## Long Cycles through Specified Vertices in a Graph

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## ABSTRACT

In this paper, we consider the length of the longest cycle through specified vertices. We show the following two results. (1) Let G be a k-connected graph of order at least 2k and circumference l. Suppose m < k. Then for any m vertices of G, G has a cycle which contains all of them and has length at least  $\frac{k-m}{k}l + 2m$ . (2) Let G be a 3-connected planar graph with circumference l. Then for any three vertices of G, there exists a cycle which contains all of them and has length at least  $\frac{1}{4}l + 3$ .

Here, we consider finite simple graphs. Let G be a graph. By Dirac's theorem[3] G has a cycle through specified k vertices. In [2] Dirac also showed that a 2-connected graph of order n and minimum degree at least d has a cycle of length at least min $\{n, 2d\}$ . Locke[4] and Voss[7] generalized his result by showing that under the same conditions the graph has a cycle of length at least min $\{n, 2d\}$  which contains specified two vertices.

These results lead us to the following question: Does a k-connected graph have a long cycle through specified m vertices  $(m \le k)$ ? In this paper we investigate this question.

For basic graph-theoretic terminology, we refer the reader to [1]. Let G be a graph. The circumference of G, denoted by cir(G), is the length of the longest cycle of G. We denote by w(G) the number of components of G. For  $k \ge 0$  and  $S \subset V(G)$ , we call S a k-cutset if  $w(G-S) \ge 2$  and |S| = k. We often identify a subgraph H of G with its vertex set V(H). Especially, when x is a vertex of H, we write  $x \in H$  instead of  $x \in V(H)$ . Furthermore, we write |H| instead of |V(H)|. When we consider a cycle, we always give it an orientation. Let  $C^+$  be the orientation of a cycle C and  $C^-$  be its reverse orientation. Let  $C^+ = x_0, x_1, \ldots, x_{n-1}, x_n$  be a cycle. For  $x_i, x_j \in C$ , we define a subpaths  $C^+[x_i, x_j]$  and  $C^-[x_i, x_j]$  of C by

$$C^+[x_i, x_j] = x_i, x_{i+1}, \ldots, x_{j-1}, x_j,$$

$$C^{-}[x_i, x_j] = x_i, x_{i-1}, \ldots x_{j+1}, x_j.$$

We also define  $C^+(x_i, x_j)$  and  $C^-(x_i, x_j)$  by

$$C^{+}(x_{i}, x_{j}) = C^{+}[x_{i}, x_{j}] - \{x_{i}, x_{j}\},$$

and

$$C^{-}(x_i, x_j) = C^{-}[x_i, x_j] - \{x_i, x_j\}.$$

Furthermore,  $C^+[x_i, x_j] = C^+[x_i, x_j] - \{x_j\}$ . Subpaths  $C^-[x_i, x_j)$ ,  $C^+(x_i, x_j]$ ,  $C^-(x_i, x_j]$ are defined similarly. Let  $x_1, x_2, \ldots, x_s$  be a path. We denote by end(P) the set of endvertices of P;  $end(P) = \{x_1, x_s\}$ . Let  $P = x_1, x_2, \ldots, x_s$  and  $Q = y_1, y_2, \ldots, y_t$  be paths such that  $x_s = y_1$ . We denote by  $P \cdot Q$  the walk  $x_1, x_2, \ldots, x_s = y_1, y_2, \ldots, y_t$ .

Let  $z \in V(G)$  and  $S \subset V(G) - \{z\}$ . A subgraph F of G is called a (z, S)-fan if F has the following decomposition  $F = \bigcup_{i=1}^{k} P_i$ , where

(1) each  $P_i$  is a path between z and  $a_i \in S$ , and

(2)  $P_i \cap S = \{a_i\}$ , and  $P_i \cap P_j = \{z\}$  if  $i \neq j$ .

We call k the size of the fan F. The vertices  $a_1, \ldots, a_k$  are called endvertices of F and the set of its endvertices is denoted by end(F). Since F is a tree, for any two vertices  $x, y \in F$  the path in F which joins x and y is unique. We denote this path by F[x,y]. We define F[x,y] by  $F[x,y) = F[x,y] - \{y\}$ . Paths F(x,y] and F(x,y) are defined similarly.

The following theorem is well-known, called the generalized Menger's theorem.

**THEOREM** A ([1, Theorem 6.7]). Let G be a k-connected graph,  $z \in V(G)$ , and  $S \subset V(G) - \{z\}$ . Then G has a (z, S)-fan of size min $\{|S|, k\}$ .

The following theorem was proved by Perfect[5].

