

## On weak notion of $p$ -dividing

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### Abstract

I considered the restricted notions of weak dividing. In this note, I try to define a weak notion of  $p$ -dividing (thorn-dividing).

### 1. Preliminaries

We recall some definitions.

**Definition 1** Let  $\varphi(x_0, x_1, \dots, x_{n-1})$  be a formula and  $p(x)$  be a type. We denote the type  $\{\varphi(x_0, x_1, \dots, x_{n-1})\} \cup p(x_0) \cup p(x_1) \cup \dots \cup p(x_{n-1})$  by  $[p]^\varphi$ .

Let  $A \subset B$  and  $p(x) \in S(B)$ .

$p(x)$  *divides over*  $A$  if there is a formula  $\varphi(x, b) \in p(x)$  and an infinite sequence  $\{b_i : i < \omega\}$  with  $b \equiv b_i(A)$  such that  $\{\varphi(x, b_i) : i < \omega\}$  is  $k$ -inconsistent for some  $k < \omega$ .

$p(x)$  *weakly divides over*  $A$  if there is a formula  $\varphi(\bar{x}) \in L_n(A)$  such that  $[p[A]^\varphi]$  is consistent, while  $[p]^\varphi$  is inconsistent.

We can define weak dividing for formulas.

Let  $b \notin A$ .

$\psi(x, b)$  *weakly divides over*  $A$  if there is a formula  $\varphi(\bar{x}) \in L_n(A)$  and a realization  $a$  of  $\psi(x, b)$  such that  $[tp(a/A)]^\varphi$  is consistent, while  $[\psi(x, b)]^\varphi$  is inconsistent.

And we can consider weak forking.

$p(x)$  *weakly forks over*  $A$  if there is a  $q(x, y) \in S(A)$  such that  $p(x) \cup q(x, y)$  is consistent, and any completion  $r(x, y) \in S(B)$  of  $p(x) \cup q(x, y)$  weakly divides over  $A$ .

If we exchange the role between variables and parameters in the definition of weak dividing, we could define weak forking naturally.

In this note, we call such formula " $\varphi(\bar{x})$ " in the definition above the *witness formula* of weak dividing for the sake of convenience.

I introduce an example from [3].

**Example 2** Let  $T$  be the theory of an equivalence relation with two infinite classes of the language  $L = \{ \text{a binary relation } E(x, y) \}$ . And let  $\models \neg E(a, b)$ . Then the type  $\text{tp}(a/b)$  does not divide over  $\emptyset$ , while  $\text{tp}(a/b)$  weakly divides over  $\emptyset$  by the formula  $\neg E(x, y)$ .

I tried to divide witness formulas into some classes according to their properties ago. And I told about the next characterization at the RIMS meeting last year.

**Definition 3** Let  $A \subset B$  and  $p(x) \in S(B)$ .

$p(x)$   $\mathcal{M}$ -weakly divides over  $A$  if there is a formula  $\varphi(\bar{x}) \in L_n(A)$  and a Morley sequence  $I = \{a_i : i < n+1\}$  of  $p \upharpoonright A$  such that  $\models \varphi(a_0, a_1, \dots, a_{n-1})$ , while the type  $[p]^\varphi$  is inconsistent.

**Theorem 4** Let  $T$  be simple.

Then  $T$  is stable if and only if  $\mathcal{M}$ -weak dividing over models is symmetric.

## 2. Weak notion of $p$ -dividing

In recent years another variant of dividing, "thorn"-dividing has been characterized in rosy theory (see e.g. [4]). I tried to define weak notion of  $p$ -dividing (thorn-dividing). We recall some definitions first.

**Definition 5** Let  $A \subset B$  and  $p(x) \in S(B)$ .

$p(x)$  strongly divides over  $A$  if there is a formula  $\varphi(x, b) \in p(x)$  such that  $b \notin \text{acl}(A)$  and  $\{\varphi(x, b_i) : b_i \models \text{tp}(b/A)\}$  is  $k$ -inconsistent for some  $k < \omega$ .

$p(x)$   $\mathfrak{p}$ -divides over  $A$  if  $p(x)$  strongly divides over  $A_c$  for some parameter  $c$ .

$p(x)$   $\mathfrak{p}$ -forks over  $A$  if there is a formula  $\varphi(x, b) \in p(x)$  such that  $\varphi(x, b)$  implies a finite disjunction of formulas which  $\mathfrak{p}$ -divides over  $A$ .

