

On Reynolds-Hurkens-Coquand paradox and T -algebras

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Abstract

Recently, Coquand [2] presented a variant on the paradox, as a variation of Reynolds “paradox” [5] in terms of weakly initial T -algebras. We first formalize the notion of the powerful universe U in λHOL such as $\Gamma_U := \{U : \square, in : T(U) \rightarrow U, out : U \rightarrow T(U)\}$ with a functor T where $T(U) := \text{Pow}(\text{Pow}(U))$ and $\text{Pow}(U) := U \rightarrow *$. Then rather than a retraction such that $out \circ in$ is the identity on $T(U)$, we suppose logical equivalence with kind $*$ as follows:

- $Ret_1 := \Pi X : T(U). \Pi p : \text{Pow}(U). (out \circ in X p) \leftrightarrow X p$,
- $Ret_2 := \Pi X : T(U). \Pi p : \text{Pow}(U). (T \delta X p) \leftrightarrow X p$,
- $Weak_1 := \Pi X : T(U). \Pi p : \text{Pow}(U). (out \circ in X p) \leftrightarrow T \delta X p$, and
- $Weak_2 := \Pi X : T(U). \Pi p : \text{Pow}(U). (out \circ in X p) \leftrightarrow X (p \circ \delta)$,

where $\delta := (in \circ out)$. Now, λHOL can prove that $\Gamma_U \vdash Weak_1 \leftrightarrow Weak_2$; $\Gamma_U \vdash Ret_1 \wedge Ret_2 \rightarrow Weak$; $\Gamma_U \vdash Ret_1 \wedge Weak \rightarrow Ret_2$; $\Gamma_U \vdash Ret_2 \wedge Weak \rightarrow Ret_1$; $\Gamma_U \vdash \neg Ret_1$; and $\Gamma_U \vdash \neg Weak$.

Next, in λU^- , we show a dual one with respect to the weakly initial T -algebra [2], i.e., a weakly final T -coalgebra as another instance of the powerful universe.

1. Introduction

Let λHOL (Church 1940) be a PTS (Pure Type System) with the specification [1], consisting of $* : \square, \square : \Delta$ for axioms; and $(*, *)$, $(\square, *)$, (\square, \square) for rules. That is, λHOL is Girard’s system F_ω with the axiom $\square : \Delta$, which is a consistent type system and can introduce variables of kinds, e.g., $X : \square$. It is well-known that in System F, weakly initial T -algebras and final T -coalgebras are impredicatively represented as the categorical data types.

Let λU^- be λHOL together with the rule (Δ, \square) , which can introduce polymorphic kinds like polymorphic types in system F. The use of the rule makes it possible to define powerful universes. Moreover, we work with the rules (\square, \square) and (Δ, \square) to represent categorical data types, just like system F of $(*, *)$ and $(\square, *)$.

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2. Encoding powerful universe in λHOL

As usual, define logical connectives for types and formulae with the kind $*$. Also define power sets and data types as kinds with the super kind \square .