**THEOREM** B (Perfect[5]). Let G be a graph,  $z \in V(G)$ , and  $S \subset V(G) - \{z\}$ . Suppose G has two (z, S)-fans  $F_1$  and  $F_2$  of size  $k_1$  and  $k_2$ , respectively. If  $k_1 \leq k_2$ , then G has a (z, S)-fan F' of size  $k_2$  such that  $end(F_1) \subset end(F')$ .

We use these two theorems in the proofs our results.

First, we show that the existence of long cycles through specified m vertices in a k-connected graph is assured if m < k. Note that a k-connected graph is hamiltonian if its order is at most 2k, by Dirac's theorem.

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and

**THEOREM 1.** Let  $k \ge 2$ ,  $0 \le m \le k$  and G be a k-connected graph of order at least 2k. For any m vertices  $x_1, \ldots, x_m$  of G, there exists a cycle such that

- (1)  $x_1, ..., x_m \in V(C)$ , and
- (2)  $|C| \ge \frac{k-m}{k}\operatorname{cir}(G) + 2m$ .

Recently, Seymour and Truemper sent me a proof which is simpler than the original one. We show their proof.

**Proof** (due to Seymour and Truemper). The proof is by induction on m. For m = 1, let  $x \in V(G)$ , and let C be a longest cycle in G. Since  $|C| \ge 2k$ ,

$$\frac{k-1}{k}\operatorname{cir}(G) + 2 = |C| - \frac{|C|}{k} + 2 \le |C|.$$

So we may assume  $x \notin V(C)$ . Now G has an (x, C)-fan of size k. The endvertices of F divide C into k paths, and any shortest one P of these paths, say  $P = C^+[u, v]$  has length at most  $\frac{1}{k}\operatorname{cir}(G)$ . So  $C^+[v, u] \cdot F[u, v]$  is a cycle which contains x and has length at least

$$|C| - \frac{\operatorname{cir}(G)}{k} + 2 = \frac{k-1}{k}\operatorname{cir}(G) + 2$$

as desired.

Suppose m > 1, and let C be a longest cycle containing at least m - 1 members of S. By the induction hypothesis,

$$|C| \ge \frac{k-m+1}{k} \operatorname{cir}(G) + 2(m-1)$$
  
=  $\frac{k-m}{k} \operatorname{cir}(G) + 2m + \frac{\operatorname{cir}(G)}{k} - 2$   
 $\ge \frac{k-m}{k} \operatorname{cir}(G) + 2m.$  (\*)

So we may assume that exactly one member x of S does not lie on C. Since  $cir(G) \ge 2k$ ,  $|C| \ge 2k$ . So G has an (x, C)-fan of size k. The endvertices of F divide C into k paths. We call such a path bad if it contains some member of S internally, and we call it good if it is not bad. Let b represent the number of bad paths, and let L be the sum of lengths of the bad paths. Then some good path  $P = C^+[u, v]$  has length at most

$$\frac{|C| - L}{k - b}$$

(, where  $|C| \ge 2k$  and  $k \ge m-1$ ). Keeping |C| and k fixed, and under the conditions  $L \ge 2b$  and  $b \le m-1$ , this is maximized when L = 2b and b = m-1. Hence,

$$|P| \leq \frac{|C| - 2(m-1)}{k - m + 1}.$$

A cycle  $C^+[v, u] \cdot F[u, v]$  contains S, and from (\*) it has length at least

$$|C| - \frac{|C| - 2(m-1)}{k - m + 1} + 2 \ge \frac{k - m}{k} \operatorname{cir}(G) + 2m$$

as desired.

Theorem 1 is sharp. Let,  $k \ge 2$ ,  $s \ge 1$ , and  $0 \le m \le k$ . Let  $H_0, H_1, \ldots, H_k$  and  $H'_0$  be graphs such that  $H_1 \simeq \cdots \simeq H_k \simeq K_s$ ,  $H_0 \simeq \overline{K_m}$  and  $H'_0 \simeq \overline{K_k}$ . Suppose vertex sets  $V(H_0), \ldots, V(H_k)$  and  $V(H'_0)$  are disjoint. Define G(k, m, s) by  $G(k, m, s) = (H_1 \cup \cdots \cup H_k \cup H_0) + H'_0$ . Then G(k, m, s) is k-connected,  $|G(k, m, s)| = ks + k + m \ge 2k$ , and  $\operatorname{cir}(G(k, m, s)) = ks + k$ . On the other hand, the length of the longest cycle through  $V(H_0)$  is (k - m)s + k + m. The above example shows that large circumference does not assure the existence of long cycles through specified k vertices in k-connected graphs.

