## A note on lowness for Robinson theories

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#### **Abstract**

We show following two theorems. Theorem A: for thick simple existentially universal domain, the equality of Lascar strong types is definable by an existential type. Theorem B: for thick low existentially universal domain, Lascar strong types equal strong types. Theorem A is already proved by Ben-Yaacov [2].

#### 1 Preliminaries

**Definition 1.1** We say that an L-structure M is  $\kappa$ -existentially universal domain (e.u.domain) if

- if  $\Sigma(x)$  is a partial existential type over A ( $|A| < \kappa$ ) which is finitely satisfiable in M, then  $\Sigma$  is satisfiable in M, and
- for  $|A|, |B| < \kappa$ , and  $f: A \to B$ : a bijection such that  $\exp(a) \subset \exp(f(a))$  for all tuples a from A, f extends to an automorphism of M.

**Remark 1.1** An e.u.domain M is an existentially closed model for the universal theory of M,  $\operatorname{Th}(M)_{\forall}$ .

Let  $\mathcal{M}$  be a  $\kappa$ -e.u.domain for a enoughly big cardinal  $\kappa$ . Put  $T = \operatorname{Th}_{\forall}(\mathcal{M})$ .  $M, N, \ldots$  denote existentially closed models of  $T, a, b, \ldots$  denote finite tuples in  $\mathcal{M}$ , and  $A, B, \ldots$  denote small subsets of  $\mathcal{M}$ .

**Definition 1.2** Let  $\Sigma(x, B)$  be an existential type over B.

1. We sat that  $\Sigma(x, B)$  divides over A if there exists an existentially indiscernible sequence  $(B_i : i < \omega)$  over A with  $B_0 = B$  such that  $\bigcup_{i < \omega} \Sigma(x, B_i)$  is not realized in  $\mathcal{M}$ .

- 2. We say that  $\Sigma(x)$  forks over A if there exists a small set of dividing (/A) existential formulas  $\Psi$  (with parameters) such that  $\mathcal{M} \models \Sigma \to \bigvee \Psi$ .
- **Remark 1.2** If  $\Sigma(x)$  divides over A, then there is an existential formula  $\varphi(x)$  such that  $\Sigma \vdash \varphi(x)$  and  $\varphi(x)$  divides over A.
  - It is not known whether that if  $\Sigma$  forks over A, then there are an existential formula  $\theta$  where  $\Sigma \vdash \theta$  and dividing (/A) existential formulas  $\psi_1, \ldots, \psi_n$  such that  $\mathcal{M} \models \theta \to \bigvee_{i=1}^n \psi_i$ .

**Definition 1.3** We say that  $\mathcal{M}$  is simple if for all  $a \in \mathcal{M}$ ,  $A \subset \mathcal{M}$ , there exists  $B \subset A$  with  $|B| \leq |T| + \aleph_0$  such that  $\exp(a/A)$  does not fork over B.

**Fact 1.1** [3] Suppose that  $\mathcal{M}$  is simple. Then,  $\Sigma$  forks over A if and only if  $\Sigma$  divides over A.

- **Definition 1.4** 1. We say that lstp(a) = lstp(b) if for any bounded  $\emptyset$ invariant equivalence relation E(x, y), E(a, b) holds.
  - 2. We say that  $d(a,b) \leq 1$  if there is an existentially indiscernible sequence I such that  $a,b \in I$ .
  - 3. We say that  $d(a, b) \leq n$  if there exist  $a_0, \ldots, a_n$  with  $a_0 = a, a_n = b$  such that  $d(a_i, a_{i+1}) \leq 1$  for any i < n.
  - 4. We say that  $d(a,b) < \omega$  if  $d(a,b) \le n$  for some  $n < \omega$ .

Fact 1.2  $/3/\operatorname{lstp}(a) = \operatorname{lstp}(b)$  if and only if  $d(a, b) < \omega$ .

**Fact 1.3** [3] If  $(a_i : i < \lambda)$  is an enoughly long sequence and  $A \subset \mathcal{M}$ , then there is an existentially indiscernible sequence  $(b_i : i < \omega)$  such that for any  $n < \omega$ , there are  $i_0 < \cdots < i_{n-1} < \lambda$  such that  $\exp(b_0, \ldots, b_{n-1}/A) = \exp(a_{i_0}, \ldots, a_{i_{n-1}}/A)$ .

**Fact 1.4** [3] Suppose that  $\mathcal{M}$  is simple. Then, for all  $a, A \subset B$ , there exists a' such that

- lstp(a'/A) = lstp(a/A) and
- etp(a'/B) does not fork over A.

