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# On competition numbers of complete multipartite graphs with partite sets of equal size

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## Abstract

Let  $D$  be an acyclic digraph. The competition graph of  $D$  is a graph which has the same vertex set as  $D$  and has an edge between  $u$  and  $v$  if and only if there exists a vertex  $x$  in  $D$  such that  $(u, x)$  and  $(v, x)$  are arcs of  $D$ . For any graph  $G$ ,  $G$  together with sufficiently many isolated vertices is the competition graph of some acyclic digraph. The competition number  $k(G)$  of  $G$  is the smallest number of such isolated vertices. In general, it is hard to compute the competition number  $k(G)$  for a graph  $G$  and it has been one of important research problems in the study of competition graphs to characterize a graph by its competition number.

In this paper, we compute the competition numbers of a complete multipartite graph in which each partite set has two vertices and a complete multipartite graph in which each partite set has three vertices.

**Keywords:** competition graphs, competition numbers, complete multipartite graphs

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## 1 Introduction

Suppose  $D$  is an acyclic digraph (for all undefined graph-theoretical terms, see [1] and [15]). The *competition graph* of  $D$ , denoted by  $C(D)$ , has the same set of vertices as  $D$  and an edge between vertices  $u$  and  $v$  if and only if there is a vertex  $x$  in  $D$  such that  $(u, x)$  and  $(v, x)$  are arcs of  $D$ . Roberts [14] observed that if  $G$  is any graph,  $G$  together with sufficiently many isolated vertices is the competition graph of an acyclic digraph. Then he defined the *competition number*  $k(G)$  of a graph  $G$  to be the smallest number  $k$  such that  $G$  together with  $k$  isolated vertices added is the competition graph of an acyclic digraph.

The notion of competition graph was introduced by Cohen [3] as a means of determining the smallest dimension of ecological phase space. Since then, various variations have been defined and studied by many authors (see, for example, [2, 7, 8, 9, 11, 16]). Besides an application to ecology, the concept of competition graph can be applied to the study of communication over noisy channel (see Roberts [14] and Shannon [17]) and to problem of assigning channels to radio or television transmitters (see Cozzens and Roberts [4], Hale [6], or Opsut and Roberts [13]).

Roberts [14] observed that characterization of competition graph is equivalent to computation of competition number. It does not seem to be easy in general to compute  $k(G)$  for all graphs  $G$ , as Opsut [12] showed that the computation of the competition number of a graph is an NP-hard problem (see [8, 9] for graphs whose competition numbers are known). It has been one of important research problems in the study of competition graphs to characterize a graph by its competition number.

In this paper, we shall compute competition numbers of a complete multipartite graph in which each partite set has two vertices and a complete multipartite graph in which each partite set has three vertices.

We denote by  $K_n^m$  the complete  $m$ -partite graph in which each partite set has  $n$  vertices.

For a digraph  $D$ , a sequence  $v_1, v_2, \dots, v_n$  of the vertices of  $D$  is called an *acyclic ordering* of  $D$  if  $(v_i, v_j) \in A(D)$  implies  $i > j$ . It is well-known that a digraph  $D$  is acyclic if and only if there exists an acyclic ordering of  $D$ .

An *edge clique cover* (or an ECC for short) of a graph  $G$  is a family of cliques such that each edge of  $G$  is contained in some clique in the family. The smallest size of an ECC of  $G$  is called the *edge clique cover number* (or the ECC number for short), and is denoted by  $\theta_e(G)$ . Opsut [12] gave

bounds of  $k(G)$  for a graph  $G$  by showing that  $\theta_e(G) - |V(G)| + 2 \leq k(G) \leq \theta_e(G)$ . Dutton and Brigham [5] characterized the competition graphs of acyclic digraphs in terms of an ECC as follows:

**Theorem 1** ([5]). *A graph  $G$  is the competition graph of an acyclic digraph if and only if there exists an ordering  $v_1, \dots, v_n$  of the vertices of  $G$  and an ECC  $\{S_1, \dots, S_n\}$  of  $G$  such that  $v_i \in S_j$  implies  $i < j$ .*

For a vertex  $v$  in a graph  $G$ , let the open neighborhood of  $v$  be denoted by

$$N_G(v) = \{u \mid u \text{ is adjacent to } v\}$$

and the closed neighborhood of  $v$  be denoted by  $N_G[v] = N_G(v) \cup \{v\}$ . We denote the subgraph of  $G$  induced by  $N_G(v)$  (resp.  $N_G[v]$ ) by the same symbol  $N_G(v)$  (resp.  $N_G[v]$ ). For a digraph  $D$ , we define  $N_D^-(v) = \{w \in V(D) \mid (w, v) \in A(D)\}$  and  $N_D^+(v) = \{w \in V(D) \mid (v, w) \in A(D)\}$ .

