数理解析研究所講究録
918 巻 1995 年 109-117

NVNF-sequentiality of Left-linear Term Rewriting Systems

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Abstract

This paper introduces an extension of NV-sequentiality defined by Oyamaguchi in [6]. We call the extension NVNF-sequentiality. It is showed that the class of NVNF-sequential systems properly includes the class of NV-sequential systems, and a indices with respect to NVNF-sequentiality are computed for a given term when a term rewriting system is NVNF-sequential.

1 introduction

Term rewriting systems can be regarded as a model for computation in which terms are reduced using a set of directed equations, called rewrite rules. Term rewriting systems play an important role in various fields of computer science such as abstract data type specifications, implementations of functional programming languages and automated deduction.

In a term rewriting system, there are possibilities that a term having normal forms has infinite reduction sequences starting with it. We require some strategies telling us which redex to contract in order to get the desired result. Therefore, it is important to have a normalizing strategy which is guaranteed to find the normal form of terms whenever their normal forms exist. Huet and Lévy [2] showed that the needed reduction strategy is normalizing for every orthogonal (i.e., left-linear and non-overlapping) term rewriting system. The needed reduction strategy always rewrites one of needed redexes which have to be rewritten in order to reach a normal form. Unfortunately, it is undecidable in general whether a redex is needed or not. However, they show that for strong sequential orthogonal term rewriting systems at least one of the needed redexes in a term not in normal form can be efficiently computed. The work of Huet and Lévy was extended to several kinds of systems. Toyama [8] extended the notion of strong sequentiality to left-linear term rewriting systems. Decidability of strong sequentiality was showed for left-linear systems by Jouannaud and Sadfi [3]. Oyamaguchi [6] introduced NV-sequentiality which is also decidable.

In this paper, we introduce an extension of NV-sequentiality. This sequentiality is called NVNF-sequentiality [5]. Like NV-sequentiality, NVNF-sequentiality is based on the analysis of left-hand sides and the non-variable parts of the right-hand side of rewrite rules. However, the reachability to the normal form is considered in NVNF-sequentiality. We first show that the class of NVNF-sequential systems properly includes the class of NV-sequential systems. Next we prove that, for a given term t, it is decidable whether an occurrence u in t is an index with respect to NVNF-sequentiality. This problem can be reduced to the problem to decide whether a term has a normal form for left-linear and right-ground term rewriting systems. When a term rewriting system is left-linear and right-ground, there exists an upper bound of the least hight of normal forms which a term t can reduce to if t has a normal form. The reachability problem for left-linear and right-ground term rewriting systems has been shown to be decidable in [1, 7]. Hence an index with respect to NVNF-sequentiality can be computed.

2 Definition

We mainly follow the notation of [2, 4]. Let $\mathcal{F}$ be a finite set of function symbols and let $\mathcal{V}$ be a countably infinite set of variables where $\mathcal{F} \cap \mathcal{V} = \phi$. The set of all terms built from $\mathcal{F}$ and $\mathcal{V}$ is denoted $\mathcal{T}(\mathcal{F}, \mathcal{V})$. The set $\mathcal{T}(\mathcal{F}, \mathcal{V})$ is sometimes denoted by $\mathcal{T}$. Terms not containing variables are called ground terms. Identity of terms is denoted by $\equiv$.

For any term t, we define the set $O(t)$ of occurrences as follow:

$$O(t) = \begin{cases} \{\varepsilon\} & \text{if } t \in \mathcal{V}, \\ \{\varepsilon\} \cup \{\, i.u \mid 1 \le i \le n,\ u \in O(t_i) \,\} & \text{if } t \equiv F(t_1, \ldots, t_n). \end{cases}$$

If $u \in O(t)$ then the subterm t/u of t at u is defined by

$$t/u \equiv \begin{cases} t & \text{if } u = \varepsilon, \\ t_i/v & \text{if } t \equiv F(t_1, \ldots, t_n) \text{ and } u = i.v. \end{cases}$$