Next, we confine ourselves to planar graphs. Even if we consider only planar graphs, the length of the longest cycle through specified two vertices in a 2-connected graph is independent of its circumference. Let  $C = x_0, x_1, \ldots, x_m = x_0$  be a cycle of length m  $(m \ge 4)$ . Add a new vertex y and join  $yx_1$  and  $yx_{m-1}$ . Then this graph has circumference m, but the unique cycle through y and  $x_0$  has length four. On the other hand, by Tutte's theorem[6] 4-connected planar graphs are hamiltonian, and hence the length of the longest cycle through four specified vertices in a 4-connected planar graph is equal to its circumference. On a planar graph of connectivity three, we show the following theorem.

**THEOREM 2.** Let G be a 3-connected planar graph. Then any three vertices of G lie on a cycle of length at least  $\frac{1}{4}\operatorname{cir}(G) + 3$ .

The proof of Theorem 2 is given by the following two lemmas.

**LEMMA 1.** Let G be a 3-connected planar graph. Then for any two vertices x, y, there exists a cycle C such that

(1)  $x, y \in V(C)$ . (2)  $|C| \ge \frac{1}{2} \operatorname{cir}(G) + 2$ .

**LEMMA 2.** Let G be a 3-connected planar graph,  $x, y, z \in V(G)$  and C be a cycle of G such that  $x, y \in V(C)$ . Then there exists a cycle C' such that (1)  $x, y, z \in V(C')$ . (2)  $|C'| \ge \frac{1}{2}|C| + 2$ . **Proof of Lemma 1.** If G is hamiltonian, then the lemma clearly holds. So we may assume that G is not hamiltonian, which implies  $|G| \ge 7$  and  $\operatorname{cir}(G) \ge 6$ . Let C be a longest cycle of G. We consider three cases.

Case 1.  $\{x, y\} \subset V(C)$ .

This case is trivial.

Case 2.  $|\{x, y\} \cap V(C)| = 1.$ 

We may assume that  $x \in V(C)$  and  $y \notin V(C)$ . Consider a (y, C)-fan F of size three. Let  $\operatorname{end}(F) = \{y_1, y_2, y_3\}$ . If  $x \in \{y_1, y_2, y_3\}$ , say  $x = y_1$ , then we have two cycles  $C^+[x, y_2] \cdot F[y_2, x]$  and  $C^-[x, y_2] \cdot F[y_2, x]$ , one of which has length at least  $\frac{1}{2}|C| + 2 = \frac{1}{2}\operatorname{cir}(G) + 2$  and contains both x and y. Next, assume  $x \notin \{y_1, y_2, y_3\}$ . We may assume  $x \in C^+(y_3, y_1)$ . Then one of the two cycles  $C^+[y_3, y_2] \cdot F[y_2, y_3]$  and  $C^-[y_1, y_2] \cdot F[y_2, y_1]$  has the desired properties.

Case 3.  $\{x, y\} \cap V(C) = \emptyset$ .

First, we show the following claims.

Claim 1. Suppose there exists a path P in G such that

(1) P joins two distinct vertices of C and P intersects C only at its endvertices. (2)  $x, y \in V(P)$ .

Then the Lemma follows.

*Proof.* Let a and b be endvertices of P. Then one of the two cycles  $P[a, b] \cdot C^+[b, a]$  and  $P[a, b] \cdot C^-[b, a]$  satisfies the desired properties.

Claim 2. Suppose there exist two paths P and Q such that

(1)  $V(P) \cap V(Q) = \emptyset$ .

(2) Both P and Q join two vertices of C.

(3)  $V(P) \cap V(C) = \operatorname{end}(P)$  and  $V(Q) \cap V(C) = \operatorname{end}(Q)$ .

(4) Vertices of end(P) and vertices of end(Q) appear alternately around  $C^+$ .

(5)  $x \in V(P)$  and  $y \in V(Q)$ .

Then the lemma follows.

*Proof.* Let  $end(P) = \{x_1, x_2\}$  and  $end(Q) = \{y_1, y_2\}$ . We may assume  $x_1, y_1, x_2$  and  $y_2$  appear in this order around  $C^+$ . Then one of the two cycles

 $C^+[x_1, y_1] \cdot Q[y_1, y_2] \cdot C^-[y_2, x_2] \cdot P[x_2, x_1]$ 

and

$$C^{-}[x_1, y_2] \cdot Q[y_2, y_1] \cdot C^{+}[y_1, x_2] \cdot P[x_2, x_1]$$

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has the desired properties.