A *vertex clique cover* of a graph  $G$  is a family of cliques such that each vertex of  $G$  is contained in some clique in the family. The smallest size of a vertex clique cover of  $G$  is called the *vertex clique cover number*, and is denoted by  $\theta_v(G)$ . Opsut [12] showed the following:

**Proposition 2** ([12]). *Let  $G$  be a graph. Then we have*

$$\min\{\theta_v(N_G(v)) \mid v \in V(G)\} \leq k(G).$$

This proposition is true even if each open neighborhood is replaced with the closed neighborhood.

**Proposition 3.** *Let  $G$  be a graph. Then we have*

$$\min\{\theta_v(N_G[v]) \mid v \in V(G)\} \leq k(G).$$

*Proof.* Let  $t = \min\{\theta_v(N_G[v]) \mid v \in V(G)\}$  and  $k = k(G)$ . Let  $D$  be an acyclic digraph such that  $C(D) = G \cup I_k = G \cup \{z_1, \dots, z_k\}$ . Let  $z_1, \dots, z_k, v_1, \dots, v_n$  be an ordering of  $D$  so that  $(u, v)$  is an arc of  $D$  only if  $u$  is on the right hand side of  $v$  in the sequence. Then we have  $|N_D^+(v_1)| \geq t$  since  $\theta_v(N_G[v_1]) \geq t$ . Since  $N_D^+(v_1) \subseteq \{z_1, \dots, z_k\}$ , we have  $t \leq k$  and thus the proposition holds.  $\square$

For some special graph families, we have explicit formulae for computing competition numbers. For example, if  $G$  is a choral graph with the minimum degree  $\geq 1$  then  $k(G) = 1$ , and if  $G$  is a triangle-free connected graph then

$$k(G) = |E(G)| - |V(G)| + 2$$

(see [14]). From this formula, it follows that for a complete bipartite graph  $K_{n_1, n_2}$ , we have  $k(K_{n_1, n_2}) = n_1 n_2 - (n_1 + n_2) + 2$ . Kim and Sano [10] gave the exact competition number of a complete tripartite graph  $K_n^3$ :

**Theorem 4** ([10]). *For  $n \geq 2$ ,*

$$k(K_n^3) = n^2 - 3n + 4.$$

Then they proposed an open problem to compute competition numbers of various types of complete multipartite graphs. To answer their question, we study  $K_n^m$ . We give a lower bound for the competition number of  $K_n^m$ . Then we make use of it to compute the competition numbers of  $K_2^m$  and  $K_3^m$ .

## 2 A lower bound for the competition number of $K_n^m$

In this section, we give a lower bound for  $k(K_n^m)$  for integers  $m \geq 2$  and  $n \geq 1$ .

For each positive integer  $n$ , we denote the set  $\{1, \dots, n\}$  by  $[n]$ .

**Proposition 5.** *For any vertex  $v$  of  $K_n^m$  with  $m \geq 2$ ,*

$$\theta_v(N_{K_n^m}[v]) \geq n.$$

*Proof.* Let  $P_k$  denote the  $k$ th partite set of  $K_n^m$  and let  $P_k = \{v_{k1}, \dots, v_{kn}\}$  for each  $k \in [m]$ . From the fact that  $v_{11}$  is adjacent to all the vertices in  $V(K_n^m) \setminus P_1$  and that any two vertices in  $P_2$  are not adjacent, we know that at least  $n$  cliques are needed to cover the  $n$  edges  $v_{11}v_{21}, \dots, v_{11}v_{2n}$  which are incident to  $v_{11}$ . Therefore we have  $\theta_v(N_{K_n^m}[v_{11}]) \geq n$ . By symmetry, we can conclude  $\theta_v(N_{K_n^m}[v]) \geq n$  for any  $v \in V(K_n^m)$ .  $\square$

Given a digraph  $D = (V, A)$ , we use a symbol  $u \rightarrow v$  for arc  $(u, v)$  in  $A$ . In addition, for  $S \subset V$ , we denote by  $S \rightarrow w$  the arc set  $\{(x, w) \mid x \in S\}$ .

**Theorem 6.** *For an integer  $m \geq 2$ ,*

$$k(K_n^m) \geq 2n - 2.$$

*Proof.* Let  $k$  be the competition number of  $K_n^m$ . Then there exists an acyclic digraph  $D$  such that  $C(D) = K_n^m \cup I_k$  where  $I_k = \{a_1, a_2, \dots, a_k\}$ . Also, let  $a_1, a_2, \dots, a_k, v_1, v_2, \dots, v_{mn}$  be an acyclic ordering of  $D$ . By Proposition 5,

we have  $\theta_v(N_{K_n^m}[v_i]) \geq n$  for  $i = 1, \dots, mn$ . Thus  $v_i$  has at least  $n$  distinct prey, that is,

$$|N_D^+(v_i)| \geq n. \quad (1)$$

However, since  $v_1$  and  $v_2$  are the vertices of the lowest index and the second lowest index, respectively,

$$N_D^+(v_1) \subset I_k \quad \text{and} \quad N_D^+(v_2) \subset I_k \cup \{v_1\}.$$

Thus,

$$N_D^+(v_1) \cup N_D^+(v_2) \subset I_k \cup \{v_1\}. \quad (2)$$

Let  $P_1$  and  $P_2$  be the partite sets to which  $v_1$  and  $v_2$  belong, respectively. Then  $P_1 = P_2$  if  $v_1$  and  $v_2$  are nonadjacent, and  $P_1 \neq P_2$  if  $v_1$  and  $v_2$  are adjacent.

Firstly, assume that  $v_1$  and  $v_2$  are nonadjacent. Then  $v_1$  and  $v_2$  do not share a prey, that is,

$$N_D^+(v_1) \cap N_D^+(v_2) = \emptyset. \quad (3)$$

By (1), (2) and (3),

$$k \geq |N_D^+(v_1) \cup N_D^+(v_2)| - 1 = |N_D^+(v_1)| + |N_D^+(v_2)| - 1 \geq 2n - 1.$$

Now suppose that  $v_1$  and  $v_2$  are adjacent. Then  $P_1 \neq P_2$ . Let

$$A_1 = \{a \mid a \text{ is a common prey of } v_1 \text{ and } v \text{ for } v \in P_2 \setminus \{v_2\}\},$$

$$A_2 = \{a \mid a \text{ is a common prey of } v_2 \text{ and } v \text{ for } v \in P_1 \setminus \{v_1\}\}.$$

Let  $b$  be a common prey of  $v_1$  and  $v_2$ . Then  $b \notin A_1 \cup A_2$ . For, otherwise,  $v_1, v_2, v$  form a triangle for some  $v \in P_1 \cup P_2$ , which is a contradiction. Then  $A_1 \cup \{b\} \subset N_D^+(v_1)$  and  $A_2 \cup \{b\} \subset N_D^+(v_2)$ . Since  $v_1$  is adjacent to every vertex of  $P_2$  and any pair of vertices in  $P_2$  is nonadjacent,  $|A_1| \geq n - 1$ . For the same reason,  $|A_2| \geq n - 1$ . Suppose that there is a vertex  $a$  belonging to  $A_1 \cap A_2$ . Then  $v \rightarrow a, v_1 \rightarrow a, w \rightarrow a, v_2 \rightarrow a$  for some  $v \in P_2 \setminus \{v_2\}$  and  $w \in P_1 \setminus \{v_1\}$ . Then  $v_1, v_2, v, w$  form  $K_4$ . This implies that  $v_1$  is adjacent to  $w$  and we reach a contradiction. Thus,  $A_1 \cap A_2 = \emptyset$ . Hence

$$k \geq |N_D^+(v_1) \cup N_D^+(v_2)| - 1 \geq (|A_1 \cup A_2| + 1) - 1 = |A_1| + |A_2| \geq 2n - 2.$$

□

### 3 The competition number of $K_2^m$

In this section, we compute the competition number of  $K_2^m$ , which is often called a ‘cocktail party graph’.

**Theorem 7.** For  $m \geq 2$ ,

$$k(K_2^m) = 2.$$

*Proof.* By Theorem 6, it is true that  $k(K_2^m) \geq 2$ . Now we show that  $k(K_2^m) \leq 2$ . If  $m = 2$ , then  $K_2^2$  is a 4-cycle and it is well-known that  $k(K_2^2) = 2$ . Now suppose that  $m \geq 3$ . Let  $\{x_i, y_i\}$  be the  $i$ th partite set of  $K_2^m$  for each  $i \in [m]$ . We let

$$\begin{aligned} S_1 &= \{x_1, x_2, x_3, \dots, x_m\}; \\ S_2 &= \{x_1, y_2, y_3, \dots, y_m\}; \\ S_3 &= \{y_1, x_2, y_3, \dots, y_m\}; \\ S_4 &= \{y_1, y_2, x_3, \dots, y_m\}; \\ &\vdots \\ S_{m+1} &= \{y_1, y_2, y_3, \dots, x_m\}. \end{aligned}$$

