Polynomial Time Learnabilities of Tree Patterns with Internal Structured Variables from Queries

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Abstract

We give the polynomial time learnabilities of two classes of ordered tree patterns with internal structured variables, in the query learning model of Angluin (1988). An ordered tree pattern with internal structured variables, called a term tree, is a rooted tree pattern which consists of tree structures, ordered children and internal structured variables. A term tree is suited for representing structural features in semistructured or tree structured data such as HTML/XML files. We show the polynomial time learnabilities of two classes of term trees using membership and restricted subset queries and one positive example.

1 Introduction

Large amount of Web documents such as HTML/XML files are available. Such documents are called semistructured data and considered tree structured data, which are represented by rooted trees with ordered children and edge labels [1]. As an example of a representation of tree structured data, we give a rooted tree \( T \) in Fig. 1. This work is motivated from data mining of tree structured patterns from semistructured data.

As a representation of a tree structured pattern, we use an ordered tree pattern with internal structured variables, called a term tree. A term tree is a rooted tree pattern which consists of tree structures, ordered children and internal structured variables. A variable in a term tree is a list of vertices and it can be substituted by an arbitrary tree. A term tree is more powerful than or incomparable to other representations of tree structured patterns, which were proposed in computational learning theory, such as ordered tree patterns [2] and ordered gapped tree patterns [7]. We can show that a term tree is more powerful than an ordered tree pattern, which is also called a first order term in formal logic. Consider the example in Fig. 1. The tree pattern \( f(b, x, g(a, z), y) \) can be represented by the term tree \( s \), but the term tree \( t \) cannot be represented by any ordered tree pattern because of the existence of internal structured variables represented by \( z_2 \) and \( z_3 \) in \( t \). The variable represented by \( z_3 \) in \( t \) is a list of vertices \( [v_8,v_7,v_9] \).
For a set of edge labels $\Lambda$, the term tree language $L_{\Lambda}(t)$ of a term tree $t$ with $\Lambda$, which denotes the representing power of $t$, is the set of all labeled trees which are obtained from $t$ by substituting arbitrary labeled trees for all variables in $t$. The subtrees which are obtained from $t$ by removing the variables in $t$ represent the common subtree structures in the trees in $L_{\Lambda}(t)$. A term tree $t$ is said to be regular if all variable labels in $t$ are mutually distinct.

$\mathcal{OTT}_{\Lambda}$ denotes the set of all regular term trees with $\Lambda$ as a set of edge labels. Let $t$ be a term tree in $\mathcal{OTT}_{\Lambda}$. For a variable $h = [u_0, u_1, \ldots, u_t]$ in $t$, we define $\text{parent}(h) = u_0$ and $\text{child}(h) = \{u_1, \ldots, u_t\}$. We denote by $z\mathcal{OTT}_{\Lambda}$ the set of all term trees $t \in \mathcal{OTT}_{\Lambda}$ with a variable set $H_t$ such that $\text{parent}(h) \notin \text{child}(h_2)$ for any $h_1$ and $h_2$ in $H_t$.

In query learning model, a learning algorithm accesses to oracles, which answer specific kinds of queries, and collect information about a target term tree $t_*$. We consider the following oracles. Membership oracle: The input is a term tree $t$ having no variable. The output is “yes” if $t \in L_{\Lambda}(t_*)$, and “no” otherwise. Restricted subset oracle: The input is a term tree $t$ in $\mathcal{OTT}_{\Lambda}$. The output is “yes” if $L_{\Lambda}(t) \subseteq L_{\Lambda}(t_*)$, and “no” otherwise. The former is called a membership query and the latter is called a restricted subset query. In this model, a learning algorithm is said to exactly identify a target term tree $t_*$ if it outputs a term tree $t$ such that $L_{\Lambda}(t) = L_{\Lambda}(t_*)$ and halts, after it uses some queries.

In this paper, we assume $|\Lambda| \geq 2$. We show that any term tree in $\mathcal{OTT}_{\Lambda}$ is exactly identifiable in polynomial time using at most $n^2 + n$ membership queries, at most $n$ restricted subset queries and one positive example, where $n$ is the size of the positive example. Moreover, we show that any term tree in $z\mathcal{OTT}_{\Lambda}$ is exactly identifiable in polynomial time using at most $n^2 + 2n$ membership queries and one positive example, where $n$ is the size of the positive example.

