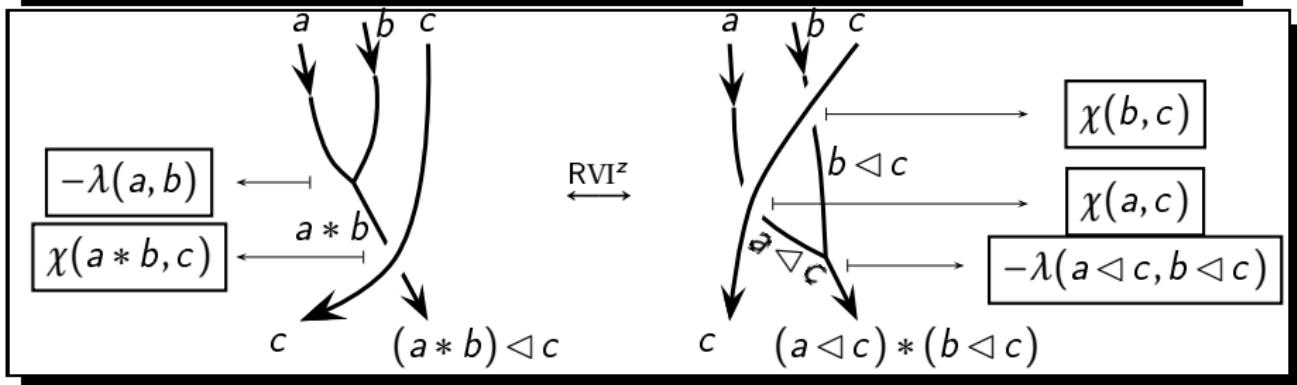
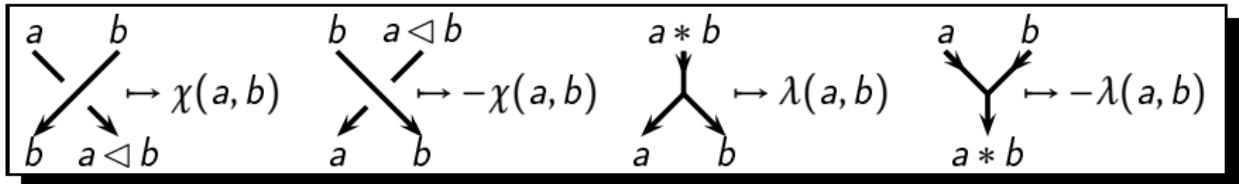


$$\text{RIV} \leftrightarrow \chi(a, b * c) = \chi(a, b) + \chi(a \triangleleft b, c)$$

$$\text{RVI} \leftrightarrow \chi(a * b, c) + \lambda(a \triangleleft c, b \triangleleft c) = \chi(a, c) + \chi(b, c) + \lambda(a, b)$$

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Qualgebra cocycle invariants for 3-graphs



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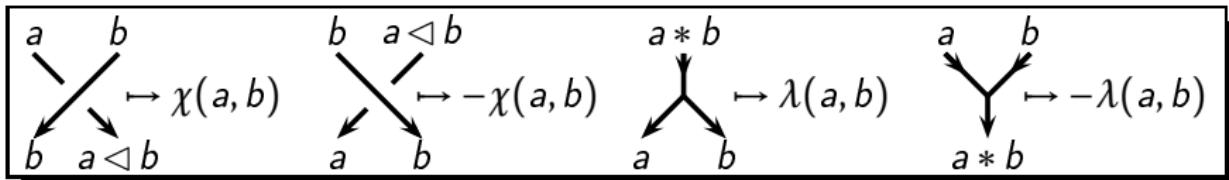
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**Qualgebra
2-cocycle**

RI-RIII are automatic.

Qualgebra cocycle invariants for 3-graphs



$$\left. \begin{array}{l} \text{RIV} \leftrightarrow \chi(a, b * c) = \chi(a, b) + \chi(a \triangleleft b, c) \\ \text{RVI} \leftrightarrow \chi(a * b, c) + \lambda(a \triangleleft c, b \triangleleft c) = \chi(a, c) + \chi(b, c) + \lambda(a, b) \\ \text{RV} \leftrightarrow \chi(a, b) + \lambda(a, b) = \lambda(b, a \triangleleft b) \end{array} \right\} \begin{array}{l} \text{Qualgebra} \\ \text{2-cocycle} \end{array}$$

RI-RIII are automatic.

3-graph invariants $\stackrel{\text{colorings}}{\sim} \stackrel{\text{weights}}{\sim}$ qualgebra & 2- or 3-cocycle

Qualgebra cocycles: example

Example 4

$$Q = \{p, q, r, s\}$$

$$x \triangleleft r = \overline{x}, \quad x \triangleleft y = x \text{ for other } y;$$

* is commutative,

$$r * x = r \text{ for } x \neq r, \quad r * r = s * s = s,$$

$$p * p = \overline{q * q} \in \{p, q, s\}, \quad p * q = s,$$

$$p * s = \overline{q * s} \in \{p, q, s\}.$$

- ✿ $Z^2(Q) \cong \mathbb{Z}^8$
- ✿ $B^2(Q) \cong \mathbb{Z}^4$
- ✿ $H^2(Q) \cong \mathbb{Z}/2\mathbb{Z} \oplus \mathbb{Z}^4$

Qualgebra cocycle invariants: diverse questions

Qualgebra 2-coboundaries:

$$\phi : S \rightarrow \mathbb{Z} \rightsquigarrow \quad \chi(a, b) = \phi(a \triangleleft b) - \phi(a), \quad \rightsquigarrow \text{trivial}$$

$$\lambda(a, b) = \phi(a) + \phi(b) - \phi(a * b) \quad \text{invariants}$$

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Question: Define and study qualgebra homology in higher degrees?

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- ✿ Region coloring and shadow cocycle invariants.

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Enhancements

- ✿ Region coloring and shadow cocycle invariants.
- ✿ Distinguish zip- and unzip-vertices:

