TOPOLOGY OF THE SPACE ON WHICH CELLULAR AUTOMATA WORKS

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Abstract. We investigate:

- (A) Topology of the product of discrete topological spaces and
- (B) Homeomorphism of two topologies (T_p) and (T_m) of cell space $C = S^{\mathbb{Z}^n}$ of a cellular automaton $\mathcal{A} = (\mathbb{Z}, S, N, f)$.

1. Product of Discrete Topolpgical Spaces

Definition.

1

 $\circ X_i$: discrete topological space for i in I

 \mathcal{O}_i : the system of open sets of X_i ,

 $\circ \ \ X = \prod_{i \in I} X_i \ : \ \mbox{the product of} \ X_i \ \mbox{of weak topology}$

 \mathcal{O} : the system of open sets of X,

and so

 $\circ \quad \mathcal{O}_0 = \{ \prod_{j \in J} O_j \ \times \prod_{i \in I \setminus J} X_i \mid O_j \in \mathcal{O}_j, \ J \subseteq I, \ |J| < \infty \}$

: an open basis of \mathcal{O} .

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Now we state our theorem.

Theorem A. For X_i 's descrete the following are equivalent:

- (a) $X = \prod_{i \in I} X_i$ is discrete.
- (b) $|\{i \in I \mid |X_i| \ge 2\}| < \infty$, that is, $|X_i| = 1$ for almost all i in I.

Corollary. The product of discrete topological spaces is discrete if and only if it is homeomorphic to a product of a finite number of discrete topological spaces. In particular, if S is a discrete topological space containing at least two elements, then its infinite product is not discrete.

Proof.

(a)
$$\Rightarrow X \ni {}^\forall x$$
: open, (since X is discrete)
$$\Rightarrow X \ni {}^\forall x = \cup_{\lambda \in \Lambda} O_{\lambda} \quad \text{for some} \quad O_{\lambda}\text{'s} \in \mathcal{O}_{0},$$
 (for \mathcal{O}_{0} is an open basis of \mathcal{O})
$$= O_{\lambda} \qquad \text{for any } \lambda \quad \text{in} \quad \Lambda,$$
 (since $|\{x\}| = 1$)

$$\Rightarrow \mathbf{setting} \\ O_{\lambda} = \prod_{j \in J} x_j \ \times \ \prod_{i \in I \setminus J} X_i \ \mathbf{for \ some} \ J, \ |J| < \infty \\ \mathbf{and} \\ x \ = \ \prod_{j \in J} x_j \ \times \ \prod_{i \in I \setminus J} x_i,$$

we have

$$\Rightarrow$$
 $X_i = x_i$ for i in $I \setminus J$ with $|J| < \infty$, \Rightarrow (b).

(b) \Rightarrow For

$$J = \{i \in I \mid |X_i| \ge 2\}$$

we have

$$|J| < \infty$$

 \Rightarrow Choose

$$X \ni {}^{\forall} x = \prod_{j \in J} x_j \times \prod_{i \in I \setminus J} x_i, \quad x_i \in X_i, \quad x_j \in X_j,$$

Then by the definition of J

$$\forall i \in I \setminus J, \quad x_i = X_i$$

and so

$$X\ni {}^\forall x=\prod_{j\in J}x_j \ \times \prod_{i\in I\setminus J}X_i\in \mathcal{O}_0, \ x_j\in X_j$$

$$\Rightarrow \ \textbf{(a)}$$
 Q.E.D.

2. The topology of the configuration set C of a cellular automaton $\mathcal A$

$$\circ~\mathcal{A}=(\mathbb{Z},S,N,f)~$$
 : a cellular automaton
$$\mathbb{Z}=\{0,\pm 1,\pm 2,\cdots\}~\text{: rational integers,}$$

$$\circ \ \mathcal{Z}^n = \mathbb{Z} \times \mathbb{Z} \times \cdots \times \mathbb{Z}$$
: cell space

$$\circ \ S = \{s_1, s_2, \cdots, s_m\} :$$
 states

$$\circ f: S^l \to S$$
: local map

Further we define

$$\circ \ C = S^{\mathbb{Z}^n} = \operatorname{Map}(\mathbb{Z}^n, S) :$$
 configurations

where we give

 \circ S: the discrete topology

Then $C = S^{\mathbb{Z}^n}$ has two topologies (\mathbf{T}_p) and (\mathbf{T}_m) following :

$$(\mathbf{T}_m)$$

$$\forall c \in C = S^{\mathbb{Z}^n},$$

$$\forall \epsilon = 2^{-\lambda}, \quad \lambda \in \mathbb{N}$$

define

$$U_{\epsilon}(c) = \{c' \in C \mid d(c,c') < \epsilon \}$$

: the ϵ -neighbourhood of c

where

$$d(c,c') = 2^{-\min\{\delta(0,i) \mid c(i) \neq c'(i), i \in I\}} \quad for \quad c,c' \quad in \quad C,$$

$$\delta(0,i) = \sqrt{{i_1}^2 + {i_2}^2 + \dots + {i_n}^2} \quad \textbf{for} \quad i = (i_1,i_2,\dots,i_n) \quad \textbf{in} \quad \mathbb{Z}^n. .$$