We write  $a \downarrow b$  to mean that etp(a/b) does not fork over  $\emptyset$ .

Fact 1.5 (Independence theorem for simple e.u.domain, [3]) Suppose that  $\mathcal{M}$  is simple and  $a_1, a_2, b_1, b_2$  satisfy the following:

- $lstp(a_1) = lstp(a_2),$
- $a_1 \downarrow b_1$ ,  $a_2 \downarrow b_2$ ,  $b_1 \downarrow b_2$ .

Then, there exists a such that

- $a \models \exp(a_1/b_1) \cup \exp(a_2/b_2)$
- $a \downarrow b_1b_2$ .

### 2 Proof of Thorem A

In this section, we prove Tehorem A. For simplicity, we show over  $\emptyset$ .

**Definition 2.1** We say that  $\mathcal{M}$  is thick if " $d(x,y) \leq 1$ " is definable by an existential type. If  $\mathcal{M}$  is thick, then we assume that  $q_1(x,y)$  defines " $d(x,y) \leq 1$ ".

**Lemma 2.1** Suppose that  $\mathcal{M}$  is thick. Then, " $d(x,y) \leq 2$ " is definable by an existential type.

*Proof:* It is defined by  $\{\exists z \varphi(x, z) \land \varphi(z, y) | \varphi(x, y) \in q_1(x, y)\}.$ 

**Lemma 2.2** Suppose that  $\mathcal{M}$  is thick and simple. Then, the following are equivalent:

- 1. lstp(a) = lstp(b)
- 2.  $d(a,b) \leq 2$
- 3.  $q_1(x,a) \cup q_1(x,b)$  does not fork over  $\emptyset$

*Proof:*  $(3 \to 2 \to 1)$  is trivial.  $(1 \to 2)$  Let c be a tuple such that lstp(c) = lstp(a) = lstp(b) and  $c \downarrow ab$ . Take a' such that etp(a'a) = etp(ac). Then lstp(a') = lstp(a) and  $a' \downarrow a$ . So, by independence theorem, we can get  $a_2$  such that  $a_2 \models etp(a/c) \cup etp(a'/a)$  and  $a_2 \downarrow ac$ .

Iterating this, we can get a sequence  $(a_i : i < \omega)$  such that  $\exp(a_i a_j) = \exp(ac)$  for each  $j < i < \omega$ . By compactness and Fact 1.3, we can assume this sequence is existentially indiscernible. So, we get existentially indiscernible sequences I, J such that  $a, c \in I$  and  $b, c \in J$ .

**Theorem A** [2] Suppose that  $\mathcal{M}$  is thick and simple. Then, "lstp(x) = lstp(y)" is definable by an existential type.

*Proof:* By above lemmas.

#### 3 Proof of Theorem B

In this section, we prove Tehorem B. Again for simplicity, we show over  $\emptyset$ .

**Definition 3.1** We say that stp(a) = stp(b) if for any definable (by an existential formula over  $\emptyset$ ) finite equivalence relation E(x, y), E(a, b) holds.

- **Definition 3.2** 1. Let  $\varphi(x,y)$  be an existential formula. An existential formula  $\psi(y_0,\ldots,y_{k-1})$  where  $\mathrm{lh}(y_i)=\mathrm{lh}(y)$  for each i< k is said to be a k-inconsistency witness for  $\varphi$  if  $\mathcal{M}\models\forall y_0\ldots y_{k-1}(\psi(y_0,\ldots,y_{k-1})\to \neg\exists x\bigwedge_{i\leq k}\varphi(x,y_i)).$ 
  - 2. Let  $\Sigma(x)$  be an existential type and  $\varphi(x,y)$  be an existential formula.
    - We say that  $D(\Sigma, \varphi) \geq 0$  if  $\Sigma$  is satisfiable.
    - We say that  $D(\Sigma, \varphi) \geq n + 1$  if there is a natural number k, a k-inconsistency witness  $\psi$ , and an existentially indiscernible sequence  $(b_i : i < \omega)$  such that  $D(\Sigma(x) \cup \{\varphi(x, b_i)\}, \varphi) \geq n$  for each  $i < \omega$  and  $\mathcal{M} \models \psi(b_{i_0}, \dots b_{i_{k-1}})$  for all  $i_0, \dots, i_{k-1} < \omega$ .
  - 3. We say that  $\mathcal{M}$  is low if
    - $\mathcal{M}$  is simple and
    - $D(x = x, \varphi) < \omega$  for any existentiall formula  $\varphi$ .

