Introduction to Inter-universal Teichmüller Theory I

- An Approximate Statement of the Main Theorem -

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Introduction to IUT I

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Notation and Terminology

$$\mathcal{O}^{\boldsymbol{\mu}} \stackrel{\text{def}}{=} (\mathcal{O}^{\times})_{\text{tor}} \subseteq \mathcal{O}^{\times} \stackrel{\text{def}}{=} \{|z| = 1\}$$
$$\subseteq \mathcal{O}^{\rhd} \stackrel{\text{def}}{=} \{0 < |z| \le 1\} \subseteq \mathcal{O} \stackrel{\text{def}}{=} \{|z| \le 1\}$$

$\mathcal{O}^{ imes \mu} \stackrel{\mathrm{def}}{=} \mathcal{O}^{ imes} / \mathcal{O}^{\mu}$

an isomorph of $A \stackrel{\mathrm{def}}{\Leftrightarrow}$ an object which is isomorphic to A

 $R_+:$ the underlying additive module of a ring R

- F: a number field, i.e., $[F:\mathbb{Q}]<\infty$, s.t. $\sqrt{-1}\in F$
- $\mathbb{V}(-)$: the set of primes of (-)
- E: an elliptic curve over F which has

either good or split multiplicative reduction at $\forall v \in \mathbb{V}(F)$ $q_v \in \mathcal{O}_{F_v}^{\rhd}$: the q-parameter of E at $v \in \mathbb{V}(F)$ $q_E \stackrel{\text{def}}{=} (q_v)_{v \in \mathbb{V}(F)} \in \prod_{v \in \mathbb{V}(F)} \mathcal{O}_{F_v}^{\rhd}$ $\Rightarrow \deg q_E (= [F : \mathbb{Q}]^{-1} \log(\prod \sharp(\mathcal{O}_{F_v}/q_v \mathcal{O}_{F_v}))) \quad (\approx 6 \cdot \text{ht}_E)$

The Szpiro Conjecture for Elliptic Curves over Number Fields A certain upper bound of ht_E , i.e., $\deg q_E$ Suppose that the following (*) holds:

(*): $\exists N \geq 2$, $\exists C \geq 0$ s.t. $\deg q_E^N \leq \deg q_E + C$

Then since $\deg q_E^N = N \cdot \deg q_E$, one may conclude that

$$\deg q_E \leq \frac{C}{N-1}.$$

In order to establish (*), let us

- take two isomorphs ${}^{\dagger}\mathfrak{S}$, ${}^{\ddagger}\mathfrak{S}$ of (a part of) scheme theory,
- consider a "link" between these two isomorphs

$$\Theta_{\text{naive}}$$
: ${}^{\dagger}\mathfrak{S} \ni {}^{\dagger}q_E^N \mapsto {}^{\ddagger}q_E \in {}^{\ddagger}\mathfrak{S}$, and

 compare, via Θ_{naive}, the computation of deg of [†]q^N_E (in [†]S) with the computation of deg of [‡]q_E (in [‡]S).

$$\Theta_{\text{naive}} \colon {}^{\dagger}\mathfrak{S} \to {}^{\ddagger}\mathfrak{S} \colon {}^{\dagger}q_E^N \mapsto {}^{\ddagger}q_E$$

Very roughly speaking, the main theorem of IUT asserts that: Relative to such a link, the computation of $\deg^{\dagger}q_E^N$ is, up to mild indeterminacies, *compatible* with the computation of $\deg^{\ddagger}q_E$.

$$(\Rightarrow \deg q_E^N \stackrel{\mathsf{ind}.}{=} \deg q_E \Rightarrow (*) \Rightarrow \mathsf{the Szpiro Conjecture})$$

Terminology

- a(n) (arithmetic) holomorphic structure
 def def a (structure which determines a) ring structure
- a mono-analytic structure

 $\stackrel{\rm def}{\Leftrightarrow}$ an "underlying" ("non-holomorphic") structure of a hol. str.