We denote by  $\mathcal{S}$  the family of those sets just defined. Since no two vertices in  $S_i$  belong to the same partite set, it is true that  $S_i$  forms a clique in  $K_2^m$  for each  $i \in [m+1]$ . Now we take an edge  $e$  of  $K_2^m$ . Then  $e = x_i x_j$ ,  $e = x_i y_j$ , or  $e = y_i y_j$  for some  $i, j \in [m]$ ,  $i \neq j$ . If  $e = x_i x_j$ , then  $e$  is covered by the set  $S_1$ . If  $e = x_i y_j$ , then  $e$  is covered by  $S_{i+1}$ . If  $e = y_i y_j$ , then  $e$  is covered by  $S_k$  for some  $k \in [m] \setminus \{1, i+1, j+1\}$ . Thus  $\mathcal{S}$  is an ECC of  $K_2^m$ .

We construct a digraph  $D$  as follows:

$$\begin{aligned} V(D) &= V(K_2^m) \cup \{a, b\}; \\ A(D) &= \{(u, a) \mid u \in S_1\} \cup \{(u, b) \mid u \in S_2\} \cup \bigcup_{i=3}^{m+1} \{(u, x_{i-2}) \mid u \in S_i\}. \end{aligned}$$

Since  $x_i$  is not contained in  $S_j$  for any  $j \geq i+2$ , it is true that  $D$  is acyclic. For each  $x \in V(D)$ , either  $N_D^-(x) = \emptyset$  or  $N_D^-(x) = S_i$  for some  $i \in [m+1]$ . Thus  $C(D) = K_2^m \cup \{a, b\}$ .  $\square$

## 4 The competition number of $K_3^m$

In this section, we compute  $k(K_3^m)$  for any  $m \geq 2$ . As  $K_3^2$  is triangle-free, it is true that  $k(K_3^2) = 9 - 6 + 2 = 5$ . For  $m \geq 3$ , we present the following theorem.

**Theorem 8.** For  $m \geq 3$ ,

$$k(K_3^m) = 4.$$

*Proof.* Given an integer  $m \geq 3$ , there exist a positive integer  $t$  and an integer  $r$  such that  $m = 3t + r$  and  $r \in \{0, 1, 2\}$ . Now we take  $3t$  partite sets of  $K_3^m$  and put them into  $t$  groups so that each group contains 3 partite sets. For each  $i \in [t]$ , we denote the  $j$ th partite set in the  $i$ th group by  $P_{ij} = \{x_{ij}, y_{ij}, z_{ij}\}$  for each  $j = 1, 2, 3$ . In case of  $r > 0$ , we have  $r$  remaining partite sets and denote them by  $P_{t+1,j} = \{x_{t+1,j}, y_{t+1,j}, z_{t+1,j}\}$  for  $j \in [r]$ .

Let  $Q_i = P_{i1} \cup P_{i2} \cup P_{i3}$  for each  $i \in [t]$  and  $Q_{t+1} = P_{t+1,1} \cup \cdots \cup P_{t+1,r}$ . Note that, for  $i \in [t]$ , the subgraph of  $K_3^m$  induced by  $Q_i$  is isomorphic to  $K_3^3$ , and the subgraph of  $K_3^m$  induced by  $Q_{t+1}$  is isomorphic to  $K_3^r$ .

For each  $i \in [t]$ , we let

$$\begin{aligned} S(x_{i1}) &= \{x_{i1}, y_{i2}, y_{i3}\}, & S(y_{i1}) &= \{y_{i1}, z_{i2}, z_{i3}\}, & S(z_{i1}) &= \{z_{i1}, x_{i2}, x_{i3}\}, \\ S(x_{i2}) &= \{x_{i2}, y_{i1}, y_{i3}\}, & S(y_{i2}) &= \{y_{i2}, z_{i1}, z_{i3}\}, & S(z_{i2}) &= \{z_{i2}, x_{i1}, x_{i3}\}, \\ S(x_{i3}) &= \{x_{i3}, y_{i1}, y_{i2}\}, & S(y_{i3}) &= \{y_{i3}, z_{i1}, z_{i2}\}, & S(z_{i3}) &= \{z_{i3}, x_{i1}, x_{i2}\}. \end{aligned}$$