As our previous works, we showed the learnabilities of graph structured patterns\cite{8}, term trees with unordered children\cite{10} in the framework of polynomial time inductive inference from positive data\cite{4}. Also our work\cite{11} showed that the class $\mathcal{OTT}_{\Lambda}$, a subclass of $\mathcal{OTT}_{\Lambda}$, is polynomial time inductively inferable from positive data. As an application\cite{9}, we gave a data mining method from semistructured data by using a learning algorithm for term trees. As other related works, the works\cite{2,3,6,7} showed the learnabilities of tree structured patterns in query learning model. The tree structured patterns and learning models of this work are incomparable to those of all the other related works.

This paper is organized as follows. In Section 2, we explain term trees as tree structured patterns. In Section 3, we explain the query learning model. In Section 4, we show the above learnabilities of the two classes $\mathcal{OTT}_{\Lambda}$ and $z\mathcal{OTT}_{\Lambda}$.

2 Preliminaries

Let $T = (V_T, E_T)$ be an ordered tree with a vertex set $V_T$ and an edge set $E_T$. A list $h = [u_0, u_1, \ldots, u_t]$ of vertices in $V_T$ is called a variable of $T$ if $u_1, \ldots, u_t$ are consecutive children of $u_0$, i.e., $u_0$ is the parent of $u_1, \ldots, u_t$ and $u_{j+1}$ is the next sibling of $u_j$ for any $j$ with $1 \leq j < t$. We call $u_0$ the parent port of the variable $h$ and $u_1, \ldots, u_t$ the child ports of $h$. Two variables $h = [u_0, u_1, \ldots, u_t]$ and $h' = [u_0', u_1', \ldots, u_{t'}]$ are said to be disjoint if $\{u_1, \ldots, u_t\} \cap \{u_1', \ldots, u_{t'}\} = \emptyset$. For a set $S$, we denote by $|S|$ the number of elements in $S$.

**Definition 1** Let $T = (V_T, E_T)$ be an ordered tree and $H_T$ a set of pairwise disjoint variables of $T$. An ordered term tree obtained from $T$ and $H_T$ is a triplet $t = (V_t, E_t, H_t)$, where $V_t = V_T$, $E_t = E_T - \bigcup_{u_0, u_1, \ldots, u_t \in H_T} \{u_0, u_1\} \in E_T \mid 1 \leq i \leq t\}$ and $H_t = H_T$.

For two vertices $u, u' \in V_t$, we say that $u$ is the parent of $u'$ in $t$ if $u$ is the parent of $u'$ in $T$. Similarly we say that $u'$ is a child of $u$ in $t$ if $u'$ is a child of $u$ in $T$. In particular, for a vertex $u \in V_t$ with no child, we call $u$ a leaf of $t$. We define the order of the children of each vertex $u$ in $t$ as the order of the children of $u$ in $T$. We often omit the description of the ordered tree $T$ and variable set $H_T$ because we can find them from the triplet $t = (V_t, E_t, H_t)$. We define the size of $t$ as the number of vertices in $t$ and denote it by $|t|$.
For example, the ordered term tree $t$ in Fig. 1 is obtained from the tree $T = (V_T, E_T)$ and the set of variables $H_T$ defined as follows. $V_T = \{v_1, \ldots, v_{11}\}$, $E_T = \{(v_1, v_2), (v_1, v_3), (v_1, v_4), (v_4, v_5), (v_5, v_6), (v_6, v_7), (v_8, v_9), (v_9, v_10), (v_10, v_11)\}$ with the root $v_1$ and the sibling relation displayed in Fig. 1. $H_T = \{(v_4, v_6), (v_1, v_6), (v_6, v_7, v_9)\}$.

For any ordered term tree $t$, a vertex $u$ of $t$, and two children $u'$ and $u''$ of $u$, we write $u' <_u u''$ if $u'$ is smaller than $u''$ in the order of the children of $u$. We assume that every edge and variable of an ordered term tree is labeled with some words from specified languages. A label of a variable is called a \textit{variable label}. $\Lambda$ and $X$ denote a set of edge labels and a set of variable labels, respectively, where $\Lambda \cap X = \phi$. An ordered term tree $t = (V_t, E_t, H_t)$ is called \textit{regular} if all variables in $H_t$ have mutually distinct variable labels in $X$.