 $\begin{array}{ll} \mbox{(e.g.:} & \mathbb{Q}_p, \ \pi_1^{\text{\'et}}(\mathbb{P}^1_{\mathbb{Q}_p} \setminus \{0,1,\infty\}) \text{: hol.; } & (\mathbb{Q}_p)_+, \ \mathbb{Q}_p^{\times}, \ G_{\mathbb{Q}_p} \text{: mono-an.)} \end{array}$

 $F_{\mathrm{mod}}\subseteq F$: the field of moduli of E

l: a prime number

 $K \stackrel{\text{def}}{=} F(E[l](\overline{F}))$

 $\underline{\mathbb{V}} \subseteq \mathbb{V}(K): \text{ the image of a splitting of } \mathbb{V}(K) \twoheadrightarrow \mathbb{V}_{\text{mod}} \stackrel{\text{def}}{=} \mathbb{V}(F_{\text{mod}})$ Suppose that F/F_{mod} and K/F_{mod} are *Galois*. $S \stackrel{\text{def}}{=} \left[\text{Spec } \mathcal{O}_K/\text{Gal}(K/F_{\text{mod}}) \right] \quad (\text{the stack-theoretic quotient})$ $\Rightarrow \quad \text{The arith. div. on } \mathcal{O}_F \text{ determined by } q_E \text{ can be descended to an arith. div. on } S, \text{ i.e., by considering the arith. div. on } S \text{ det'd by}$

$$\mathfrak{q} \stackrel{\text{def}}{=} (q_{\underline{v}} \stackrel{\text{def}}{=} q_{\underline{v}|_F} \in \mathcal{O}_{F_{\underline{v}|_F}}^{\rhd} \subseteq \mathcal{O}_{K_{\underline{v}}}^{\rhd})_{\underline{v} \in \underline{\mathbb{V}}} \in \prod_{\underline{v} \in \underline{\mathbb{V}}} \mathcal{O}_{K_{\underline{v}}}^{\rhd}.$$

Note that $\deg \mathfrak{q} \stackrel{\text{def}}{=} [F_{\text{mod}} : \mathbb{Q}]^{-1} \log(\prod_{\underline{v} \in \mathbb{V}} \sharp(\mathcal{O}_{K_{\underline{v}}}/q_{\underline{v}}\mathcal{O}_{K_{\underline{v}}})) = \deg q_E.$

Recall An arithmetic line bundle on \mathcal{O}_K

 $= \quad \text{a certain pair of a l.b. } \mathcal{L} \text{ on } \mathcal{O}_K \text{ and a metric on } \mathcal{L} \times_{\mathbb{Z}} \mathbb{C}$

 $1 \to \boldsymbol{\mu}(K) \to K^{\times} \stackrel{\text{ADiv}}{\to} \bigoplus_{w \in \mathbb{V}(K)} (K_w^{\times} / \mathcal{O}_{K_w}^{\times}) \to \operatorname{APic} \mathcal{O}_K \to 1$

Categories of Arithmetic Line Bundles on S

 $\mathcal{F}^{\circledast}_{\mathrm{mod}}:$ the Frobenioid of arithmetic line bundles on S

Module-theoretic Description

 $\mathcal{F}_{\mathfrak{mod}}^{\circledast}$: the Frobenioid of collections $\{a_{\underline{v}}\mathcal{O}_{K_{\underline{v}}}\}_{\underline{v}\in\mathbb{V}}$ s.t.

 $a_{\underline{v}} \in K_{\underline{v}}^{\times}, \quad a_{\underline{v}} \in \mathcal{O}_{K_{\underline{v}}}^{\times} \text{ for almost } \underline{v} \in \underline{\mathbb{V}}$

Multiplicative Description

 $\mathcal{F}^{\circledast}_{\mathrm{MOD}}: \text{ the Frobenioid of pairs } (T, \ \{t_{\underline{v}}\}_{\underline{v} \in \underline{\mathbb{V}}}) \text{ s.t.}$

$$T: \text{ an } F_{\mathrm{mod}}^{\times} \text{-torsor,} \quad t_{\underline{v}} \in T \times^{F_{\mathrm{mod}}^{\times}} K_{\underline{v}}^{\times} / \mathcal{O}_{K_{\underline{v}}}^{\times}$$

