ON 3-AMPLENESS IN ROSY THEORIES

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ABSTRACT. In rosy theories we introduce a geometric notion of independence, strong non-3-ampleness, and we show that strong non-3-ampleness implies non-3-ampleness, and non-2-ampleness(=CM-triviality) implies strong non-3-ampleness.

1. INTRODUCTION

There is a simple characterization of CM-triviality. By using the characterization, we could show that any rosy CM-trivial theory has weak canonical bases, and CM-triviality in the real sort implies geometric elimination of imaginaries [Y]. We want to know whether these results can be extended in case of non-3-ampleness or not, so our first motivation is to find a simple characterization of non-3-ampleness like in case of CM-triviality. This paper is organized as follows. In the second section, we review the definiton of CM-triviality(=non-2-ampleness) and non-n-ampleness for each $n < \omega$ in rosy theories. In the third section, trying to find a simple charactrization, we offer another geometric notion (we call it strong non-3-ampleness). We show that strong non-3-ampleness implies non-3-ampleness, and non-2-ampleness(=CM-triviality) implies strong non-3-ampleness. But, for now there are no examples of non-3-ample and 2-ample structures. We also raise up some problems on non-3-ampleness.

Our notation is standard. Let T be a rosy theory. (i.e. having a good independence relation \downarrow) We work in \mathcal{M}^{eq} , the eq-structure, consisting of imaginary elements, where \mathcal{M} is a sufficiently saturated model of T. $\bar{a}, \bar{b}, \ldots \subset_{\omega} \mathcal{M}$ denote finite sequences in \mathcal{M}^{eq} . A, B, \ldots denote small subsets of \mathcal{M}^{eq} and AB denotes $A \cup B$. For $a \in \mathcal{M}^{eq}$ and $A \subset \mathcal{M}^{eq}$, we write $a \in \operatorname{acl}^{eq}(A)$ if the orbit of a by automorphisms fixing A pointwise is finite. In rosy theories [A], we have that $a \downarrow_b c$ implies $\operatorname{acl}^{eq}(ab) \cap \operatorname{acl}^{eq}(bc) = \operatorname{acl}^{eq}(b)$.

2. REVIEW OF ROSY CM-TRIVIALITY AND NON-n-AMPLENESS

CM-triviality is a geometric notion of the nonforking independence relation. In 1988, it is introduced by Hrushovski where he disproves Zilber's conjecture on strongly minimal sets [H]. CM-triviality forbids a point-line-plane incident system. Hrushovski also offered three characterizations of CM-triviality in stable theories. The following is the simplest characterization for rosy CM-triviality.

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Definition 2.1. A rosy theory T is CM-trivial, if $\bar{a} \downarrow_{A \cap \operatorname{acl}^{\operatorname{eq}}(\bar{a},B)} B$ holds for any $\bar{a}, A, B \subset \mathcal{M}^{\operatorname{eq}}$ such that $\bar{a} \downarrow_A B$ and A, B are algebraically closed.

The Weak Caninical Base $\operatorname{wcb}(\bar{a}/B)$ of $\operatorname{tp}(\bar{a}/B)$ has the following properties, where B is an algebraically closed subset of $\mathcal{M}^{\operatorname{eq}}$:

- $\bar{a} \perp_{\operatorname{wcb}(\bar{a}/B)} B$
- $\operatorname{wcb}(\bar{a}/B)$ is algebraically closed.
- $\bar{a} \downarrow_A B \Rightarrow \operatorname{wcb}(\bar{a}/B) \subseteq \operatorname{acl}^{eq}(A) \subseteq \mathcal{M}^{eq}$

The weak canocal base is the smallest algebraically closed subset C of B such that $\bar{a} \downarrow_C B$. As in [P2], rosy theories do not necessarily have weak canonical bases. But any rosy CM-trivial theory has weak canonical bases, so we have the following characterization [Y].

Fact 2.2. Let T be rosy. The following are equivalent.

