

# A note on strictly stable generic structures

Koichiro Ikeda \*

Faculty of Business Administration  
Hosei University

## Abstract

We show that there is a generic structure in a finite language such that the theory is strictly stable and not  $\omega$ -categorical, and has finite closures.

## 1 The class $\mathbf{K}$

It is assumed that the reader is familiar with the basics of generic structures. For details, see Baldwin-Shi [1] and Wagner [3].

Let  $R, S$  be binary relations with irreflexivity, symmetricity and  $R \cap S = \emptyset$ . Let  $L = \{R, S\}$ .

**Definition 1.1** Let  $\mathbf{K}_0$  be the class of finite  $L$ -structures  $A$  with the following properties:

1.  $A \models R(a, b)$  implies that  $a, b$  are not  $S$ -connected;
2. If  $A \models R(a, b) \wedge R(b, c)$ , then  $a, c$  are not  $S$ -connected;
3. If  $A \models R(a, b) \wedge R(b', c)$  and  $b, b'$  are  $S$ -connected, then  $a, c$  are not  $S$ -connected;
4.  $A$  has no  $S$ -cycles.

**Definition 1.2** Let  $A \in \mathbf{K}_0$ .

- For  $a, b \in A$ ,  $aEb$  means that  $a$  and  $b$  are  $S$ -connected.
- For  $a \in A$ , let  $a_E = a/E$ , and let  $A_E = \{a_E : a \in A\}$ .

---

\*The author is supported by Grants-in-Aid for Scientific Research (No. 17K05350).

- A binary relation  $R_E$  on  $A_E$  is defined as follows: for any  $a, b \in A$ ,  $A_E \models R_E(a_E, b_E)$  iff there are some  $a', b' \in A$  with  $a'Ea, b'Eb$  and  $A \models R(a', b')$ . By Definition 1.1, the structure  $A_E = (A_E, R_E)$  can be considered as an  $R$ -structure (or an  $R$ -graph) with irreflexivity and symmetricity.

**Notation 1.3** Let  $A \in \mathbf{K}_0$ .

- Let  $s(A)$  denote the number of the  $S$ -edges in  $A$ .
- Let  $x(A) = |A| - s(A)$ .
- Let  $r(A)$  denote the number of the  $R$ -edges in  $A$ .
- For  $\alpha$  with  $0 < \alpha \leq 1$ , let  $\delta(A) = x(A) - \alpha \cdot r(A)$ .

**Definition 1.4** Let  $A, B, C \in \mathbf{K}_0$ .

- Let  $\delta(B/A)$  denote  $\delta(BA) - \delta(A)$ .
- For  $A \subset B$ ,  $A$  is said to be closed in  $B$ , denoted by  $A \leq B$ , if  $\delta(X/A) \geq 0$  for any  $X \subset B - A$ .
- For  $A = B \cap C$ ,  $B$  and  $C$  are said to be free over  $A$ , denoted by  $B \perp_A C$ , if  $R^{B \cup C} = R^B \cup R^C$  and  $S^{B \cup C} = S^B \cup S^C$ .
- When  $B \perp_A C$ , we write  $B \oplus_A C$  for an  $L$ -structure  $B \cup C$ .

**Lemma 1.5**  $(\mathbf{K}_0, \leq)$  has the free amalgamation property, i.e., whenever  $A \leq B \in \mathbf{K}_0$ ,  $A \leq C \in \mathbf{K}_0$  and  $B \perp_A C$  then  $B \oplus_A C \in \mathbf{K}_0$ .

**Proof.** Let  $D = B \oplus_A C$ . We have to check that  $D$  satisfies conditions 1-4 in Definition 1.1. Here, for simplicity, we see condition 2 in Definition??. Take any  $a, b, c \in D$  with  $R(a, b) \wedge R(b, c)$ . If  $abc$  is contained in either  $B$  or  $C$ , then it is clear that  $a$  and  $c$  are not  $S$ -connected. So we can assume that  $a \in B - A, b \in A$  and  $c \in C - A$ . Suppose for a contradiction that  $a$  and  $c$  are  $S$ -connected. Then there is some  $d \in A$  with  $R(d, c)$ . So  $\delta(c/A) \leq 1 - (\alpha + 1) < 0$ , and hence  $A \not\leq C$ , a contradiction. Hence  $a$  and  $c$  are not  $S$ -connected.