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 \Rightarrow The holomorphic structure of $F_{\rm mod}$ determines

 $\begin{array}{l} \mathcal{F}_{\mathrm{mod}}^{\circledast\mathbb{R}}, \ \mathcal{F}_{\mathrm{mod}}^{\circledast\mathbb{R}}, \ \mathcal{F}_{\mathrm{MOD}}^{\circledast\mathbb{R}}: \ \text{the resp. realifications of } \mathcal{F}_{\mathrm{mod}}^{\circledast}, \ \mathcal{F}_{\mathrm{mod}}^{\circledast}, \ \mathcal{F}_{\mathrm{MOD}}^{\circledast}, \\ \text{i.e., obtained by replacing } \bigoplus_{\underline{v}}(K_{\underline{v}}^{\times}/\mathcal{O}_{K_{\underline{v}}}^{\times}) \ \text{by } \bigoplus_{\underline{v}}((K_{\underline{v}}^{\times}/\mathcal{O}_{K_{\underline{v}}}^{\times})\otimes\mathbb{R}) \\ (\Rightarrow \ \text{The hol. str. of } F_{\mathrm{mod}} \ \text{determines } \mathcal{F}_{\mathrm{mod}}^{\circledast\mathbb{R}} \xrightarrow{\sim} \mathcal{F}_{\mathrm{mod}}^{\circledast\mathbb{R}} \xrightarrow{\sim} \mathcal{F}_{\mathrm{MOD}}^{\circledast\mathbb{R}}) \end{array}$

The multiplication by 1/N on $\bigoplus_{\underline{v}}((K_{\underline{v}}^{\times}/\mathcal{O}_{K_v}^{\times})\otimes \mathbb{R})$ determines

 $\Theta_{\mathrm{naive}}\colon \ ^{\dagger}\!\mathcal{F}_{\mathrm{MOD}}^{\circledast\mathbb{R}} \ \overset{\sim}{\longrightarrow} \ ^{\ddagger}\!\mathcal{F}_{\mathrm{MOD}}^{\circledast\mathbb{R}} \quad \text{which maps} \quad ^{\dagger}\!\mathfrak{q}^{N} \mapsto {^{\ddagger}\!\mathfrak{q}}.$

Remark

- $\bullet~\Theta_{naive}$ may be regarded as a "deformation of value groups".
- The link "Θ_{naive}" will be eventually established by means of nonarchimedean theta functions (cf. p.21).



Goal

Relative to a link such as Θ_{naive} , the computation of $\deg^{\dagger}\mathfrak{q}^{N}$ is, up to mild indeterminacies, *compatible* with the comp. of $\deg^{\ddagger}\mathfrak{q}$.

Note that \nexistsa ring automorphism of $K_{\underline{v}}$ s.t. $q_{\underline{v}}^N \mapsto q_{\underline{v}}$ (if $q_{\underline{v}} \neq 1$). Thus, Θ_{naive} cannot be compatible with the holomorphic structures, i.e., Θ_{naive} may be compatible with only certain mono-analytic str. (For instance, Θ_{naive} is compatible with the local Galois group $G_{\underline{v}} \stackrel{\text{def}}{=} \text{Gal}(\overline{F}_{\underline{v}}/K_{\underline{v}})$ for each finite $\underline{v} \in \underline{\mathbb{V}}$ — cf. Θ_{naive} "=" a deformation of value groups.) On the other hand:

Remark

The "degree computation" is, at least a priori, performed by means of the holomorphic structure under consideration.

Thus, in order to obtain a certain compatibility of the degree computations, we have to establish a "*multiradial representation*" of the degree computations whose coric data consist of suitable mono-analytic structures.