If s is a subterm of t then we write $s \subseteq t$. If $u \in O(t)$ then the term $t[u \leftarrow s]$ obtained by replacing t/u with s is defined as follow:

$$t\,[u \leftarrow s] \equiv \begin{cases} s & \text{if } u = \varepsilon, \\ F(t_1, \ldots, t_i[v \leftarrow s], \ldots, t_n) & \text{if } t \equiv F(t_1, \ldots, t_n) \text{ and } u = i.v. \end{cases}$$

Occurrences are partially ordered by the prefix ordering $\le$, i.e. $u \le v$ if there exists w such that $u.w = v$. In this case we define v/u as w. If $u \not\le v$ and $v \not\le u$ then we say that u and v are disjoint, and write $u \perp v$. If $u_1, \ldots, u_n$ are pairwise disjoint, we use $t\,[\, u_i \leftarrow s_i \mid 1 \le i \le n \,]$ to denote $t\,[u_1 \leftarrow s_1] \cdots [u_n \leftarrow s_n]$.

A substitution θ is a mapping from $\mathcal{V}$ into $\mathcal{T}(\mathcal{F}, \mathcal{V})$. Substitutions are extended into homomorphisms from $\mathcal{T}(\mathcal{F}, \mathcal{V})$ into $\mathcal{T}(\mathcal{F}, \mathcal{V})$. In following we write $t\theta$ instead of $\theta(t)$.

A term rewriting system is a pair $(\mathcal{F}, \mathcal{R})$ consisting of a set $\mathcal{F}$ of function symbols and a finite set $\mathcal{R}$ of rewrite rules. A rewrite rule is a pair $\langle l, r \rangle$ of terms such that $l \notin \mathcal{V}$ and any variable in r also occurs in l. We write $l \to r$ for $\langle l, r \rangle$. An instance of a left-hand side of a rewrite rule is a redex. The rewrite rules of a term rewriting system $(\mathcal{F}, \mathcal{R})$ define a reduction relation $\to_{\mathcal{R}}$ on $\mathcal{T}(\mathcal{F}, \mathcal{V})$ as follow: $t \to_{\mathcal{R}} s$ if there exist a rewrite rule $l \to r \in \mathcal{R}$, a occurrence $u \in O(t)$ and a substitution θ such that $t/u \equiv l\theta$ and

$s \equiv t[u \leftarrow s]$. The transitive-reflexive closure of $\to_{\mathcal{R}}$ is denoted by $\to^*_{\mathcal{R}}$. $\to^+_{\mathcal{R}}$ denotes the transitive closure of $\to_{\mathcal{R}}$ and $\to^{\equiv}_{\mathcal{R}}$ denotes the reflexive closure of $\to_{\mathcal{R}}$. A normal form is a term without redexes. We say t has a normal form if $t \to^*_{\mathcal{R}} s$ for some normal form s. The set of normal forms of a term rewriting system $\mathcal{R}$ is denoted by $\mathrm{NF}_{\mathcal{R}}$. When no confusion can arise, we omit the subscript $\mathcal{R}$.

A term rewriting system $\mathcal{R}$ is left-linear if for any $l \to r \in \mathcal{R}$, every variable in l occurs only once. In this paper we restrict ourselves to the class of left-linear term rewriting systems.

3 NVNF-sequentiality

In this section we will explain NVNF-sequentiality. In order to define this concept, we need some preliminaries.

Let Ω be a new constant symbol representing an unknown part of a term. The set $\mathcal{T}(\mathcal{F} \cup \{\Omega\}, \mathcal{V})$ is abbreviated to $\mathcal{T}_\Omega$. Elements of $\mathcal{T}_\Omega$ are called Ω-terms. An Ω-normal form is an Ω-term without redexes, and the set of all Ω-normal forms is denoted by NF_Ω. Only terms containing neither redexs nor Ω's are said to be normal form. t_Ω denotes the Ω-term obtained from t by replacing all variables in t by Ω, and t_x denotes the term obtained from t by replacing all Ω by x. $O_\Omega(t)$ denotes the set of Ω-occurrences of t, i.e. $O_\Omega(t) = \{u \in O(t) \mid t/u \equiv \Omega\}$. The prefix ordering $\leq$ on $\mathcal{T}_\Omega$ is defined as follows:

(i) $\Omega \leq t$ for all $t \in \mathcal{T}_\Omega$,
(ii) $F(s_1, \ldots, s_n) \leq F(t_1, \ldots, t_n)$ if $s_i \leq t_i$ $(1 \leq i \leq n)$,
(iii) $x \leq x$ for all $x \in \mathcal{V}$.