- (1) T is CM-trivial.
- (2) T has weak canonical bases and $\operatorname{wcb}(\bar{a}/A) \subseteq \operatorname{wcb}(\bar{a}/B)$ holds for any $\bar{a}, A, B \subset \mathcal{M}^{\operatorname{eq}}$ such that $\operatorname{acl}^{\operatorname{eq}}(\bar{a}, A) \cap B = A$ with $A = \operatorname{acl}^{\operatorname{eq}}(A)$ and $B = \operatorname{acl}^{\operatorname{eq}}(B)$.

We use the following notations to briefly write the definition of n-ampleness.

$$a \wedge b := \operatorname{acl}^{\operatorname{eq}}(a) \cap \operatorname{acl}^{\operatorname{eq}}(b)$$

 $a_{< i} := a_0, a_1, \dots, a_{i-1}$
 $a_{< 0} := \emptyset$

The following is proved in [P1].

Remark 2.4. (1) one-basedness $(a \downarrow_{a \wedge b} b \text{ for any } a, b)$ is equivalent to non-1-ampleness: $ca_0 \wedge ca_1 = \operatorname{acl^{eq}}(c) \ (a_1 \downarrow_{a_0} a_0) \text{ implies } a_1 \downarrow_c a_0$.

- (2) CM-triviality $(a \downarrow_b c \Rightarrow a \downarrow_{b \land ac} c)$ is equivalent to non-2-ampleness: $ca_0 \land ca_1 = \operatorname{acl^{eq}}(c), ca_0 a_1 \land ca_0 a_2 = \operatorname{acl^{eq}}(ca_0)$ $a_2 \downarrow_{a_1} a_0, \ a_1 \downarrow_{a_0} a_0 \text{ imply } a_2 \downarrow_c a_0.$
- (3) Non-n-ampleness implies non-(n+1)-ampleness for each $n < \omega$.

3. STRONG NON-3-AMPLENESS

The following remark is a non-3-ample's version of Fact 2.2 under assuming the existence of weak canonical bases.

Remark 3.2. If T has weak canonical bases, then the following are equivalent.

(1) T is not 3-ample.

(2) $\operatorname{wcb}(a_3/ca_0) \subseteq \operatorname{acl^{eq}}(\operatorname{wcb}(a_3/ca_0a_1a_2)c)$ holds for any $a_0, a_1, a_2, a_3, c \in \mathcal{M}^{eq}$ such that $a_0c \wedge a_1c = \operatorname{acl^{eq}}(c), a_0a_1c \wedge a_0a_2c = \operatorname{acl^{eq}}(a_0c), a_0a_1a_2c \wedge a_0a_1a_3c = \operatorname{acl^{eq}}(a_0a_1c) \ a_3 \downarrow_{a_2c} a_0a_1, a_2 \downarrow_{a_1c} a_0$

Proof. $(1) \Rightarrow (2)$: Clear.

(1) \Leftarrow (2): We may assume $c = \emptyset$. As $a_3 \downarrow_{a_2} a_0 a_1$ and $\operatorname{wcb}(a_3/a_0) \subseteq \operatorname{wcb}(a_3/a_0 a_1 a_2)$, we have $\operatorname{wcb}(a_3/a_0) \subseteq a_0 \wedge a_2$. On the other hand, as $a_2 \downarrow_{a_1} a_0$, we have $a_0 \wedge a_2 \subseteq a_0 a_1 \wedge a_1 a_2 \subseteq \operatorname{acl}^{\operatorname{eq}}(a_1)$. As $a_0 \wedge a_1 = \operatorname{acl}^{\operatorname{eq}}(\emptyset)$, we have $\operatorname{wcb}(a_3/a_0) \subseteq a_0 \wedge a_2 \subseteq a_0 \wedge a_1 = \operatorname{acl}^{\operatorname{eq}}(\emptyset)$. \square

Now we consider the following notion.