Given a formula  $\varphi$ , a set  $\Delta$  of formulas in variables  $x, y$ , a set of formulas  $\Pi$  in variables  $y, z$ , and a number  $k$ , we define  $\mathfrak{p}(\varphi, \Delta, \Pi, k)$  (thorn-rank) inductively as follows :

- (1)  $\mathfrak{p}(\varphi, \Delta, \Pi, k) \geq 0, \infty, \lambda$  for limit ordinal  $\lambda$  is defined as usual.
- (2)  $\mathfrak{p}(\varphi, \Delta, \Pi, k) \geq \alpha + 1$  if and only if there is a  $\delta \in \Delta$ , some  $\pi(y, z) \in \Pi$  and parameters  $c$  such that

(a)  $\mathfrak{p}(\varphi \wedge \delta(x, a), \Delta, \Pi, k) \geq \alpha$  for infinitely many  $a \models \pi(y, c)$

(b)  $\{\delta(x, a)\}_{a \models \pi(y, c)}$  is  $k$ -inconsistent.

For a type  $p$ , we define  $\mathfrak{p}(p, \Delta, \Pi, k) = \min\{\mathfrak{p}(\varphi, \Delta, \Pi, k) \mid \varphi \in p\}$ .

A theory  $T$  is rosy if for any type  $p(x)$ , any finite sets of formulas  $\Delta$  and  $\Pi$ , and any finite  $k$ ,  $\mathfrak{p}(\varphi, \Delta, \Pi, k)$  is finite.

**Remark 6** (1) In rosy theories,  $\mathfrak{p}$ -forking satisfies the independence axioms.

(2) If  $a \models \varphi(x, b)$  and  $\varphi(x, b)$   $\mathfrak{p}$ -divides over  $C$  by the set  $\{b_i \models \theta(y, d)\}$ , then  $b \in \text{acl}(Cda) - \text{acl}(Cd)$ .

Weak notions of  $\mathfrak{p}$ -dividing could be defined in many ways. By the definition,  $\mathfrak{p}$ -dividing implies dividing. So we expect that weak  $\mathfrak{p}$ -dividing implies weak dividing.

**Definition 7** Let  $b \notin A$ .

$\psi(x, b)$  *weakly  $\mathfrak{p}$ -divides over  $A$*  if there is a formula  $\varphi(\bar{x}) = \exists y \bigwedge_{i < n} \theta(x_i, y) \in L_n(A)$  and a realization  $a$  of  $\psi(x, b)$  such that  $[\text{tp}(a/A)]^\varphi$  is consistent, while  $[\psi(x, b)]^\varphi$  is inconsistent.

We define weak  $\mathfrak{p}$ -dividing( $\mathfrak{p}$ -forking) for types just like weak dividing(forking).

We can check the next fact easily.

**Fact 8** *Let  $T$  be rosy. Then  $\mathfrak{p}$ -forking implies weak  $\mathfrak{p}$ -forking.*

### 3. Weak $\mathfrak{p}$ -dividing and NIP theories

**Definition 9** A formula  $\varphi(x, y)$  has the *independence property* if for every  $n < \omega$ , there are sequences  $a_l$  ( $l < n$ ) such that for every  $w \subset n$ ,  $\models (\exists x) \left[ \bigwedge_{l < n} \varphi(x, a_l)^{\text{if } (l \in w)} \right]$ .

A theory  $T$  is *NIP* if no formula  $\varphi(x, y)$  has the independence property.

Weak  $\mathfrak{p}$ -dividing is a kind of algebraic extension.

**Lemma 10** ( *$T$  is any theory.*)  $A \subset B$ .

*Then  $\text{tp}(a/B)$  does not weakly  $\mathfrak{p}$ -divide over  $A$  if and only if*

*for any  $n < \omega$ , any  $C$  and any extension  $q(x, C, A)$  of  $\text{tp}(a/A)$  over  $AC$ , if  $\bigcup_{i < n} q(x_i, C, A)$  is consistent, then  $\bigcup_{i < n} q(x_i, Z, A) \cup \bigcup_{i < n} r(x_i, Y, A)$  is consistent where  $\text{tp}(a/B) := r(x, B, A)$ .*

By the lemma above, we can prove the next fact.

**Proposition 11** *Let  $T$  be NIP and unstable.*

*Then weak  $\mathfrak{p}$ -dividing is not symmetric.*

## References

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