Two Ω-terms t and s are compatible, written $t \uparrow s$, if there exists an Ω-term r such that $t \leq r$ and $s \leq r$. In this case the least upper bound of t and s is denoted by $t \sqcup s$.

Definition 3.1 ([2]) *Let P be a predicate on $\mathcal{T}_\Omega$. An Ω-occurrence u of t is index with respect to P if for all Ω-term s, $s \geq t$ and $P(s)$ imply $s/u \not\equiv \Omega$.*

The set of indices of t with respect to P is denoted by $I_P(t)$. Intuitively, if u is an index with respect to P in t, then the term at u has to been evaluated in order to make the predicate P true.

Definition 3.2 ([6])

(1) *The reduction relation is defined as follow: $t \to_{nv} s$ iff there exists $l \to r \in \mathcal{R}$, $u \in O(t)$ such that $t/u \geq l_\Omega$ and $s \equiv t[u \leftarrow s']$ for some $s' \geq r_\Omega$.*

(2) *The predicate nvnf on $\mathcal{T}_\Omega$ as follow: $nvnf(t)$ holds iff $t \to^*_{nv} s$ for some s in normal form.*

Definition 3.3 *A left-linear term rewriting system is NVNF-sequential if every Ω-normal form containing at least one occurrence of Ω has index with respect to nvnf.*

In definition of NV-sequentiality, a predicate *term* is used, where $term(t)$ holds iff $t \to_{nv}^* s$ for some $s \in \mathcal{T}$. Note that s may be a reducible term.

Example 3.4 Let

$$\mathcal{R} = \begin{cases} F(A,B,x) \to A \\ F(B,x,A) \to B \\ F(x,A,B) \to C \\ C \to C. \end{cases}$$

Consider the Ω-term $t \equiv F(\Omega,\Omega,\Omega)$. $I_{nvnf}(t) = \{1\}$ because there does not exist an Ω-term s such that $s \geq t$, $s/1 \equiv \Omega$ and $s \to_{nv}^* s'$ for some normal form s'. Note that $\mathcal{R}$ is not NV-sequential since $F(\Omega,\Omega,\Omega)$ has no indices with respect to the predicate *term*.

Next we show that $\mathcal{R}$ in Example 3.4 is NVNF-sequential system. For this purpose we need the following lemma.

Lemma 3.5 *Let $t \in \mathcal{T}_\Omega$. If $u \in I_{nvnf}(t)$, $t \leq s$ and $s/u \equiv \Omega$ then $u \in I_{nvnf}(s)$.*

Proof. If $u \notin I_{nvnf}(s)$ then there exists $s' \geq s$ such that $s'/u \equiv \Omega$ and $nvnf(s')$ is *true*. Since $s' \geq t$, $u \notin I_{nvnf}(t)$. □

Lemma 3.6 *$\mathcal{R}$ of Example 3.4 is NVNF-sequential system.*

Proof. We first prove the following claim: If $u \in I_{nvnf}(t)$ and $v \in I_{nvnf}(s)$ then $u.v \in I_{nvnf}(t[u \leftarrow s])$.