**Remark 1.6** In [2], Hrushovski proved that there were an  $\alpha \in (0, 1)$  and a function  $f : \mathbb{N} \rightarrow \mathbb{R}$  such that

1.  $f(0) = 0, f(1) = 1$ ;
2.  $f$  is unbounded and convex;

3.  $f'(n) \leq \min\{r : r = \frac{p - q\alpha}{m} > 0, m \leq n \text{ and } m, p, q \in \omega\}$  for each  $n \in \omega$ .

**Definition 1.7** For a function  $f$  in Remark 1.6, let  $\mathbf{K} = \{A \in \mathbf{K}_0 : \delta(A') \geq f(x(A')) \text{ for any } A' \subset A\}$ .

**Lemma 1.8**  $(\mathbf{K}, \leq)$  has the free amalgamation property.

**Proof.** Let  $A, B, C \in \mathbf{K}$  be such that  $A \leq B, A \leq C$  and  $B \perp_A C$ . Let  $D = B \oplus_A C$ . We want to show that  $D \in \mathbf{K}$ . By Lemma 1.5, we have  $D \in \mathbf{K}_0$ . So it is enough to see that  $f(|D|) \leq \delta(D)$ . Without loss of generality, we can assume that  $\delta(C/A) \geq \delta(B/A)$ . By Remark??, we have  $\frac{\delta(B) - \delta(A)}{|B| - |A|} \geq f'(|B|)$ . On the other hand, since  $B \in \mathbf{K}$ , we have  $\delta(B) \geq f(|B|)$ . Hence we have  $\delta(D) \geq f(|D|)$ .

**Definition 1.9** • Let  $\overline{\mathbf{K}}$  denote the class of  $L$ -structure  $A$  satisfying  $A_0 \in \mathbf{K}$  for every finite  $A_0 \subset A$ .

- For  $A \subset B \in \overline{\mathbf{K}}$ ,  $A \leq B$  is defined by  $A \cap B_0 \leq B_0$  for any finite  $B_0 \subset B$ .
- For  $A \subset B \in \overline{\mathbf{K}}$ , we write  $\text{cl}_B(A) = \bigcap \{C : A \subset C \leq B\}$ .
- It can be checked that there exists a countable  $L$ -structure  $M$  satisfying
  1. if  $M \in \overline{\mathbf{K}}$ ;
  2. if  $A \leq B \in \mathbf{K}$  and  $A \leq M$ , then there exists a copy  $B'$  of  $B$  over  $A$  with  $B' \leq M$ ;
  3. if  $A \subset_{\text{fin}} M$ , then  $\text{cl}_M(A)$  is finite.

This  $M$  is called a  $(\mathbf{K}, \leq)$ -generic structure.

## 2 Theorem

In what follows, let  $M$  be the  $(\mathbf{K}, \leq)$ -generic structure,  $T = \text{Th}(M)$  and  $\mathcal{M}$  a big model of  $T$ .

**Lemma 2.1**  $T$  has finite closures, i.e., for any finite  $A \subset \mathcal{M}$ ,  $\text{cl}_{\mathcal{M}}(A)$  is finite.

**Proof.** For each  $t \in \mathbf{R}$ , let  $H_t = \{(x, y) : x, r \in \omega, y = x - \alpha r, f(x) \leq y \leq t\}$ . Since  $f$  is unbounded, each  $H_t$  is finite. Hence any  $A \subset_{\text{fin}} \mathcal{M}$  has finite closures.

**Lemma 2.2**  $T$  is not  $\omega$ -categorical.

**Proof.** Let  $a_0, a_1, \dots$  be vertices with the relations  $S(a_0, a_1), S(a_1, a_2), \dots$ . Since  $a_0 a_1 \dots \in \overline{\mathbf{K}}$ , we can assume that  $a_0 a_1 \dots \subset \mathcal{M}$ . It can be checked that  $\text{tp}(a_0 a_n) \neq \text{tp}(a_0 a_m)$  for each distinct  $m, n \in \omega$ . Then  $S_2(T)$  is infinite. Hence  $T$  is not  $\omega$ -categorical.

For  $A \subset_{\text{fin}} \mathcal{M}$  and  $n \in \omega$ ,  $A$  is said to be  $n$ -closed, if  $\delta(X/A) \geq 0$  for any  $X \subset \mathcal{M} - A$  with  $|X| \leq n$ .