1. Bottom $\perp := \Pi X : *. X$ with kind $*$ by $(\square, *)$.
2. Negation $\neg := \lambda X : *. X \rightarrow \perp$ with kind $(* \rightarrow *)$ by $(\square, *)$, (\square, \square) .
3. Conjunction $\wedge := \lambda A, B : *. \Pi X : *. (A \rightarrow B \rightarrow X) \rightarrow X$ with kind $(* \rightarrow * \rightarrow *)$ by (\square, \square) .
Write $A \leftrightarrow B$ for $(A \rightarrow B) \wedge (B \rightarrow A)$ with *pair*, *fst*, *snd* where $A, B : *$.
4. For $X : \square$, power set $\text{Pow}(X) := X \rightarrow *$ with super kind \square by (\square, \square) .
Write $\mathsf{T}(X)$ for $\text{Pow}(\text{Pow}(X))$ where $X : \square$.
5. Composition $g \circ f := \lambda x : X. g(fx)$
 $X, Y, Z : \square, f : (X \rightarrow Y), g : (Y \rightarrow Z) \vdash \lambda x : X. g(fx) : X \rightarrow Z$ by (\square, \square) .
6. Contravariant functor $\text{mon}_{X \rightarrow Y}(f) := \lambda q : \text{Pow}(Y). q \circ f$
 $X, Y : \square, f : X \rightarrow Y \vdash \lambda q : \text{Pow}(Y). q \circ f : \text{Pow}(Y) \rightarrow \text{Pow}(X)$ by (\square, \square) .
7. Covariant functor $\mathsf{T}_{X \rightarrow Y}(f) := \text{mon}_{\text{Pow}(Y) \rightarrow \text{Pow}(X)}(\text{mon}_{X \rightarrow Y} f)$
 $X, Y : \square, f : X \rightarrow Y \vdash \lambda F : \mathsf{T}(X). F \circ (\text{mon}_{X \rightarrow Y} f) : \mathsf{T}(X) \rightarrow \mathsf{T}(Y)$ by (\square, \square) .
8. Powerful universe $U : \square$ and a well-formed context Γ by $\square : \Delta$ and (\square, \square) :
 $\Gamma_U := \{U : \square, \text{in} : (\mathsf{T}(U) \rightarrow U), \text{out} : (U \rightarrow \mathsf{T}(U))\}$
 $\delta := \text{in} \circ \text{out}$
9. Retraction by $(\square, *)$:
 $\text{Ret1} := (\Pi X : \mathsf{T}(U). \Pi p : \text{Pow}(U). (\text{out} \circ \text{in}) X p \leftrightarrow X p)$
 $\Gamma_{\text{ret1}} := \Gamma_U \cup \{\text{ret1} : (\Pi X : \mathsf{T}(U). \Pi p : \text{Pow}(U). (\text{out} \circ \text{in}) X p \leftrightarrow X p)\}$
 $\alpha_e X p := \text{fst}(\text{ret1 } X p) : (\text{out} \circ \text{in}) X p \rightarrow X p$
 $\beta_e X p := \text{snd}(\text{ret1 } X p) : X p \rightarrow (\text{out} \circ \text{in}) X p$
 $\text{Ret2} := (\Pi X : \mathsf{T}(U). \Pi p : \text{Pow}(U). \mathsf{T} \delta X p \leftrightarrow X p)$
 $\Gamma_{\text{ret2}} := \Gamma_U \cup \{\text{ret2} : (\Pi X : \mathsf{T}(U). \Pi p : \text{Pow}(U). \mathsf{T} \delta X p \leftrightarrow X p)\}$
10. U is weakly initial (final) by $(\square, *)$:
 $\Gamma_U \vdash \Pi X : \mathsf{T}(U). \Pi p : \text{Pow}(U). (\text{out} \circ \text{in}) X p \leftrightarrow \mathsf{T}(\text{in} \circ \text{out}) X p : *$ by $(\square, *)$.
 $\text{Weak1} := \Pi X : \mathsf{T}(U). \Pi p : \text{Pow}(U). (\text{out} \circ \text{in}) X p \leftrightarrow \mathsf{T}(\delta) X p$
 $\text{Weak2} := (\Pi X : \mathsf{T}(U). \Pi p : \text{Pow}(U). (\text{out} \circ \text{in}) X p \leftrightarrow X (p \circ \delta))$
 $\Gamma_{\text{weak1}} := \Gamma_U \cup \{w1 : (\Pi X : \mathsf{T}(U). \Pi p : \text{Pow}(U). (\text{out} \circ \text{in}) X p \leftrightarrow \mathsf{T}(\delta) X p)\}$
 $\Gamma_{\text{weak2}} := \Gamma_U \cup \{w2 : (\Pi X : \mathsf{T}(U). \Pi p : \text{Pow}(U). (\text{out} \circ \text{in}) X p \leftrightarrow X (p \circ \delta))\}$
 $\alpha_w X p := \text{fst}(\text{weak2 } X p) : (\text{out} \circ \text{in}) X p \rightarrow X (p \circ \delta)$
 $\beta_w X p := \text{snd}(\text{weak2 } X p) : X (p \circ \delta) \rightarrow (\text{out} \circ \text{in}) X p$

Lemma 1 *Let $X, Y, Z: \square, f: X \rightarrow Y, g: Z \rightarrow X, h: W \rightarrow Z, p: \text{Pow}(Y), P: \mathbb{T}(X)$. Then we have the following elementary facts under $=_\beta$.*

1. $(f \circ g) \circ h =_\beta f \circ (g \circ h)$
2. $(\text{mon } g) \circ (\text{mon } f) =_\beta \text{mon}(f \circ g)$
3. $(\mathbb{T} f) \circ (\mathbb{T} g) =_\beta \mathbb{T}(f \circ g)$
4. $\text{mon } f p =_\beta p \circ f$
5. $\mathbb{T} f P =_\beta P \circ (\text{mon } f)$

We write simply $\Gamma \vdash A$ if $\Gamma \vdash M : A$ for some M .