Proof of the claim. Suppose $u.v \notin I_{nvnf}(t[u \leftarrow s])$. Then there exists $t' \geq t[u \leftarrow s]$ such that $t'/u.v \equiv \Omega$ and $nvnf(t')$ is *true*. Hence there exists a reduction

$$t' \equiv t_0 \xrightarrow{u_0}_{nv} t_1 \xrightarrow{u_1}_{nv} \cdots\cdots \xrightarrow{u_{n-1}}_{nv} t_n \in \mathrm{NF}.$$

We distinguish two case:

(1) $u_i \not< u$ for all i $(0 \leq i \leq n-1)$. Because there exists no u_i such that $u_i < u$, we have $t'/u \to_{nv}^* t_n/u \in \mathrm{NF}$. Hence $nvnf(t'/u)$ is *true*. Clearly $t'/u > s$ and $t'/u.v \equiv \Omega$. This contradicts the assumption $v \in I_{nvnf}(s)$.

(2) $u_i < u$ for some i. Let j be the smallest number satisfying $u_j < v$. Note that $t'/u \to_{nv}^* t_j/u$. t_j/u_j is a redex but $t_j/u_j \not\equiv C$. Moreover $t_j \not\equiv A$ and $t_j \not\equiv B$ because $t'/u \geq s$, $t'/u.v \equiv \Omega$ and $v \in I_{nvnf}(s)$. Thus $t_j[u \leftarrow \Omega] \xrightarrow{u_j}_{nv} t_{j+1}$. We can obtain the following reduction: $t'[u \leftarrow \Omega] \to_{nv}^* t_j[u \leftarrow \Omega] \to_{nv} t_{j+1} \to_{nv}^* t_n$. Hence $t'[u \leftarrow \Omega] \geq t$, $nvnf(t'[u \leftarrow \Omega])$ is *true*. But this is contradictory to $u \in I_{nvnf}(t)$.

Therefore the claim follows.

Let t be an Ω-normal form containing at least one occurrence of Ω. We prove, by induction on size of t that t has an index with respect to *nvnf*. When $t \equiv \Omega$, it is clear that t has an index. Induction step:

(1) $t \equiv F(t_1,t_2,t_3)$. If t_1 contains Ω then by induction hypothesis, t_1 has an index. By $1 \in I_{nvnf}(F(\Omega,\Omega,\Omega))$ and Lemma 3.5, $1 \in I_{nvnf}(F(\Omega,t_2,t_3))$. Therefore, from the claim it follows that t has an index. Otherwise we distinguish three cases:

(1-1) $t_1 \equiv A$. If t_2 contains Ω then by induction hypothesis, t_2 has an index. By $2 \in I_{nvnf}(F(A,\Omega,\Omega))$ and Lemma 3.5, $2 \in I_{nvnf}(F(A,\Omega,t_3))$. From the claim, t has an index. Otherwise t_3 contains Ω. By induction hypothesis, t_3 has an index. We can obtain $3 \in I_{nvnf}(F(A,t_2,\Omega))$. Therefor by the claim, t has an index.

(1-2) $t_1 \equiv B$. Similar to (1-1).

(1-3) Otherwise we have $I_{nvnv}(F(t_1,\Omega,\Omega)) = \{2, 3\}$. By induction hypothesis, t_2 or t_3 has an index. By Lemma 3.5 and the claim, t has an index.

(2) $t \equiv G(t_1,\cdots,t_n)$ $(G \not\equiv F)$. Suppose t_i contains Ω. Then by induction hypothesis, t_i has an index. Because $i \in I_{nvnf}(G(t_1,\cdots,t_{i-1},\Omega,t_{i+1},\cdots,t_n))$, t has an index from the claim. □

By Example 3.4 and Lemma 3.6 we have the following theorem.

Theorem 3.7 *The class of NVNF-sequential term rewriting systems properly includes the class of NV-sequential systems.*

Proof. NV-sequential system is NVNF-sequential because an index with respect to *term* is also an index with respect to *nvnf*. $\mathcal{R}$ of Example 3.4 is NVNF-sequential but not NV-sequential. Therefore the inclusion is proper. □

4 Indices with respect to NVNF-sequentiality

In this section we show that for a given $t \in \mathcal{T}_\Omega$, it is decidable whether $u \in O_\Omega(t)$ is an index with respect to *nvnf* in t. We introduce the reduction $\to_\omega$ which is given in [6].