**Lemma 3.1** Suppose that  $\mathcal{M}$  is thick and low. Then,

- 1.  $\{a: \varphi(x,a) \text{ divides over } \emptyset\}$  is definable by an existential type.
- 2.  $\{(a,b): \varphi(x,a) \land \varphi(x,b) \text{ does not divide over } \emptyset\}$  is definable by an existential type if it is restricted to  $(p \otimes p)^{\mathcal{M}} = \{(a,b): a,b \models p,a \downarrow b\}$ . So, it is definable by an existential universal formula if it is restricted to  $(p \otimes p)^{\mathcal{M}}$
- *Proof:* (1) Note that by lowness, for any  $\varphi(x,y)$  there is an existentiall formula  $\psi$  such that for all a, if  $\varphi(x,a)$  divides over  $\emptyset$ , then  $\varphi$  devides by an existentially indiscernible sequence in which any k-elements satisfies  $\psi$ .
  - (2) For  $a, b \models p$  where  $a \downarrow b$ , the following are equivalent:

- 1.  $\varphi(x,a) \wedge \varphi(x,b)$  does not divide over  $\emptyset$
- 2. there exist  $a^*$  and  $b^*$  such that
  - $\mathcal{M} \models \varphi(a^*, a)$  and  $a^* \downarrow a$ ;
  - $\mathcal{M} \models \varphi(b^*, b)$  and  $b^* \downarrow b$ ;
  - $lstp(a^*) = lstp(b^*)$

By Theorem A, "lstp( $a^*$ ) = lstp( $b^*$ )" is expressible by an existential type. " $a^* \downarrow a$ " is expressible by " $D(\text{etp}(a/a^*), \varphi, \psi) \geq D(p, \varphi, \psi)$ " for any  $\varphi, \psi$ .

We sat that  $E_{p(x),\varphi(x,y)}(b,c)$  if for all  $a \models p$  with  $a \stackrel{\downarrow}{\smile} bc$ ,  $\varphi(x,a) \land \varphi(x,b)$  does not divide over  $\emptyset$  if and only if  $\varphi(x,a) \land \varphi(x,c)$  does not divide over  $\emptyset$ .

**Lemma 3.2** Suppose that  $\mathcal{M}$  is thick and low. For any  $a \models p$  where  $\varphi(x, a)$  does not divide over  $\emptyset$ ,  $E_{p(x),\varphi(x,y)}$  is a definable (by an existential formula) finite equivalence relation on  $(p^2)^{\mathcal{M}}$ .

Proof: We can check that  $E_{p,\varphi}$  is a bounded equivalence relation boundedness is by " $lstp(x) = lstp(y) \Rightarrow E_{p,\varphi}(x,y)$ ". On the other hand, by the above lemma  $\neg E_{p,\varphi}$  is definable by an existential type. So,  $E_{p,\varphi}$  is a finite equivalence relation. Let  $a_1, \ldots, a_n$  be representations of classes. Then  $\bigcup \{\neg E(x,a_i) : i \leq n\}$  is not satisfiable. For simplicity, we assume n=3. There exists an eixistential formula  $\varphi(x,y)$  such that

- 1.  $\neg E(x, a_i) \vdash \varphi(x, a_i)$  for each  $i \leq 3$
- 2.  $\mathcal{M} \models \neg \exists x \varphi(x, a_1) \land \varphi(x, a_2) \land \varphi(x, a_3)$ .

Put  $\psi(x,y) = \neg \varphi(x,y)$ . Note that  $\mathcal{M} \models \forall x(\psi(x,a_1) \leftrightarrow \varphi(x,a_2) \land \varphi(x,a_3))$ . So,  $\psi(x,a_1)$  is also existential. By a symmetric argument,  $\psi(x,a_2), \psi(x,a_3)$  are all existential. Then we have

$$E(x,y) \leftrightarrow \bigwedge_{i < 3} (\psi(x,a_i) \leftrightarrow \psi(y,a_i)).$$

We can omit parameters  $a_i$ 's because this does not depend on a choice of representations and  $\psi(x, a_i)$  is existential universal.

**Theorem B** Suppose that  $\mathcal{M}$  is thick and low. Then, stp = lstp

*Proof:* If stp(a) = stp(b), then by the above lemma  $a, b \models E_{p,\varphi}$  for any  $\varphi$ . Take c such that lstp(c) = lstp(a) and  $c \downarrow ab$ . Then,  $q_1(x,a) \cup q_1(x,c)$  does not divide by Lemma 3. Then,  $q_1(x,b) \cup q_1(x,c)$  does not divide by  $E_{p,\varphi}(a,b)$ . Again by Lemma 3, we have lstp(b) = lstp(c).

# References

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