2.1. Retraction and Γ_{ret} in λHOL

Definition 1 1. $X_0 := \lambda p: (\text{Pow } U). \Pi u: U. p u \rightarrow \neg(\text{out } u p)$ with $\mathbb{T}(U)$.

2. $q_0 := \lambda u: U. \Pi p: (\text{Pow } U). p u \rightarrow \neg(\text{out } u p)$ with $\text{Pow}(U)$.
3. $x_0 := \text{in } X_0$ with U .

Under Γ_{ret1} , λHOL proves the following proposition.

Proposition 1 1. $\Gamma_{ret1} \vdash l_1 : X_0 q_0$ for some l_1 .

2. $\Gamma_{ret1} \vdash l_2 : q_0 x_0$ for some l_2 .

Proof. (1) $l_1 := \lambda u: U. \lambda h: q_0 u. h q_0 h$ proves $X_0 q_0$, i.e., $\Pi u: U. q_0 u \rightarrow \neg(\text{out } u q_0)$.

(2) $l_2 := \lambda p: \text{Pow}(U). \lambda h: p x_0. \lambda h_1: X_0 p. (\alpha_e X_0 p) h_1 x_0 h h_1$ proves $q_0 x_0$, i.e., $\Pi p: \text{Pow}(U). p x_0 \rightarrow \neg(\text{out } x_0 p)$.

Proposition 2 1. λHOL proves that $\Gamma_{ret1} \vdash \perp$, and then $\Gamma_U \vdash \neg \text{Ret1}$ and $U: \square \vdash \Pi \text{in}: (\mathbb{T}U \rightarrow U). \Pi \text{out}: (U \rightarrow \mathbb{T}U). \neg \text{Ret1}$.

2. Let $F := (\alpha_e X_0 q_0) \circ (\beta_e X_0 q_0) : X_0 q_0 \rightarrow X_0 q_0$. λHOL never enjoy the reduction such that $F l \rightarrow_\beta l$.

Proof. (1) λHOL proves that $\Gamma_{ret1} \vdash l_1 x_0 l_2 (\beta_e X_0 q_0 l_1) : \perp$, and hence $\Gamma_U \vdash \lambda \text{ret1} : \text{Ret1}. l_1 x_0 l_2 (\beta_e X_0 q_0 l_1) : \neg \text{Ret1}$, and then $\Pi \text{in}: (\mathbb{T}U \rightarrow U). \Pi \text{out}: (U \rightarrow \mathbb{T}U). \neg \text{Ret1}$ by $(\square, *)$.

(2) Towards a contradiction, suppose that $\alpha_e X_0 q_0 (\beta_e X_0 q_0 l) \rightarrow_\beta l$, i.e., type isomorphism $(\text{out} \circ \text{in}) X p \simeq X p$. Then the head reduction gives

$$\begin{aligned} & l_1 x_0 l_2 (\beta_e X_0 q_0 l_1) \\ & \rightarrow_\beta^* l_2 q_0 l_2 (\beta_e X_0 q_0 l_1) \\ & \rightarrow_\beta^* \alpha_e X_0 q_0 (\beta_e X_0 q_0 l_1) x_0 l_2 (\beta_e X_0 q_0 l_1) \\ & \rightarrow_\beta^* l_1 x_0 l_2 (\beta_e X_0 q_0 l_1) \rightarrow_\beta^* \dots, \end{aligned}$$

leading to a contradiction to the fact that λHOL is strongly normalizing.

It should be remarked that we had $F l_1 x_0 l_2 (\beta_e X_0 q_0 l_1) =_\beta l_1 x_0 l_2 (\beta_e X_0 q_0 l_1)$ in the above so that the proof of the retract gave a certain fixed point of $F : X_0 q_0 \rightarrow X_0 q_0$. This simple example shows that an existence of type isomorphism $(\text{out} \circ \text{in}) X_0 p_0 \simeq X_0 p_0$ leads to a fixed point of functions with $X_0 q_0 \rightarrow X_0 q_0$.

$$\begin{array}{c}
\frac{[h_u : p_0 u] \quad \frac{\text{stablep}_0 : \Pi u : U. p_0 u \rightarrow p_0(\delta u)}{\text{stablep}_0 u h_u : p_0(\delta u)}}{\text{stablep}_0 u h_u : p_0(\delta u)} \quad \frac{[h_z : p_0 u] \quad \frac{h_z : \Pi p : \text{Pow}(U). p(\delta u) \rightarrow \neg(\text{out } u p)}{h_z p_0 : p_0(\delta u) \rightarrow \neg(\text{out } u p_0)}}{h_z p_0 : p_0(\delta u) \rightarrow \neg(\text{out } u p_0)} =_{\beta} \\
\frac{\frac{h_u p_0 (\text{stablep}_0 u h_u) : \neg(\text{out } u p_0)}{\lambda h_u : p_0 u. h_u p_0 (\text{stablep}_0 u h_u) : p_0 u \rightarrow \neg(\text{out } u p_0)} (*, *)}{\lambda u : U. \lambda h_u : p_0 u. h_u p_0 (\text{stablep}_0 u h_u) : \Pi u : U. p_0 u \rightarrow \neg(\text{out } u p_0)} (\square, *)
\end{array}$$

Let $\text{loop} := \text{lem1 } p_0 \text{ lem2 lem3}$.