Definition 4.1 ([6]) *The reduction relation $\to_\omega$ is defined as follows: $t \to_\omega s$ iff there exists $l \to r \in \mathcal{R}$, $u \in O(t)$ such that $t/u \uparrow l_\Omega$, $t/u \not\equiv \Omega$ and $s \equiv t[u \leftarrow r_\Omega]$.*

We explain a relationship between this reduction $\to_\omega$ and $\to_{nv}$, and show the condition for ensuring that Ω-occurrence is an index with respect to *nvnf*.

Lemma 4.2

(1) *If $t \to_{nv}^* s$ and $t' \le t$ then $t' \to_\omega^* s'$ for some $s' \le s$.*

(2) *If $t \to_\omega^* s$ then $t' \to_{nv}^* s$ for some $t' \ge t$.*

Proof.

(1) We prove that if $t \to_{nv} s$ and $t' \le t$ then $t' \to_\omega^= s'$ for some $s' \le s$. If $t \to_{nv} s$ then there exist $l \to r \in \mathcal{R}$, $u \in O(t)$ such that $t/u \ge l_\Omega$ and $s \equiv t[u \leftarrow s_1]$ for some $s_1 \ge r_\Omega$. If $u \in O(t')$, $t'/u \not\equiv \Omega$ then $t'/u \uparrow l_\Omega$. Hence $t' \to_\omega t'[u \leftarrow r_\Omega]$ and $t'[u \leftarrow r_\Omega] \le s$. Otherwise it is clear that $t' \le s$. Using this fact, we can prove (1) by induction on length of $t \to_{nv}^* s$.

(2) This is proved by induction on length of $t \to_\omega^* s$. The case of zero is trivial. Assume that $t \to_\omega s_1 \to_\omega^* s$ where $t/u \uparrow l_\Omega$, $t/u \not\equiv \Omega$ and $s_1 \equiv t[u \leftarrow r_\Omega]$ for some $l \to r \in \mathcal{R}$ and $u \in O(t)$. From induction hypothesis, for some $s_2 \geq s_1$, $s_2 \to_{nv}^* s$. Let $t_1 \equiv t/u \sqcup l_\Omega$. Define $t' \equiv s_2[u \leftarrow t_1]$. Because $s_2 \geq s_1 \equiv t[u \leftarrow r_\Omega]$, $t' \equiv s_2[u \leftarrow t_1] \geq t$. We have $t' \to_{nv} s_2$ by $s_2/u \geq r_\Omega$. Therefore $t' \to_{nv}^* s$. □

Lemma 4.3 *Let $t \in \mathcal{T}_\Omega$ and $u \in O_\Omega(t)$. Let $\bullet$ be a fresh constant symbol. Then $u \notin I_{nvnf}(t)$ iff there exists $s \in \mathrm{NF}_\Omega$ such that $t[u \leftarrow \bullet] \to_\omega^* s$ and $\bullet \not\subseteq s$.*

Proof. *only-if part.* If $u \notin I_{nvnf}(t)$ then there exists $t' \geq t$ such that $t'/u \equiv \Omega$ and $nvnf(t')$ is *true*. Thus $t' \to_{nv}^* s'$ for some normal from s'. From $\Omega \not\subseteq s'$ and left-linearity, we can obtain $t'[u \leftarrow \bullet] \to_{nv}^* s'$ and $\bullet \not\subseteq s'$. Using Lemma 4.2 (1), we obtain $s \leq s'$ such that $t[u \leftarrow \bullet] \to_{nv}^* s$. Because s' is a normal form, s is an Ω-normal form and $\bullet \not\subseteq s$.

if part. If $t[u \leftarrow \bullet] \to_\omega^* s \in \mathrm{NF}_\Omega$ and $\bullet \not\subseteq s$, then there exists $t' \geq t[u \leftarrow \bullet]$ such that $t' \to_{nv}^* s$ by Lemma 4.2 (2). Let $t'' \equiv t'_x[u \leftarrow \Omega]$, $s' \equiv s_x$. We can transform the reduction $t' \to_{nv}^* s$ into $t'' \to_{nv}^* s'$. Because s is an Ω normal from, s' is a normal form and hence $nvnf(t'')$ is *true*. Clearly $t'' \geq t$ and $t''/u \equiv \Omega$. Therefore $u \notin I_{nvnf}(t)$. □

Next we show that for any $t \in \mathcal{T}_\Omega$, there exists an upper bound of the least hight of Ω-normal form which t can reduce by $\to_\omega$ when it exists.