Multiradial Algorithm

Suppose that we are given

- a mathematical object R, i.e., a radial data,
- \bullet an "underlying" object C of R, i.e., a coric data, and
- a func'l algorithm Φ whose input data is (an isomorph of) R. Example
 - $\mathit{R}:$ the one-dimensional complex linear space $\mathbb C$
 - $C{:}$ the underlying two-dimensional real linear space $\mathbb{R}^{\oplus 2}$
 - R: the field \mathbb{Q}_p C: the underlying additive module $(\mathbb{Q}_p)_+$
 - R: the étale fundamental gp $\pi_1^{\text{ét}}(V)$ of a hyperbolic curve V/\mathbb{Q}_p C: the absolute Galois group $G_{\mathbb{Q}_p}$ of \mathbb{Q}_p

Roughly speaking, we shall say that the algorithm Φ is:

- coric if Φ depends only on C
- multiradial if Φ (is rel'd to R but) may be described in terms of C
- *uniradial* if Φ is not multiradial, i.e., essentially depends on R

If one starts with a coric data "C" and applies the alg'm Φ , then:

- uniradial ⇒ the output depends on the choice of a "spoke"
- multiradial \Rightarrow the output is unaffected by alterations in a "spoke"



(Tautological) Example $(R, C) \cong (\mathbb{C}, \mathbb{R}^{\oplus 2})$

- $\bullet~ \Phi(R) =$ the holomorphic structure on $R \Rightarrow$ uniradial
- $\Phi(R) =$ the real analytic structure on $R \Rightarrow$ coric
- $\Phi(R) = \text{the } \operatorname{GL}_2(\mathbb{R})\text{-orbit of the hol. str. on } R \Rightarrow \text{multiradial}$



Summary

- We want to obtain a certain compatibility of the degree computations relative to a link such as Θ_{naive} .
- Θ_{naive} cannot be compatible w/ the holomorphic str., i.e.,
 Θ_{naive} is compatible w/ only certain mono-an. str., e.g., G_v.
- On the other hand, the degree computation is, at least a priori, performed by means of the holomorphic structure.
- Thus, we have to establish a *multiradial representation* of the degree computations whose coric data are suitable mono-an. str.

An Approximate Statement of the Main Theorem of IUT (tentative) ∃A suitable multiradial algorithm whose output data consist of the following three objects ∽ mild indeterminacies

•
$$\{(\mathcal{O}_{K_{\underline{v}}})_*\}_{\underline{v}\in\mathbb{V}}$$
 (* = + if \underline{v} is finite, * = \emptyset if \underline{v} is infinite)

•
$$\mathfrak{q}^N \curvearrowright \prod_{\underline{v} \in \underline{\mathbb{V}}} (\mathcal{O}_{K_{\underline{v}}})_*$$

•
$$F_{\text{mod}} \curvearrowright \prod_{\underline{v} \in \underline{\mathbb{V}}} ((K_{\underline{v}})_* \text{``via} (\mathcal{O}_{K_{\underline{v}}})_* \text{''})$$

Moreover, this algorithm is *compatible* with

$$\Theta_{\text{naive}} \colon {}^{\dagger}\mathcal{F}_{\text{MOD}}^{\circledast \mathbb{R}} \xrightarrow{\sim} {}^{\ddagger}\mathcal{F}_{\text{MOD}}^{\circledast \mathbb{R}}; \, {}^{\dagger}\mathfrak{q}^N \mapsto {}^{\ddagger}\mathfrak{q}.$$

Main Theorem of IUT \Rightarrow Szpiro Conjecture

$$\begin{array}{l} \Theta_{\text{naive}} \colon {}^{\dagger}\!\mathcal{F}_{\text{MOD}}^{\circledast\mathbb{R}} \xrightarrow{\sim} {}^{\ddagger}\!\mathcal{F}_{\text{MOD}}^{\circledast\mathbb{R}}; \quad {}^{\dagger}\!\mathfrak{q}^N \mapsto {}^{\ddagger}\!\mathfrak{q} \\ \xrightarrow{\mathsf{Th'm}} \end{array}$$