Proposition 5 1. λHOL proves that $\Gamma_{\text{weak}} \vdash \text{loop} : \perp$.

2. λHOL proves that $\Gamma_U \vdash \neg \text{Weak}$ and then $U : \square \vdash \Pi \text{in} : \text{T}U \rightarrow U. \Pi \text{out} : U \rightarrow \text{T}U. \neg \text{Weak}$.

3. Encoding of sets and algebras in λU^-

We define instances of the powerful universe, i.e., a weakly initial T -algebra \mathbf{A} as in [2] and in addition here, a weakly final T -coalgebra \mathbf{B} , as follows.

1. For $A, B : \square$, product $A \times B := \Pi X : \square. (A \rightarrow B \rightarrow X) \rightarrow X$ with \square .
 $\langle M, N \rangle := \lambda X : \square. \lambda z : (A \rightarrow B \rightarrow X). z M N$ with $A \times B$
 $\pi_1 M := M A (\lambda x : A. \lambda y : B. x)$ with A where $M : A \times B$
 $\pi_2 M := M B (\lambda x : A. \lambda y : B. y)$ with B where $M : A \times B$
2. For $A, B : \square$, coproduct $A + B := \Pi X : \square. (A \rightarrow X) \rightarrow (B \rightarrow X) \rightarrow X$ with \square .
3. For $G(X) : \square$ with $X : \square$,
 existential quantifier $\Sigma X : \square. G(X) := \Pi Y : \square. (\Pi X : \square. G(X) \rightarrow Y) \rightarrow Y$ with \square .
 $[B, M] := \lambda Y : \square. \lambda z : (\Pi X : \square. G(X) \rightarrow Y). z B M$ with $\Sigma X : \square. G(X)$.
 $(\text{let } [X, z] = M \text{ in } N) := M C (\lambda X : \square. \lambda z : G(X). N)$ with $C : \square$,
 where $N : C$ and $M : (\Sigma X : \square. G(X))$.
 Note that $(\text{let } [X, z] = [B, M] \text{ in } N) \rightarrow_{\beta} N[X := B, z := M]$.

Now the context Γ_U in λHOL can be established by the following algebras in λU^- , so that the condition *Weak* holds in λU^- under the weakly initial and cofinal algebras, respectively.

3.1. Weakly initial algebra

We introduce the powerful universe \mathbf{A} with `fold`, `intro`, and `match`. Then the following Proposition 6 can be established [2].

1. Powerful universe $\mathbf{A} := \Pi X : \square. (T(X) \rightarrow X) \rightarrow X$ with super kind \square .
 Let $\text{fold}_X^{\mathbf{A}} := \lambda f. \lambda a^{\mathbf{A}}. a X f$ with $(T(X) \rightarrow X) \rightarrow (\mathbf{A} \rightarrow X)$.
 Let $\text{intro} := \lambda u : T(\mathbf{A}). \lambda X : \square. \lambda f. f(T \mathbf{A} X (\text{fold}_X^{\mathbf{A}} f) u)$ with $T(\mathbf{A}) \rightarrow \mathbf{A}$.
 Let $\text{match} := \text{fold}_{T(\mathbf{A})}^{\mathbf{A}} (T(T \mathbf{A}) \text{intro})$ with $\mathbf{A} \rightarrow T(\mathbf{A})$, i.e., $\text{match}(a) = a(T \mathbf{A})(T \text{intro})$.
2. **Proposition 6** $\langle \mathbf{A}, \text{intro} \rangle$ is weakly initial: for any $f : T(X) \rightarrow X$, the following is commutative.

$$\begin{aligned}
& ((\text{fold}_X^{\mathbf{A}} f) \circ \text{intro}) =_{\beta} (f \circ (T(\text{fold}_X^{\mathbf{A}} f))) \\
& (\text{match} \circ \text{intro}) =_{\beta} ((T \text{intro}) \circ (T \text{match}))
\end{aligned}$$

3.2. Weakly final co-algebra

We here define the powerful universe \mathbf{B} together with `unfold`, `in`, and `out` as well. Then the following Proposition 7 introduces a weakly final T -coalgebra.