For given a term rewriting system $\mathcal{R}$, $Rh_\mathcal{R}$ is defined with $Rh_\mathcal{R} = \{ r_\Omega \mid l \to r \in \mathcal{R}\}$, and $Rhn_\mathcal{R}$ is the smallest set such that $Rh_\mathcal{R} \subseteq Rhn_\mathcal{R}$ and if $t \in Rhn_\mathcal{R}$, $u \in O(t)$ and $r \in Rh_\mathcal{R}$ then $t[u \leftarrow r] \in Rhn_\mathcal{R}$. It is clear that if $r \in Rh_\mathcal{R}$ and $r \to_\omega^* t$ then $t \in Rhn_\mathcal{R}$. In the sequel we often omit the subscript $\mathcal{R}$.

Lemma 4.4 *If $t \to_\omega^+ s$ then there exists $u_1, \ldots, u_n \in O(t)$ which are pairwise disjoint, and the following conditions hold.*

(i) $s \equiv t[\,u_i \leftarrow s/u_i \mid 1 \leq i \leq n\,]$,

(ii) $t/u_i \to_\omega^+ r_i \to_\omega^* s/u_i$ *for some* $r_i \in Rh_\mathcal{R}$, $1 \leq i \leq n$.

Proof. By $t \to_\omega^+ s$, there exists a reduction

$$t \equiv t_0 \overset{u_0}{\to}_\omega t_1 \overset{u_1}{\to}_\omega \cdots\cdots \overset{u_{n-1}}{\to}_\omega t_n \equiv s \quad (n > 0)\,.$$

Let $\{u_{i_1}, \ldots, u_{i_k}\}$ be the set of minimal redex occurrences of $\{u_0, u_1, \cdots, u_{n-1}\}$. Then $u_{i_1}, \ldots, u_{i_k} \in O(t)$ are pairwise disjoint. By minimality of $u_{i_1}, \ldots, u_{i_k}$, (i), (ii) hold. □

We use $|\,u\,|$ for the length of the word u. The height $|\,t\,|$ of an Ω-term t is defined by $|\,t\,| = \max\{\,|\,u\,| \mid u \in O(t)\,\}$. The maximum hight of the left-hand sides and right-hand sides of a given $\mathcal{R}$ is denote by $\rho_\mathcal{R}$. We write ρ when confusion does not occur. $(t)_\rho$ is a prefix term of t whose hight is ρ, i.e., $(t)_\rho \equiv t[\,u \leftarrow \Omega \mid u \in O(t) \wedge |\,u\,| = \rho\,]$.

Lemma 4.5 *Let $r \in Rh_\mathcal{R}$, $r \to_\omega^* s$ where $|\,s\,| > \rho \times n$ $(n > 0)$. Then, there exists $\varepsilon < u_0 < \cdots < u_{n-1} \in O(s)$ and for any i $(\,0 \leq i \leq n-1\,)$, the following condition holds: $r \to_\omega^* s[u_i \leftarrow r_i]$ and $r_i \to_\omega^* s/u_i$ for some $r_i \in Rh_\mathcal{R}$.*

Proof. The proof is by induction on n. By $r \rightarrow_\omega^* s$, there exists a reduction

$$r \equiv t_0 \xrightarrow{v_0}_\omega t_1 \xrightarrow{v_1}_\omega \cdots\cdots \xrightarrow{v_{m-1}}_\omega t_m \equiv s.$$