$$\begin{pmatrix} \{(^{\dagger}\mathcal{O}_{K_{\underline{v}}})_{*}\}_{\underline{v}\in\underline{\mathbb{V}}} \\ {}^{\dagger}\mathfrak{q}^{N} \curvearrowright \prod (^{\dagger}\mathcal{O}_{K_{\underline{v}}})_{*} \\ {}^{\dagger}F_{\mathrm{mod}} \curvearrowright \prod (^{\dagger}K_{\underline{v}})_{*} \end{pmatrix} \xrightarrow{\mathrm{ind.}} \begin{pmatrix} \{(^{\dagger}\mathcal{O}_{K_{\underline{v}}})_{*}\}_{\underline{v}\in\underline{\mathbb{V}}} \\ {}^{\ddagger}\mathfrak{q}^{N} \curvearrowright \prod (^{\ddagger}\mathcal{O}_{K_{\underline{v}}})_{*} \\ {}^{\ddagger}F_{\mathrm{mod}} \curvearrowright \prod (^{\ddagger}K_{\underline{v}})_{*} \end{pmatrix}$$

•
$$c_{\text{mod}} : {}^{\dagger}\mathcal{F}_{\text{mod}}^{\circledast} \xrightarrow{\simeq} {}^{\ddagger}\mathcal{F}_{\text{mod}}^{\circledast}$$
 which maps $\{{}^{\dagger}q_{\underline{v}}^{N\dagger}\mathcal{O}_{K_{\underline{v}}}\} \mapsto \{{}^{\ddagger}q_{\underline{v}}^{N\ddagger}\mathcal{O}_{K_{\underline{v}}}\}$
• $c_{\Box} : {}^{\Box}\mathcal{F}_{\text{mod}}^{\circledast} \xrightarrow{\sim} {}^{\Box}\mathcal{F}_{\text{MOD}}^{\circledast}$ which maps $\{{}^{\Box}q_{\underline{v}}^{(N)}{}^{\Box}\mathcal{O}_{K_{\underline{v}}}\} \mapsto {}^{\Box}\mathfrak{q}^{(N)}$

which are *compatible* with $\Theta_{\rm naive}$, i.e.,

fit into a diagram that is commutative, up to mild indeterminacies

$$\begin{array}{cccc} ^{\dagger}\mathcal{F}_{\mathfrak{mod}}^{\circledast\mathbb{R}} & \stackrel{c_{\mathfrak{mod}}}{\longrightarrow} ~^{\ddagger}\mathcal{F}_{\mathfrak{mod}}^{\circledast\mathbb{R}} \\ & c_{\dagger} & & \downarrow c_{\ddagger} \\ & ^{\dagger}\mathcal{F}_{\mathrm{MOD}}^{\circledast\mathbb{R}} & \stackrel{\Theta_{\mathrm{naive}}}{\longrightarrow} ~^{\ddagger}\mathcal{F}_{\mathrm{MOD}}^{\circledast\mathbb{R}} \\ & ^{\dagger}\mathcal{F}_{\mathrm{MOD}}^{\otimes\mathbb{R}} \stackrel{\Theta_{\mathrm{naive}}}{\longrightarrow} ~^{\ddagger}\mathcal{F}_{\mathrm{MOD}}^{\otimes\mathbb{R}} \\ & & \uparrow \mathcal{F}_{\mathrm{MOD}}^{\otimes\mathbb{R}} \\ & & \uparrow \mathcal{F}_{\mathrm{MOD}}^{\otimes\mathbb{R}} \stackrel{\otimes}{\longrightarrow} ~^{\ddagger}\mathcal{F}_{\mathrm{MOD}}^{\otimes\mathbb{R}} \\ & & \uparrow \mathcal{F}_{\mathrm{MOD}}^{\otimes\mathbb{R}} \stackrel{\otimes}{\longrightarrow} ~^{\ddagger}\mathcal{F}_{\mathrm{MOD}}^{\otimes\mathbb{R}} \\ & & \uparrow \mathcal{F}_{\mathrm{MOD}}^{\otimes\mathbb{R}} \stackrel{\otimes}{\longrightarrow} ~^{\ddagger}\mathcal{F}_{\mathrm{MOD}}^{\otimes\mathbb{R}} \\ & & \uparrow \mathfrak{q}^{N} \stackrel{\otimes}{\leftarrow} ~^{\ddagger}\mathfrak{q} \\ & & \uparrow \mathfrak{q}^{N} \stackrel{\otimes}{\leftarrow} ~^{\ddagger}\mathfrak{q} \\ & & \uparrow \mathfrak{q}^{N} \stackrel{\otimes}{\leftarrow} & \overset{\dagger}{=}\mathfrak{q} \\ & & \Rightarrow ~^{\dagger}\mathfrak{q}_{\underline{v}}^{\dagger}\mathcal{O}_{K_{\underline{v}}} \\ & & \Rightarrow ~^{\dagger}\mathfrak{q} \stackrel{\otimes}{\longrightarrow} C_{\mathrm{noises}} ~^{\dagger}\mathfrak{q} \\ & & \Rightarrow ~^{\dagger}\mathfrak{q} \stackrel{\otimes}{\longrightarrow} C_{\mathrm{noises}} \\ & & & \bullet ~^{\dagger}\mathfrak{q} \stackrel{\otimes}{\longrightarrow} C_{\mathrm{noises}} \\ & & & \bullet ~^{\dagger}\mathfrak{q} \stackrel{\otimes}{\longrightarrow} C_{\mathrm{noises}} \\ & & & \bullet ~^{\dagger}\mathfrak{q} \stackrel{\otimes}{\longrightarrow} C_{\mathrm{noises}} \\ & & & & \bullet ~^{\dagger}\mathfrak{q} \stackrel{\otimes}{\longrightarrow} C_{\mathrm{noises}} \\ & & & & & \bullet ~^{\dagger}\mathfrak{q} \stackrel{\otimes}{\longrightarrow} C_{\mathrm{noises}} \\ & & & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & & \\ &$$