\textbf{Note.} In this paper, we treat only regular ordered term trees, and then we call a regular ordered term tree a \textit{term tree}, simply. In particular, an ordered term tree with no variable is called a \textit{ground term tree} and considered to be a tree with ordered children.

$OT_\Lambda$ denotes the set of all ground term trees with $\Lambda$ as a set of edge labels. Let $OTT_\Lambda$ be the set of all term trees. In particular, for a positive integer $L$, we denote by $OTT_\Lambda^L$ the set of all term trees with $\Lambda$ as a set of edge labels such that each variable in $t$ has at most $L$ child ports.

Let $f = (V_f, E_f, H_f)$ and $g = (V_g, E_g, H_g)$ be term trees. We say that $f$ and $g$ are \textit{isomorphic}, denoted by $f \equiv g$, if there is a bijection $\varphi$ from $V_f$ to $V_g$ such that (i) the root of $f$ is mapped to the root of $g$ by $\varphi$, (ii) $\{u, u'\} \in E_f$ if and only if $\{\varphi(u), \varphi(u')\} \in E_g$ and the two edges have the same edge label, (iii) $\{u_0, u_1, \ldots, u_\ell\} \in H_f$ if and only if $[\varphi(u_0), \varphi(u_1), \ldots, \varphi(u_\ell)] \in H_g$, and (iv) for any vertex $u$ in $f$ which has more than one child, and for any two children $u'$ and $u''$ of $u$, $u' <_u u''$ if and only if $\varphi(u') <_{\varphi(u)} \varphi(u'')$. We say that an edge $\{u, u'\} \in E_f$ \textit{corresponds} to an edge $\{v, v'\} \in E_g$ if $v = \varphi(u)$ and $v' = \varphi(u')$.

Let $f$ and $g$ be term trees with at least two vertices. Let $h = [v_0, v_1, \ldots, v_j]$ be a variable in $f$ with the variable label $x$ and $\sigma = [v_0, u_1, \ldots, u_L]$ a list of $L + 1$ distinct vertices in $g$ where $v_0$ is the root of $g$ and $u_1, \ldots, u_L$ are leaves of $g$. The form $x := [g, \sigma]$ is called a \textit{binding} for $x$. A new term tree $f' = f(x := [g, \sigma])$ is obtained by applying the binding $x := [g, \sigma]$ to $f$ in the following way. For the variable $h = [v_0, v_1, \ldots, v_j]$, we attach $g$ to $f$ by removing the variable $h$ from $H_f$ and by identifying the vertices $v_0, v_1, \ldots, v_j$ with the vertices $u_0, u_1, \ldots, u_L$ of $g$ in this order. We define a new ordering $<_t'$ on every vertex $v$ in $f'$ in the following natural way. Suppose that $v$ has more than one child and let $v'$ and $v''$ be two children of $v$ in $f'$. We note that $v_i = u_i$ for any $0 \leq i \leq L$. (1) If $v, v', v'' \in V_g$ and $v' <_t v''$, then $v' <_{t'} v''$. (2) If $v, v', v'' \in V_f$ and $v' <_t v''$, then $v' <_{t'} v''$. (3) If $v = v_0(= u_0)$,
v' ∈ V_f - {v_1, . . . , v_t}, v'' ∈ V_g, and v' <'_v v_1, then v' <'_v v''. (4) If v = v_0 (= u_0), v' ∈ V_f - {v_1, . . . , u_t}, v'' ∈ V_g, and u_2 <'_v v', then v'' <'_v v'. In Fig. 2, we give an example of the new ordering on vertices in a term tree.