Because $|s| > \rho \times n \ (n > 0)$, we can obtain $j < m$ such that $t_j \in Rh_\mathcal{R}$ and $v_i \neq \varepsilon$ for all $i \ (j \leq i \leq m-1)$. When $n = 1$, using Lemma 4.4, we can easily show that there exists $u \in O(s)$ such that $u \neq \varepsilon$ and $t_j \rightarrow_\omega^* s[u \leftarrow r']$, $r' \rightarrow_\omega^* s/u$ for some $r' \in Rh_\mathcal{R}$. By $r \rightarrow^* t_j$, $r \rightarrow_\omega^* s[u \leftarrow r']$. Suppose $n > 1$. Using Lemma 4.4, we can obtain $u \neq \varepsilon$ such that $t_j \rightarrow_\omega^* s[u \leftarrow r']$, $r' \rightarrow_\omega^* s/u$ and $|s/u| > \rho \times (n-1)$ for some $r' \in Rh_\mathcal{R}$ By induction hypothesis, there exists $\varepsilon < u'_1 < \cdots < u'_{n-1} \in O(s)$ and for any $i \ (1 \leq i \leq n-1)$, the following condition holds: $r' \rightarrow_\omega^* (s/u)[u'_i \leftarrow r_i] \equiv s[u.u'_i \leftarrow r_i]/u$ and $r_i \rightarrow_\omega^* (s/u)/u'_i \equiv s/u.u'_i$ for some $r_i \in Rh_\mathcal{R}$. Let $u_0 = u$, $u_i = u.u'_i \ (1 \leq i \leq n-1)$. Clearly $\varepsilon < u_0 < \cdots < u_{n-1} \in O(s)$ and $r \rightarrow_\omega^* s[u \leftarrow r']$, $r' \rightarrow_\omega^* s/u$. For all $i \ (1 \leq i \leq n-1)$, we have $t_j \rightarrow_\omega^* s[u \leftarrow r'] \rightarrow_\omega^* s[u \leftarrow s[u.u'_i \leftarrow r_i]/u] \equiv s[u_i \leftarrow r_i]$ and $r_i \rightarrow_\omega^* s/u_i$. By $r \rightarrow_\omega^* t_j$, $r \rightarrow_\omega^* s[u_i \leftarrow r_i]$. Therefore the lemma holds. □

Lemma 4.6 *Let* $t, s \in \mathrm{NF}_\Omega$. *If* $(t/u)_\rho \equiv (s)_\rho$ *then* $t\,[u \leftarrow s] \in \mathrm{NF}_\Omega$.

Proof. Trivial. □

For given a term rewriting system $\mathcal{R}$, let τ , σ, and $k_\mathcal{R}$ be constants defined as follow: $\tau = \|\{\ (t)_\rho \mid t \in Rhn_\mathcal{R}\ \}\|$, $\sigma = \|\mathcal{R}\|$ and $k_\mathcal{R} = \rho_\mathcal{R} \times (\tau \times \sigma + 1)$, where $\|A\|$ is the cardinality of a set A.

Lemma 4.7 *Let* $t \in \mathcal{T}_\Omega$ *and* $u \in O_\Omega(t)$. *Let* $\bullet$ *be a fresh constant symbol. Then* $u \notin I_{nvnf}(t)$ *iff there exists* $s \in \mathrm{NF}_\Omega$ *such that* $|s| \leq |t| + k_\mathcal{R}$, $t[u \leftarrow \bullet] \rightarrow_\omega^* s$ *and* $\bullet \not\subseteq s$.