 (\mathbf{T}_p)

$$\forall c \in C = S^{\mathbb{Z}^n},$$

$$\forall J \subseteq \mathbb{Z}^n, \quad |J| < \infty$$

define

$$V_J(c) = (\prod_{j \in J} c(j)) \times S^{\mathbb{Z}^n \setminus J}$$

: the J-neighbourhood of c

Then, since S is discrete, we have $c_j \in \mathcal{O}_j$ for any $j \in J$. This enable us to take

$$\{V_J(c) \mid c \in C, J \subseteq \mathbb{Z}^n, |J| < \infty\}$$

as O_0 an open basis.

Then, we have:

Theorem B.

$$(\mathbf{T}_p) \simeq (\mathbf{T}_m)$$
: homeomorphic

Proof. (a) We show $(\mathbf{T}_p) \leq (\mathbf{T}_m)$, i.e., $\exists U_{\epsilon}(c) \subseteq \forall V_J(c)$

For

$$\forall V_J(c) = \prod_{j \in J} c(j) \times S^{\mathbb{Z}^n \setminus J}$$

with

$$\forall c \in C = S^{\mathbb{Z}^n},$$

$$\forall J \subseteq S^{\mathbb{Z}^n} \quad \text{and} \quad |J| < \infty$$

choose $\lambda \in \mathbb{N}$

such that (A)
$$\forall j \in J, \quad \delta(0,j) < \lambda$$

and set

$${}^\forall \epsilon = 2^{-\lambda}$$

Then,

$$c' \in U_{\epsilon}(c) \implies d(c,c') < \epsilon = 2^{-\lambda}, \text{ where}$$

$$d(c,c') = 2^{-\min\{\delta(0,i) \mid c(i) \neq c'(i)\}} \text{ for } c,c' \text{ in } C,$$

$$\Rightarrow \lambda < \min\{\delta(0,i) \mid c(i) \neq c'(i)\}$$

$$\Rightarrow \text{ By (A)}$$
 if $j \in J$, we have $\delta(0,j) < \lambda$ and so $c(j) = c'(j)$

$$\Rightarrow c' \in V_J(c) = \prod_{j \in J} c(j) \times S^{\mathbb{Z}^n \setminus J}$$

Thus,

$$U_{\epsilon}(c) \subseteq V_J(c)$$
 and so $(\mathbf{T}_p) \leq (\mathbf{T}_m)$.

(b) Next we show $(\mathbf{T}_m) \leq (\mathbf{T}_p)$, i.e. $, \exists V_J(c) \subseteq \forall U_{\epsilon}(c)$.

For

$$\forall U_{\epsilon}(c)$$

with

$$c \in C = S^{\mathbb{Z}^n},$$
 $\epsilon = 2^{-\lambda}, \quad \lambda \in \mathbb{N}$

let

(B)
$$J = \{j \in I \mid \delta(0, j) \le \lambda\}$$

Then, since

$$V_J(c) = \prod_{j \in J} c(j) \times S^{\mathbb{Z}^n \setminus J},$$

we have

$$c' \in V_J(c) \implies {}^\forall j \in J, \quad c(j) = c'(j)$$

$$\Rightarrow \quad \text{by (B)} \quad \text{if} \quad \delta(0,j) \le \lambda \text{ , we have } j \in J$$

$$\quad \text{and so } c(j) = c'(j)$$

$$\Rightarrow \quad \lambda < \min\{\delta(0,j) \mid c(j) \ne c'(j)\}$$

$$\Rightarrow \quad 2^{-\min\{\delta(0,j) \mid c(j) \ne c'(j)\}} < 2^{-\lambda} = \epsilon$$

$$\Rightarrow \quad d(c,c') < \epsilon$$

$$\Rightarrow \quad c' \in U_{\epsilon}(c)$$

Thus

$$V_J(c)\subseteq U_\epsilon(c) \quad ext{ and so } \quad (\mathbf{T}_m) \ \le (\mathbf{T}_p).$$
 Q.E.D.

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