Let  $\operatorname{end}(F_1) = \{x_1, x_2, x_3\}$ . We may assume that  $x_1, x_2, x_3$  appear in this order around  $C^+$ . If  $y \in V(F_1)$ , then the theorem follows by Claim 1. Suppose  $y \notin V(F_1)$ . Let  $D = C \cup F_1$ . Let  $F_2$  be a (y, D)-fan of size three. Let  $\operatorname{end}(F_2) = \{y_1, y_2, y_3\}$ . If  $\operatorname{end}(F_2) \cap (F_1 - \{x_1, x_2, x_3\}) \neq \emptyset$ , then the lemma follows by Claim 1. So we may assume  $\operatorname{end}(F_2) \subset V(C)$ .

Claim 3. If  $\{y_1, y_2, y_3\} \subset C^+[x_i, x_{i+1}]$  (If i = 3, we consider  $x_4 = x_1$ ), then the lemma follows.

*Proof.* We may assume  $y_1, y_2, y_3 \in C^+[x_1, x_2]$  and  $y_1, y_2$  and  $y_3$  appear in this order around  $C^+$ . Then

$$C^+[x_3, y_1] \cdot F_2[y_1, y_2] \cdot C^+[y_2, x_2] \cdot F_1[x_2, x_3]$$

or

 $C^+[x_1, y_2] \cdot F_2[y_2, y_3] \cdot C^+[y_3, x_3] \cdot F_1[x_3, x_1]$ 

has the desired properties.

By Claims 1, 2, 3, the only possible case in which the lemma would not hold is  $\{x_1, x_2, x_3\} = \{y_1, y_2, y_3\}$ . We may assume  $x_i = y_i$  (i = 1, 2, 3). Let  $D' = D \cup F_2$ . Since C is a longest cycle,  $C^+(x_1, x_2) \neq \emptyset$ . Since G is 3-connected, there exists a path P joining  $C^+(x_1, x_2)$  and  $D' - C^+[x_1, x_2]$  in  $G - \{x_1, x_2\}$ . Let  $end(P) = \{u, v\}$ ,  $u \in C^+(x_1, x_2)$  and  $v \in D' - C^+[x_1, x_2]$ . If  $v \in V(F_1) \cup V(F_2)$ , then the lemma follows by Claim 2. So we may assume  $v \in C^+(x_2, x_3]$ . Then  $F_1$ ,  $F_2$ ,  $C^+[x_1, x_2]$  and  $P[u, v] \cdot C^+[v, x_3]$  form a subdivision of  $K_{3,3}$ . This contradicts the planarity of G. Therefore, the lemma follows.

**Proof of Lemma 2.** Let  $C_0$  be a longest cycle which contains x and y. Then  $|C_0| \ge |C|$ . If G is hamiltonian, then  $C_0$  is a hamiltonian cycle, and  $|C_0| \ge 4$ . Hence the result follows. Threfore, we may assume G is not hamiltonian, and  $|G| \ge 7$ . By Lemma 1,  $|C_0| \ge \frac{1}{2} \cdot 7 + 2 \ge 5$ . So  $|C_0| \ge \frac{1}{2}|C_0| + 2 \ge \frac{1}{2}|C| + 2$ . Hence we may assume  $z \notin C_0$ . Consider a  $(z, C_0)$ -fan  $F_1$ . Let  $end(F_1) = \{z_1, z_2, z_3\}$ . We may assume that  $z_1, z_2, z_3$  appear in this order around  $C^+$ . We consider three cases.

Case 1. end $(F_1) \subset C_0^+[x, y]$  or end $(F_1) \subset C_0^+[y, x]$ .

We may assume  $\{z_1, z_2, z_3\} \subset C_0^+[x, y]$ . Then one of the two cycles  $C_0^+[z_2, z_1] \cdot F_1[z_1, z_2]$  and  $C_0^+[z_3, z_2] \cdot F_1[z_2, z_3]$  has the desired properties.

Case 2. One of end( $F_1$ ) lies on  $C_0^+(y, x)$  and the other two lie on  $C_0^+(x, y)$ .

We may assume  $z_1, z_2 \in C_0^+(x, y)$  and  $z_3 \in C_0^+(y, x)$ . Let  $C_1 = C_0^+[z_2, z_1] \cdot F_1[z_1, z_2]$ . Then  $C_0 - C_1 = C_0^+(z_1, z_2)$ . Let  $D = C_0 \cup F_1$ . By Theorem B, there exists an  $(x, D - C_0^+(z_3, z_1))$ -fan  $F_2$  of size three, such that  $z_1, z_3 \in \text{end}(F_2)$ . Let  $\text{end}(F_2) = \{z_1, z_3, a\}$ . If  $a \in F_1[z, z_1)$  or  $a \in F_1[z, z_2)$ , let

$$C_2 = C_0^+[z_1, z_3] \cdot F_1[z_3, a] \cdot F_2[a, z_1].$$

If  $a \in F_1[z, z_3)$ , let

$$C_2 = C_0^+[z_1, z_3] \cdot F_2[z_3, a] \cdot F_1[a, z_1]$$

If  $a \in C_0^+(z_2, y]$ , let

$$C_2 = C_0^+[a, z_3] \cdot F_1[z_3, z_2] \cdot C_0^-[z_2, z_1] \cdot F_2[z_1, a]$$