Proof. *if part.* By Lemma 4.3.

only-if part. If $u \notin I_{nvnf}(t)$ then by Lemma 4.3 there exists $s \in \mathrm{NF}_\Omega$ such that $t[u \leftarrow \bullet] \rightarrow_\omega^* s$ and $\bullet \not\subseteq s$. Let s be an Ω-normal form with the least size satisfying this condition. Suppose $|s| > |t| + k_\mathcal{R}$. By Lemma 4.5 and the definition of $k_\mathcal{R}$, we can show that there exists $r \in Rh$, $u_1, u_2 \in O(s)$ $(u_1 < u_2)$ such that $(s/u_1)_\rho \equiv (s/u_2)_\rho$, $t\,[u \leftarrow \bullet] \rightarrow_\omega^* s[u_i \leftarrow r]$ and $r \rightarrow_\omega^* s/u_i$ for $i = 1, 2$, see Figure 4.1. Let $s' \equiv s[u_1 \leftarrow s/u_2]$. By Lemma 4.6, $s' \in \mathrm{NF}_\Omega$ and $t[u \leftarrow \bullet] \rightarrow_\omega^* s'$, $\bullet \not\subseteq s'$. Because the size of s' is smaller than the size of s, we obtain a contradiction. □

By Lemma 4.7, in order to determine whether an Ω-occurrence is an index w.r.t. $nvnf$, we need to check the reachability to a finite number of Ω-normal form. For a term rewriting system $\mathcal{R}$, we define $\mathcal{R}^\Omega$ as follow [6]:

$$\mathcal{R}^\Omega = \{\, l \rightarrow r_\Omega \mid l \rightarrow r \in \mathcal{R}\,\} \cup \{\, \Omega \rightarrow t \mid t \subseteq l_\Omega\,,\ l \rightarrow r \in \mathcal{R}\,\}\,.$$

We can prove that in the condition of Lemma 4.7 $\rightarrow_\omega^*$ can be replaced $\rightarrow_{\mathcal{R}^\Omega}^*$ which is the transitive-reflexive closure of a usual reduction relation $\rightarrow_{\mathcal{R}^\Omega}$ defined by $\mathcal{R}^\Omega$.

Lemma 4.8 ([6])

(1) *If* $t \rightarrow_\omega^* s$ *then* $t \rightarrow_{\mathcal{R}^\Omega}^* s$.

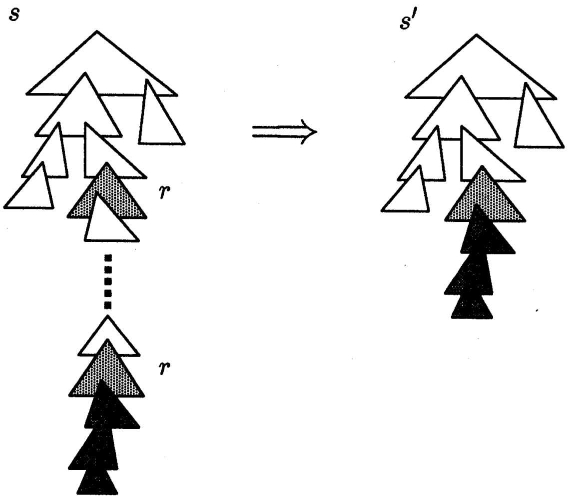

Figure 4.1

(2) *If* $t \rightarrow^*_{\mathcal{R}^\Omega} s$ *and* $t' \leq t$ *then* $t' \rightarrow^*_\omega s'$ *for some* $s' \leq s$.

Lemma 4.9 *Let* $t \in \mathcal{T}_\Omega$ *and* $u \in O_\Omega(t)$. *Then* $u \notin I_{nvnf}(t)$ *iff there exists* $s \in \mathrm{NF}_\Omega$ *such that* $|s| \leq |t| + k_{\mathcal{R}}$, $t[u \leftarrow \bullet] \rightarrow^*_{\mathcal{R}^\Omega} s$ *and* $\bullet \not\subseteq s$.

Proof. By Lemma 4.7 and Lemma 4.8 □

We assume that $\mathcal{R}$ is left-linear, so $\mathcal{R}^\Omega$ is left-linear and right-ground (i.e. r is ground term for any $l \rightarrow r \in \mathcal{R}^\Omega$). It is show that the reachability problem is decidable for left-linear and right-ground systems [1, 7]. Thus we obtain the following theorem.

Theorem 4.10 *It is decidable, for* $t \in \mathcal{T}_\Omega$, $u \in O_\Omega(t)$, *whether* u *is an index with respect to nvnf in* t.

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