1. Powerful universe $\mathbf{B} := \Sigma X : \square. (X \rightarrow \mathbf{T}(X)) \times X$ with super kind \square .

Let $\text{unfold}_X^{\mathbf{B}} := \lambda f. \lambda z : X. [X, \langle f, z \rangle]$ with $(X \rightarrow \mathbf{T}(X)) \rightarrow (X \rightarrow \mathbf{B})$.

Let $\text{out} := \lambda b : \mathbf{B}. (\text{let } [X, \beta] = b \text{ in } \mathbf{T}(\text{unfold}_X(\pi_1 \beta)) (\pi_1 \beta (\pi_2 \beta)))$
with $\mathbf{B} \rightarrow \mathbf{T}(\mathbf{B})$.

Let $\text{in} := \text{unfold}_{\mathbf{T}\mathbf{B}}^{\mathbf{B}}(\mathbf{T}(\mathbf{T}\mathbf{B}) \text{out})$ with $\mathbf{T}\mathbf{B} \rightarrow \mathbf{B}$, i.e., $\text{in}(u) = [\mathbf{T}\mathbf{B}, \langle \mathbf{T}\text{out}, u \rangle]$.

2. **Proposition 7** $\langle \mathbf{B}, \text{out} \rangle$ is weakly final: for any $f : X \rightarrow \mathbf{T}(X)$, the following is commutative.

$$(\text{out} \circ (\text{unfold}_X^{\mathbf{B}} f)) =_{\beta} ((\mathbf{T}(\text{unfold}_X^{\mathbf{B}} f)) \circ f)$$

$$(\text{out} \circ \text{in}) =_{\beta} ((\mathbf{T} \text{in}) \circ (\mathbf{T} \text{out}))$$

3.3. Computational behaviour of the proof of $\vdash \perp$ in $\lambda\mathbf{U}^-$

Now for both algebras, we obtain the same computational behaviour of the proof, which do not produce a term that reduces to itself. In the proof Proposition4, the use of Prop3 can be simply replaced with $=_{\beta}$ by removing L_1, L_2 , and then reset `loop` := `lem1 p0 lem2 lem3`. Then `loop` has type \perp , and there exists an infinite reduction path with head redexes of `lem1`, `2`, `3` [3], which means that redexes (cut-formulae) cannot be removed as follows.

$$\begin{aligned} \text{loop} &\xrightarrow{+}_{\beta} \text{lem3 } x_0 \text{ lem2 } (\text{stable}_{X_0} p_0 \text{ lem3}) \\ &\xrightarrow{+}_{\beta} \text{lem2 } p_0 (\text{stable}_{p_0} x_0 \text{ lem2}) (\text{stable}_{X_0} p_0 \text{ lem3}) \\ &\xrightarrow{+}_{\beta} \text{lem1 } (p_0 \circ \delta) (\text{stable}_{p_0} x_0 \text{ lem2}) (\text{stable}_{X_0} p_0 \text{ lem3}) \\ &\xrightarrow{+}_{\beta} \text{lem1 } (p_0 \circ \delta \circ \delta) (\text{stable}_{p_0} (\delta x_0) (\text{stable}_{p_0} x_0 \text{ lem2})) \\ &\quad (\text{stable}_{X_0} (p_0 \circ \delta) (\text{stable}_{X_0} p_0 \text{ lem3})) \\ &\xrightarrow{+}_{\beta} \dots \end{aligned}$$

4. Concluding Remarks

Independently of the concrete structure of the powerful universe, λHOL can prove \perp under the assumption Γ_{weak} or Γ_{ret} , so that λHOL derives $\Gamma_U \vdash \neg \text{Weak}$; $\Gamma_U \vdash \neg \text{Ret1}$; $\Gamma_U \vdash \text{Ret1} \wedge \text{Ret2} \rightarrow \text{Weak}$; $\Gamma_U \vdash \text{Ret1} \wedge \text{Weak} \rightarrow \text{Ret2}$; $\Gamma_U \vdash \text{Ret2} \wedge \text{Weak} \rightarrow \text{Ret1}$.

$\lambda\mathbf{U}^-$ introduces not only the weakly initial T -algebra [2], but also the dual one, i.e., a weakly final T -coalgebra as instances of the powerful universe.

It should be remarked that natural transformations can be used to characterize functions on final coalgebras having a unique fixed point [6], and this property should be applied to our final T -coalgebra.

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