**Notation 2.3** Let  $A \leq_{\text{fin}} \mathcal{M}$  and  $n \in \omega$ .

- $\text{cltp}_n(A) = \{X \cong A\} \cup \{X \text{ is } n\text{-closed}\}$
- $\text{cltp}(A) = \bigcup_{i \in \omega} \text{cltp}_i(A)$
- $E(A) = \{B \in \mathbf{K} : A \leq B\}$
- $E^+(A) = \{B \in E(A) : \text{there is a copy of } B \text{ over } A \text{ in } \mathcal{M}\}$
- $E^-(A) = E(A) - E^+(A)$
- $\text{ptp}(A) = \{\exists Y (XY \cong AB) : B \in E^+(A)\}$
- $\text{ntp}(A) = \{\neg \exists Y (XY \cong AB) : B \in E^-(A)\}$
- $\text{gtp}(A) = \text{cltp}(A) \cup \text{ptp}(A) \cup \text{ntp}(A)$
- $\text{gtp}_n(A) = \text{cltp}_n(A) \cup \text{ptp}(A) \cup \text{ntp}(A)$

**Definition 2.4** Let  $A \subset B \in \mathbf{K}_0$ . Then  $B_A$  is an  $L \cup \{R_E, S_E\}$ -structure with the following properties:

1. the universe is  $\{b_E : b \in B - A\} \cup A$ ;
2. the restriction of  $B$  on  $A$  is the  $L$ -structure  $A$ ;
3. for  $a \in A$  and  $b \in B - A$ ,  $B_A \models R_E(a, b_E)$  iff there is a  $b' \in B - A$  with  $b'Eb$  and  $B \models R(a, b')$ , and  $B_A \models R_E(b_E, a)$  iff there is a  $b' \in B - A$  with  $b'Eb$  and  $B \models R(b', a)$ ;
4. for  $a \in A$  and  $b \in B - A$ ,  $B_A \models S_E(a, b_E)$  iff there is a  $b' \in B - A$  with  $b'Eb$  and  $B \models S(a, b')$ , and  $B_A \models S_E(b_E, a)$  iff there is a  $b' \in B - A$  with  $b'Eb$  and  $B \models S(b', a)$ ;

5. for  $b, c \in B - A$ ,  $B_A \models R(b_E, c_E)$  iff there are  $b', c' \in B - A$  with  $b'Eb, c'Ec$  and  $B \models R(b', c')$ .

**Note 2.5** By the similar argument as in Definition 1.2, the structure  $B_A$  is canonically considered as an  $L$ -structure.

**Lemma 2.6** Let  $A \leq_{\text{fin}} \mathcal{M}$  and  $n \in \omega$ . Then  $\text{gtp}_n(A)$  is finitely generated.

**Proof.** Take a sequence  $(S_i)_{i \in \omega}$  of finite subsets of  $\text{gtp}_n(A)$  with  $S_0 \subset S_1 \subset \dots$  and  $\bigcup S_i = \text{gtp}_n(A)$ . For  $i \in \omega$ , let  $\sigma_i(X) = \bigwedge S_i$ . We can assume that  $\models \sigma_i(A')$  implies  $A' \cong A$ . Since  $f$  is unbounded,  $\mathcal{C}_i = \{C'_{A'} : M \models \sigma_i(A'), C' = \text{cl}_M(A')\}$  is finite. So there is some  $i_0 \in \omega$  such that  $\mathcal{C}_j = \mathcal{C}_{i_0}$  for every  $j > i_0$ . Hence  $S_{i_0}$  generates  $\text{gtp}_n(A)$ .

**Lemma 2.7** If  $\text{gtp}(A) = \text{gtp}(B)$  and  $A \leq C \leq_{\text{fin}} \mathcal{M}$ , then there is a  $D$  with  $\text{gtp}(AC) = \text{gtp}(BD)$ .

**Proof.** Let  $\Sigma(XY) = \text{gtp}(AC)$  and let  $\Sigma_n(XY) = \text{gtp}_n(AC)$  for  $n \in \omega$ . We want to show that  $\Sigma(BY)$  is consistent. To show this, it is enough to see that  $\Sigma_n(BY)$  is consistent for each  $n$ . On the other hand, by Lemma 2.6,  $\Sigma_n(XY)$  can be considered as some formula  $\sigma(XY)$ . So we want to show that  $\sigma(BY)$  has a realization. For this, we prove that  $\sigma(XY) \wedge \phi(X)$  has a realization for each  $\phi(X) \in \text{tp}(B)$ . Let  $\tau(X) = \sigma(XY)|_X$ . Note that  $\tau(X) \wedge \phi(X) \in \text{tp}(B)$  and  $\tau(X) \vdash \text{gtp}_n(A) = \text{gtp}_n(B)$ . Take  $B' \models \tau \wedge \phi$  in  $M$ . Take  $A'C' \models \sigma$  in  $M$  with  $A' \cup \text{cl}(A') \cong B' \cup \text{cl}(B')$ . Let  $DE$  be such that  $DE \cup \text{cl}(B') \cong C' \cup \text{cl}(C') \cup \text{cl}(A')$ . By genericity, we can assume that  $E \leq M$ . Then we have  $\models \sigma(B'D)$ , and hence  $\sigma(XY) \wedge \phi(X)$  has a realization.