Introduction to IUT I

An Approximate Statement of the Main Theorem of IUT (tentative) ∃A suitable multiradial algorithm whose output data consist of the following three objects ∽ mild indeterminacies

•
$$\{(\mathcal{O}_{K_{\underline{v}}})_*\}_{\underline{v}\in\mathbb{V}}$$
 (* = + if \underline{v} is finite, * = \emptyset if \underline{v} is infinite)

•
$$\mathfrak{q}^N \curvearrowright \prod_{\underline{v} \in \underline{\mathbb{V}}} (\mathcal{O}_{K_{\underline{v}}})_*$$

•
$$F_{\text{mod}} \curvearrowright \prod_{\underline{v} \in \underline{\mathbb{V}}} ((K_{\underline{v}})_* \text{``via} (\mathcal{O}_{K_{\underline{v}}})_* \text{''})$$

Moreover, this algorithm is *compatible* with

$$\Theta_{\text{naive}} \colon {}^{\dagger}\mathcal{F}_{\text{MOD}}^{\circledast \mathbb{R}} \xrightarrow{\sim} {}^{\ddagger}\mathcal{F}_{\text{MOD}}^{\circledast \mathbb{R}}; \, {}^{\dagger}\mathfrak{q}^N \mapsto {}^{\ddagger}\mathfrak{q}.$$

Recall $(\underline{v} \in \underline{\mathbb{V}}$: fin.; $p \stackrel{\text{def}}{=} p_{\underline{v}})$ $(\mathcal{O}_{K_{\underline{v}}})_+ \in \text{output}, G_{\underline{v}}$: coric Unfortunately, it is known that:

 $\exists a \ func'l \ (w.r.t. \ open \ injections) \ alg'm \ for \ rec. \ "(\mathcal{O}_{K_{\underline{v}}})_+" \ from \ "G_{\underline{v}}".$ $\mathcal{I}_{K_{\underline{v}}} \stackrel{\text{def}}{=} \frac{1}{2p} \text{Im}(\mathcal{O}_{K_{\underline{v}}}^{\times} \to \mathcal{O}_{K_{\underline{v}}}^{\times} \otimes_{\mathbb{Z}} \mathbb{Q} \xrightarrow{\log_p} (K_{\underline{v}})_+): \ \text{the } \textit{log-shell of } K_{\underline{v}}$

• $\mathcal{I}_{K_{\underline{v}}}$: a finitely generated free \mathbb{Z}_p -module