A substitution θ is a finite collection of bindings \{x_1 := [g_1, σ_1], . . . , x_n := [g_n, σ_n]\}, where x_i's are mutually distinct variable labels in X. The term tree fθ, called the instance of f by θ, is obtained by applying all the bindings x_i := [g_i, σ_i] on f. We define the root of the resulting term tree fθ as the root of f. Consider the examples in Fig. 1. An example of a term tree t is given. Let θ = \{x_1 := [g_1, u_1, w_1], x_2 := [g_2, [u_2, w_2]], x_3 := [g_3, [u_3, w_3, w_3']\} be a substitution, where g_1, g_2, and g_3 are ground term trees in Fig. 1. Then the instance θt of the term tree t by θ is isomorphic to the tree T in Fig. 1. Let t and t' be term trees. We write t ≤ t' if there exists a substitution θ such that t ≡ t'θ. If t ≤ t' and t ≠ t', then t < t'. Let Λ be a set of edge labels. The term tree language L_{Λ}(t) of a term tree t ∈ OTT_{Λ} is \{s ∈ OT_{Λ} | s ≤ t\}.

3 Learning model

In this paper, let t* be a term tree in OTT_{Λ} to be identified, and we say that the term tree t* is a target. A ground term tree t is called a positive example of L_{Λ}(t*) if t is in L_{Λ}(t*).

We introduce the exact learning model via queries due to Angluin [5]. In this model, learning algorithms can access to oracles that answer specific kinds of queries about the unknown term tree language L_{Λ}(t*). We consider the following oracles: (1) Membership oracle Mem_{Λ}: The input is a ground term tree t. The output is “yes” if t is in L_{Λ}(t*), and “no” otherwise. The query is called a membership query. (2) Restricted subset oracle rSub_{Λ}: The input is a term tree t in OTT_{Λ}. The output is “yes” if L_{Λ}(t) ⊆ L_{Λ}(t*), and “no” otherwise. The query is called a restricted subset query.

A learning algorithm A may collect information about membership and restricted subset queries of L_{Λ}(t*). We say that a learning algorithm exactly identifies a target t* if it outputs a term tree t in OTT_{Λ} with L_{Λ}(t) = L_{Λ}(t*) and halts after it uses some queries.

4 Learning using membership and restricted subset queries

We introduce an operation of a contraction which reduces the number of edges in a ground term tree. In this paper, we assume |Λ| ≥ 2.

Definition 2 Let t = (V_t, E_t, H_t) be a ground term tree and e = (u, v) an edge in E_t. We define the contraction of e to t as the following operation: If v has children v_1, . . . , v_t, then the operation removes v from V_t and replaces \{u, v\}, \{v_1, v\}, . . . , \{v_t, v\} in E_t with new edges \{u, v_1\}, . . . , \{u, v_t\}, that is, E_t = E_t - \{u, v\} + \{u, v_1\}, . . . , \{u, v_t\}. Otherwise, the operation removes v from V_t and e from E_t. We denote by t\{e\} the term tree obtained from t by applying the contraction...
Algorithm *CONTRACTION*;

Given: An oracle Mem, for a target $t_*$ in $OTT_A$ and a positive example $t$ in $L_A(t_*);
Output: A term tree $r$ in $OTT_A$ with $r \equiv \text{port}(t_*)$;
begin
begin

foreach edge $e$ in $t$ do begin Let $t' := t \setminus \{e\}$;
if $\text{Mem}(t') = \text{"yes"}$ then begin $t := t'$; break; end; end;
until $t$ does not change;
Let $r = (V_r, E_r, H_r)$ be $t$;
foreach edge $e = \{u, v\}$ in $t$ do begin
Let $t'$ be a term tree obtained from $t$ by replacing the label of $e$ with another label;
if $\text{Mem}_e(t') = \text{"yes"}$ then begin
Let $\{u', v'\}$ be an edge in $r$ which corresponds to $\{u, v\}$ in $t$.
$E_r := E_r - \{\{u', v'\}\}; H_r := H_r \cup \{[u', v']\}$; end;
end; output $r$;
end

end

Figure 3: Algorithm *CONTRACTION*

Algorithm *LEARN.OTT*;