Definition 3.3. We say that T is strongly non-3-ample, if $a_3 \downarrow_{a_0 a_2 \wedge a_1 a_2 \wedge a_0 a_1 a_3} a_0$ holds for any a_0, a_1, a_2, a_3 such that $a_3 \downarrow_{a_2} a_0 a_1, a_2 \downarrow_{a_1} a_0$

The definition of non-3-ampleness has three condition algebraically closed set and two conditions on independency. On the other hand, the definition of strong non-3-ampleness has only one condition on algebraically closed set and two conditions on independency, so it is simpler than that of non-3-ampleness.

Proposition 3.4. Strong non-3-ampleness implies non-3-ampleness.

Proof. Suppose that $a_0 \wedge a_1 = \operatorname{acl^{eq}}(\emptyset), a_0 a_1 \wedge a_0 a_2 = \operatorname{acl^{eq}}(a_0), a_0 a_1 a_2 \wedge a_0 a_1 a_3 = \operatorname{acl^{eq}}(a_0 a_1), a_3 \downarrow_{a_2} a_0 a_1, a_2 \downarrow_{a_1} a_0$, and let $b := a_0 a_2 \wedge a_1 a_2 \wedge a_0 a_1 a_3$.

We need to show $b = \operatorname{acl}^{eq}(\emptyset)$.

Claim 1 $b \subseteq \operatorname{acl}^{eq}(a_1)$: As $a_2 \downarrow_{a_1} a_0$, we have $a_0 a_1 \wedge a_1 a_2 = \operatorname{acl}^{eq}(a_1)$. Then we have

$$b = a_0 a_2 \wedge a_1 a_2 \wedge a_0 a_1 a_3$$

$$\subseteq a_0 a_1 a_2 \wedge a_1 a_2 \wedge a_0 a_1 a_3$$

$$= (a_0 a_1 a_2 \wedge a_0 a_1 a_3) \wedge a_1 a_2$$

$$= a_0 a_1 \wedge a_1 a_2 = \operatorname{acl}^{eq}(a_1)$$

Claim 2 $b \subseteq \operatorname{acl}^{eq}(a_0)$: As $b \subseteq \operatorname{acl}^{eq}(a_1)$,

$$b \subseteq a_1 \wedge a_0 a_2 \wedge a_1 a_2 \wedge a_0 a_1 a_3$$

$$\subseteq a_1 \wedge a_0 a_2 \subseteq a_0 a_1 \wedge a_0 a_2 = \operatorname{acl}^{eq}(a_0)$$

By two claims, we have $b \subseteq a_0 \wedge a_1 = \operatorname{acl}^{eq}(\emptyset)$.

Remark 3.5. Assume that $\operatorname{acl^{eq}}(A(B \wedge C)) = AB \wedge AC$ for any $A, B, C \subseteq \mathcal{M}^{eq}$. (We usually have that $\operatorname{acl^{eq}}(A(B \wedge C)) \subseteq AB \wedge AC$.) Then non-3-ampleness coincides with strong non-3-ampleness.

Proof. Let $b = a_0 a_2 \wedge a_1 a_2 \wedge a_0 a_1 a_3$. As $b \subseteq a_0 a_2 \wedge a_1 a_2$, we have

$$a_3 \underset{a_2b}{\bigcup} a_0a_1, a_2 \underset{a_1b}{\bigcup} a_0.$$

So we need to show

- $(1) \ a_0b \wedge a_1b = \operatorname{acl}^{\operatorname{eq}}(b)$
- (2) $a_0a_1b \wedge a_0a_2b = \operatorname{acl}^{eq}(a_0b)$
- (3) $a_0a_1a_2b \wedge a_0a_1a_3b = \operatorname{acl}^{eq}(a_0a_1b)$

The proof of $a_0b \wedge a_1b = \operatorname{acl}^{eq}(b)$:

$$a_0b \wedge a_1b \subseteq (a_0a_2 \wedge a_0a_1a_2 \wedge a_0a_1a_3) \wedge (a_0a_1a_2 \wedge a_1a_2 \wedge a_0a_1a_3)$$