- $(\mathcal{O}_{K_{\underline{v}}})_+$, $\log_p(\mathcal{O}_{K_{\underline{v}}}^{\times}) \subseteq \mathcal{I}_{K_{\underline{v}}} \subseteq \mathcal{I}_{K_{\underline{v}}} \otimes_{\mathbb{Z}} \mathbb{Q} = (K_{\underline{v}})_+$
- $[\mathcal{I}_{K_{\underline{v}}}:(\mathcal{O}_{K_{\underline{v}}})_+]$ (< ∞) can be comp'd by the top. gp str. of $G_{\underline{v}}$

•
$$G_{\underline{v}} \stackrel{\exists \text{func'l}}{\Rightarrow}_{\text{algorithm}}$$
 an isomorph of $\mathcal{I}_{K_{\underline{v}}}$

Thus, " $\{(\mathcal{O}_{K_{\underline{v}}})_*\}_{\underline{v}\in\underline{\mathbb{V}}}$ " in Th'm should be replaced by $\{\mathcal{I}_{\underline{v}} \stackrel{\text{def}}{=} \mathcal{I}_{K_{\underline{v}}}\}_{\underline{v}\in\underline{\mathbb{V}}}$ (where the log-shell at an infinite $\underline{v}\in\underline{\mathbb{V}}\stackrel{\text{def}}{=} \pi \cdot \mathcal{O}_{K_{\underline{v}}}$). Recall \mathfrak{q}^N , $F_{\mathrm{mod}} \in \mathsf{output}$

Unfortunately, by various indeterminacies arising from the operation of "passing from holom'c str. to mono-anal'c str.", it is difficult to obtain multiradial representations of q^N , F_{mod} themselves directly. To establish a mul'l alg'm of the desired type, we rep. multiradially a *suitable function* whose special value is q_v^N or an $\in F_{mod}$.

- " $q_{\underline{v}}^{N}$ " will be represented as a special value of a (multiradially represented) *theta function*.
- "F_{mod}" will be represented as a set of special values of (multiradially represented) κ-coric functions.

An Approximate Statement of the Main Theorem of IUT For a "general E/F",

 \exists a suitable multiradial algorithm whose output data consist of the following three objects \checkmark mild indeterminacies

- the collection of log-shells $\{\mathcal{I}_{\underline{v}}\}_{\underline{v}\in\underline{\mathbb{V}}}$
- the theta values $(= \{\mathfrak{q}^{j^2/2l}\}_{1 \leq j \leq l^*} \stackrel{\text{def}}{=} \frac{1}{2}) \curvearrowright \prod_{\underline{v} \in \underline{\mathbb{V}}} \mathcal{I}_{\underline{v}}$
- F_{mod} via κ -coric functions $\curvearrowright \prod_{\underline{v} \in \underline{\mathbb{V}}} ((K_{\underline{v}})_+ \text{"via } \mathcal{I}_{\underline{v}})$

Moreover, this alg'm is *compatible* w/ the Θ -link (more precisely, $\Theta_{LGP}^{\times \mu}$ -link) "[†] $\mathcal{F}_{MOD}^{\otimes \mathbb{R}} \xrightarrow{\sim} {}^{\ddagger}\mathcal{F}_{MOD}^{\otimes \mathbb{R}}$ "; "[†]theta values $\mapsto {}^{\ddagger}\mathfrak{q}^{1/2l}$ ".

Fundamental Strategy

 \Box is, for instance, a log-shell, a theta function, or a $\kappa\text{-coric}$ function.

- Start with a usual/existing \Box (i.e., a *Frobenius-like* \Box).
- Construct *links* by means of such Frobenius-like objects.
- Take an *étale-like* object closely related to \Box (e.g., " $\pi_1^{\text{temp}}(\underline{X}_{\underline{v}})$ " for a theta function — cf. II and III).
- Give a multiradial mono-anabelian algorithm of reconstructing □
 from the étale-like object, i.e., construct a suitable étale-like □.
- Establish "multiradial Kummer-detachment" of □, i.e.,
 - a suitable Kummer isomorphism "Frob.-like $\Box \xrightarrow{\sim}$ étale-like \Box " .