Given: Oracles $rSub_*$ and Mem, for a target $t_*$ in $OTT_A$ and a positive example $t$ in $L_A(t_*);
Output: A term tree $r$ in $OTT_A$ with $L_A(r) = L_A(t_*)$;
begin
$r := \text{CONTRACTION}(t)$ using Mem,
let $H = [h_1, \ldots, h_t]$ be the sequence of variables in $r$ by the breath-first search order;
i := 1; flag := false;
while $H$ is not empty do begin Let $S = \{h_i\}$;
repeat
Let $h_i = [u_i, v_i]$ and $h_{i+1} = [u_{i+1}, v_{i+1}]$;
if $v_{i+1}$ is the next sibling of $v_i$ then begin
Let $S := S \cup \{h_{i+1}\}; r' := \text{replace}(r, S)$;
if $rSub_*(r') = \text{"yes"}$ then begin flag := false; $i := i + 1; \text{end}$
else begin flag := true; $S := S - \{h_{i+1}\}; \text{end}$
else flag := true;
until flag
$r := \text{replace}(r, S)$; remove the variables in $S$ from $H$; $i := i + 1$;
end; output $r$;
end

end

Figure 4: Algorithm *LEARN.OTT*

Let $t = (V_t, E_t, H_t)$ be a term tree in $OTT_A$. We denote by $\text{port}(t)$ the term tree obtained from $t$ by replacing any variable $[v_0, v_1, \ldots, v_k]$ with $k$ variables $[v_0, v_1], [v_0, v_k]$. Hence, for $t \in OTT_A$, $\text{port}(t)$ is in $OTT_A$.

Let $k$ be a positive integer, $t = (V_t, E_t, H_t)$ a term tree in $OTT_A$ and $h_1 = [v_0, v_1], \ldots, h_k = [v_0, v_k]$ variables in $H_t$, where $v_{i+1}$ is the next sibling of $v_i$ for each $i = 1, \ldots, k - 1$. We denote by $\text{replace}(t, \{h_1, \ldots, h_k\})$ the term tree obtained from $t$ by replacing the variables $h_1, \ldots, h_k$ with a variable $h = [v_0, v_1, \ldots, v_k]$. That is, $\text{replace}(t, \{h_1, \ldots, h_k\}) = (V_t, E_t, H_t')$, where $H_t' = H_t \cup \{h\} - \{h_1, \ldots, h_k\}$.
Theorem 1 The algorithm *LEARN.OTT* in Fig. 4 exactly identifies any term tree $t_*$ in $OTT_A$ in polynomial time using at most $n^2 + n$ membership queries, at most $n$ restricted subset queries and one positive example $t$ in $L_A(t_*)$, where $n = |t|$.

From Theorem 1, we can identify any term tree in $OTT_A$ using membership, restricted subset queries and a positive example. Next, we show that any term tree in some subset of $OTT_A$ is identifiable using
membership queries and a positive example.

Let $t = (V_t, E_t, H_t)$ be a term tree in $OTT_\Lambda$ and $h = [u_0, u_1, \ldots, u_k]$ a variable in $H_t$. We denote by $\text{parent}(h)$ the parent node of $h$ and by $\text{child}(h)$ the set of the child ports of $h$, that is, $\text{parent}(h) = u_0$ and $\text{child}(h) = \{u_1, \ldots, u_k\}$. We denote by $zOTT_\Lambda$ the set of all term trees $t = (V_t, E_t, H_t)$ in $OTT_\Lambda$ such that $\text{parent}(h_1) \notin \text{child}(h_2)$ for any $h_1$ and $h_2$ in $H_t$.

**Theorem 2** Any term tree $t_*$ in $zOTT_\Lambda$ is exactly identifiable in polynomial time using at most $n^2 + 2n$ membership queries and one positive example $t$ in $L_\Lambda(t_*)$, where $n = |t|$.

## 5 Conclusions

We have considered the polynomial time learnabilities of $OTT_\Lambda$ and $zOTT_\Lambda$ in the query learning model. In this paper, we assume $|\Lambda| \geq 2$. We have shown that any term tree in $OTT_\Lambda$ is exactly identifiable using at most $n^2 + n$ membership queries, at most $n$ restricted subset queries and one positive example, where $n$ is the size of the positive example. Moreover, we have shown that any term tree in $zOTT_\Lambda$ is exactly identifiable using at most $n^2 + 2n$ membership queries and one positive example, where $n$ is the size of the positive example.

Suzuki et al. [11, 12] have shown that the learnabilities of $OTT_\Lambda^1$ and $OTT_\Lambda$ in the framework of polynomial time inductive inference from positive data, where $|\Lambda| \geq 1$. Thus, we will study the learnability of finite unions of term trees in $OTT_\Lambda^1$ in the same framework.

## References