 $\subseteq a_0a_2 \wedge a_1a_2 \wedge a_0a_1a_3 = \operatorname{acl}^{\operatorname{eq}}(b)$

The proof of $a_0a_1b \wedge a_0a_2b = \operatorname{acl}^{\operatorname{eq}}(a_0b)$: We use our assumption in the last equation.

$$a_0a_1b \wedge a_0a_2b \subseteq (a_0a_1a_2 \wedge a_0a_1a_3) \wedge (a_0a_2 \wedge a_0a_1a_2 \wedge a_0a_1a_2a_3)$$

 $\subseteq a_0a_2 \wedge a_0a_1a_3 = \operatorname{acl}^{eq}(a_0b)$

The proof of $a_0a_1a_2b \wedge a_0a_1a_3b = \operatorname{acl^{eq}}(a_0a_1b)$: We also use our assumption in the last equation.

$$a_0a_1a_2b \wedge a_0a_1a_3b \subseteq (a_0a_1a_2 \wedge a_0a_1a_2a_3) \wedge (a_0a_1a_2a_3 \wedge a_0a_1a_3)$$

 $\subseteq a_0a_1a_2 \wedge a_0a_1a_3 = \operatorname{acl}^{eq}(a_0a_1b)$

Remark 3.6. Non-2-ampleness (=CM-triviality) implies strong non-3-ampleness.

Proof. We will show that $a_3 \downarrow_{a_2} a_0 a_1$ implies $a_3 \downarrow_{a_0 a_2 \wedge a_1 a_2 \wedge a_0 a_1 a_3} a_0 a_1$. Put $A = \operatorname{acl^{eq}}(a_0 a_2)$ and $B = \operatorname{acl^{eq}}(a_1 a_2)$. Clearly we have $a_3 \downarrow_{a_2} AB$. As $\operatorname{acl^{eq}}(a_2) \subseteq A \cap B \subseteq AB$, we have $a_3 \downarrow_{A \cap B} AB$. In particular, we have $a_3 \downarrow_{A \cap B} a_0 a_1$. By CM-triviality, we get $a_3 \downarrow_{a_0 a_2 \wedge a_1 a_2 \wedge a_0 a_1 a_3} a_0 a_1$, as desired.

We have the following implications. non-1-ampleness(\Leftrightarrow one-basedness) \Rightarrow non-2-ampleness(\Leftrightarrow CM-triviality) \Rightarrow strong non-3-ampleness \Rightarrow non-4-ampleness \Rightarrow ···

In [E] an *n*-ample (relational) structure M_n is constructed for each $n < \omega$, but it is unknown whether M_n is not (n+1)-ample. For now, *n*-ample and non-(n+1)-ample structure is not discoverd for each $n \ge 2$. (Generic relational structures are 1-ample and non-2-ample.)

Problem 3.7. It is shown that free group is 2-ample [P3]. Is free group non-3-ample? (We need the characterization of non-forking in the free group to check non-3-ampleness.)

Problem 3.8. Does non-3-ample theory have weak canonical bases? (I think No.) We need to check Adler's criterion [A]: $a \downarrow_B C$, $a \downarrow_C B \Rightarrow a \downarrow_{B \cap C} BC$ for any a, B, C such that B and C are algebraically closed subsets of \mathcal{M}^{eq} .

Problem 3.9. Is strong non-3-ampleness with weak canonical bases equivalent to CM-triviality?

Let $T=\operatorname{Th}(\mathbb{R},+,<,\pi|_{(-1,1)}(*))$, where $\pi|_{(-1,1)}(x):=\pi\cdot x$ for $x\in(-1,1)$. T is an o-minimal theory with elimination of imaginaries. As T does not have weak canonical bases, T is 2-ample. And T does not interpret fields by [LP][PS], so it is possible that T is not n-ample for some $n<\omega$

Problem 3.10. Is T non-3-ample? (We need the charactrization of non-forking in T.)

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