**Corollary 2.8** Let  $A \leq_{\text{fin}} \mathcal{M}$ . Then  $\text{gtp}(A) \vdash \text{tp}(A)$ .

**Definition 2.9** Let  $A, B, C \subset \mathcal{M}$  with  $A = B \cap C$ . Then the notation  $B \downarrow_A^* C$  is defined as follows: for each  $n \in \omega$  and  $A^* B^* C^* \models \text{gtp}_n(ABC)$  in  $M$ ,

1.  $\text{cl}(B^*) \cap \text{cl}(C^*) = \text{cl}(A^*)$ ;
2.  $\text{cl}(B^*) \perp_{\text{cl}(A^*)} \text{cl}(C^*)$ .

**Lemma 2.10** Let  $A \leq B \leq \mathcal{M}, A \leq E \leq \mathcal{M}$  and  $E \downarrow_A^* B$ . Then  $\text{gtp}(E/A) \vdash \text{gtp}(E/B)$ .

**Proof.** For simplicity, we assume that  $A, B$  and  $E$  are finite. Take any  $E_1 \models \text{gtp}(E/A)$  with  $E_1 \downarrow_A^* B$  in  $M$ . Fix any  $n$ . Then there are realizations  $E^*A^*, E_1^*A^* \models \text{gtp}_n(EA)$  in  $M$  with  $\text{cl}(E^*) \cong_{\text{cl}(A^*)} \text{cl}(E_1^*)$ . Since  $E \downarrow_A^* B$  and  $E_1 \downarrow_A^* B$ , there is  $B^*A^* \models \text{gtp}_n(BA)$  with  $\text{cl}(E^*) \cong_{\text{cl}(B^*)} \text{cl}(E_1^*)$ . Hence  $E_1 \models \text{gtp}(E/B)$ .

**Lemma 2.11**  $T$  is strictly stable.

**Proof.** Let  $N \prec \mathcal{M}$  with  $|N| = \lambda$ . Take any  $e \in \mathcal{M} - N$ . Then there is a countable  $A \leq N$  with  $e \downarrow_A^* N$ . Let  $E = \text{cl}(eA)$ . We can assume that  $E \cap N = A$ . We want to show that  $\text{gtp}(E/A) \vdash \text{gtp}(E/N)$ . Take any  $E_1, E_2 \models \text{gtp}(E/A)$  with  $E_i \downarrow_A^* N$ . Take any countable  $N_0 \leq N$ . Take  $E_i^*A^* \subset M$  such that  $E_1^*A^*, E_2^*A^* \models \text{gtp}_n(EA)$  and  $\text{cl}(E_1^*A^*) \cong \text{cl}(E_2^*A^*)$ . Hence  $\text{gtp}(E_1/N) = \text{gtp}(E_2/N)$ . It follows that  $|S(N)| \leq 2^\omega \cdot \lambda^\omega = \lambda^\omega$ . Hence  $T$  is stable.

**Theorem 2.12** There is a generic structure  $M$  with the following properties:

1. the language is finite;
2.  $\text{Th}(M)$  is not  $\omega$ -categorical;
3.  $\text{Th}(M)$  has finite closures;
4.  $\text{Th}(M)$  is strictly stable.

## References

- [1] J. T. Baldwin and N. Shi, Stable generic structures. *Ann. Pure Appl. Log.* 79, 1-35 (1996)
- [2] E. Hrushovski, A stable  $\aleph_0$ -categorical pseudoplane. preprint (1988)
- [3] F. O. Wagner, Relational structures and dimensions. In: *Automorphisms of First-Order Structures*, 153-181, Clarendon Press, Oxford (1994)

Faculty of Business Administration  
Hosei University  
Tokyo 102-8160, Japan  
E-mail: ikeda@hosei.ac.jp