If  $a \in C_0^+(y, z_3)$ , let

$$C_2 = C_0^{-}[a, z_1] \cdot F_1[z_1, z_3] \cdot F_2[z_3, a].$$

Then in either case,  $C_0^+(z_1, z_2) \subset C_2$  and either  $C_1$  and  $C_2$  satisfies the desired properties. So the only remaining case is  $a \in C_0^+(z_1, z_2]$ . Let  $D' = D \cup F_2$ .

Next, consider a  $(y, D' - C_0^+(z_2, z_3))$ -fan  $F_3$  such that  $\{z_2, z_3\} \subset \text{end}(F_3)$ . Let  $\text{end}(F_3) = \{z_2, z_3, b\}$ . If  $b \in (F_1 - \text{end}(F_1)) \cup C_0^+(z_3, z_1)$ , then the lemma follows by the same argument. If  $b \in F_2(x, a) \cup F_2(x, z_1)$ , let

$$C_3 = F_3[b, z_2] \cdot C_0^-[z_2, z_1] \cdot F_1[z_1, z_3] \cdot F_2[z_3, b]$$

If  $b \in F_2(x, z_3)$ , let

$$C_{3} = F_{3}[b, z_{3}] \cdot F_{1}[z_{3}, z_{2}] \cdot C_{0}^{-}[z_{2}, z_{1}] \cdot F_{2}[z_{1}, b].$$

Then in either case  $C_0^+(z_1, z_2) \subset C_3$  and hence either  $C_1$  or  $C_3$  satisfies the desired properties. So the lemma follows unless  $b \in C_0^+[z_1, z_2)$ . (Possibly a = b.)

Now we consider the case  $a \in C_0^+(z_1, z_2)$  and  $b \in C_0^+(z_1, z_2)$ . If  $z_1, b, a, z_2$  appear in this order around  $C_0^+$ , let

$$C_4 = F_3[z_3, b] \cdot C_0^+[b, z_2] \cdot F_1[z_2, z_1] \cdot C_0^-[z_1, z_3]$$

and

$$C_5 = F_2[z_3, a] \cdot C_0^-[a, z_1] \cdot F_1[z_1, z_2] \cdot C_0^+[z_2, z_3]$$

$$C_4 = F_3[z_2, b] \cdot C_0^-[b, z_3] \cdot F_1[z_3, z_2]$$

and

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$$C_5 = F_2[z_1, a] \cdot C_0^+[a, z_3] \cdot F_1[z_3, z_1].$$

Then in either case we have  $\{x, y, z\} \subset C_4 \cap C_5$ ,  $C_0 \subset C_4 \cup C_5$ , and hence  $|C_4| \ge \frac{1}{2}|C_0| + 2$ or  $|C_5| \ge \frac{1}{2}|C_0| + 2$ . So the lemma follows.

Now, we may assume that  $a = z_2$  or  $b = z_1$ . If  $a = z_2$ , then  $F_1$ ,  $F_2$ ,  $F_3$  and  $C_0^-[b, z_1]$  form a subdivision of  $K_{3,3}$ . If  $b = z_1$ , then  $F_1$ ,  $F_2$ ,  $F_3$  and  $C_0^+[a, z_2]$  form a subdivision of  $K_{3,3}$ . Hence both contradicts the planarity of G. Therefore, the proof in this case is complete.

Case 3.  $|\{x, y\} \cap \operatorname{end}(F_1)\}| = |C_0^+(x, y) \cap \operatorname{end}(F_1)| = |C_0^+(y, x) \cap \operatorname{end}(F_1)| = 1.$ We may assume  $z_1 = x, z_2 \in C_0^+(x, y)$  and  $z_3 \in C_0^+(y, x)$ . Then either

$$C_6 = F_1[z_1, z_2] \cdot C_0^+[z_2, z_1],$$
 or  
 $C_7 = F_1[z_1, z_3] \cdot C_0^-[z_3, z_1]$ 

satisfies the desired properties.

Therefore, in each case, G has a cycle through x, y and z of length at least  $\frac{1}{2}|C_0| + 2$ .

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