We denote the collection of 9 sets defined above by  $\mathcal{S}_i$ . Note that any two vertices in each set in  $\mathcal{S}_i$  belong to distinct partite sets. Thus each of the above sets forms a clique in  $K_3^m$ . It is also easy to check that  $\mathcal{S}_i$  is an ECC of  $K_3^3$  induced by  $Q_i$ . If  $r > 0$ , then we define 9 more sets: If  $r = 1$ , then

$$\begin{aligned} S(x_{t+1,1}) &= \{x_{t+1,1}\}, & S(y_{t+1,1}) &= \{y_{t+1,1}\}, & S(z_{t+1,1}) &= \{z_{t+1,1}\}, \\ S(x_{t+1,2}) &= \{y_{t+1,1}\}, & S(y_{t+1,2}) &= \{z_{t+1,1}\}, & S(z_{t+1,2}) &= \{x_{t+1,1}\}, \\ S(x_{t+1,3}) &= S(x_{t+1,2}), & S(y_{t+1,3}) &= S(y_{t+1,2}), & S(z_{t+1,3}) &= S(z_{t+1,2}). \end{aligned}$$

If  $r = 2$ , then

$$\begin{aligned} S(x_{t+1,1}) &= \{x_{t+1,1}, y_{t+1,2}\}, & S(y_{t+1,1}) &= \{y_{t+1,1}, z_{t+1,2}\}, \\ S(z_{t+1,1}) &= \{z_{t+1,1}, x_{t+1,2}\}, & S(x_{t+1,2}) &= \{y_{t+1,1}, x_{t+1,2}\}, \\ S(y_{t+1,2}) &= \{z_{t+1,1}, y_{t+1,2}\}, & S(z_{t+1,2}) &= \{x_{t+1,1}, z_{t+1,2}\}, \\ S(x_{t+1,3}) &= \{y_{t+1,1}, y_{t+1,2}\}, & S(y_{t+1,3}) &= \{z_{t+1,1}, z_{t+1,2}\}, \\ S(z_{t+1,3}) &= \{x_{t+1,1}, x_{t+1,2}\}. \end{aligned}$$



It is easy to check that the above sets form an ECC of  $K_3^r$  induced by  $Q_{t+1}$ . For convenience, if  $r = 0$ , then we let

$$\begin{aligned} S(x_{t+1,1}) &= S(y_{t+1,1}) = S(z_{t+1,1}) = S(x_{t+1,2}) = S(y_{t+1,2}) = S(z_{t+1,2}) \\ &= S(x_{t+1,3}) = S(y_{t+1,3}) = S(z_{t+1,3}) = \emptyset. \end{aligned}$$

We also let  $S(x_{t+2,j}) = S(y_{t+2,j}) = S(z_{t+2,j}) = \emptyset$  for  $j = 1, 2, 3$ .

Now we define a digraph  $D$  as follows:

$$\begin{aligned} V(D) &= V(K_3^m) \cup \{a_1, a_2, a_3, a_4\} \\ A(D) &= \bigcup_{i=1}^{t+1} A_i \end{aligned}$$

where  $A_i$  is the union of arc sets

$$\begin{aligned} \bigcup_{\ell=1}^{t+1} S(x_{\ell 1}) \rightarrow a_1, & \quad S(x_{i2}) \cup \bigcup_{\ell=i+1}^{t+1} S(y_{\ell 1}) \rightarrow x_{i1}, \\ S(x_{i3}) \cup \bigcup_{\ell=i+1}^{t+1} S(z_{\ell 1}) \rightarrow x_{i2}, & \quad S(y_{i1}) \cup \bigcup_{\ell=i+1}^{t+1} S(x_{\ell 1}) \rightarrow x_{i3}, \\ S(y_{i2}) \cup \bigcup_{\ell=i+1}^{t+1} S(y_{\ell 1}) \rightarrow y_{i1}, & \quad S(y_{i3}) \cup \bigcup_{\ell=i+1}^{t+1} S(z_{\ell 1}) \rightarrow y_{i2}, \\ S(z_{i1}) \cup \bigcup_{\ell=i+1}^{t+1} S(x_{\ell 1}) \rightarrow z_{i-1,1}, & \quad S(z_{i2}) \cup \bigcup_{\ell=i+1}^{t+1} S(y_{\ell 1}) \rightarrow z_{i-1,2}, \\ S(z_{i3}) \cup \bigcup_{\ell=i+1}^{t+1} S(z_{\ell 1}) \rightarrow z_{i-1,3}. \end{aligned}$$

where  $z_{01} = a_2$ ,  $z_{02} = a_3$ , and  $z_{03} = a_4$ .

We denote by  $D_i$  the subdigraph of  $D$  induced by  $Q_i \cup \{a_1, z_{i-1,1}, z_{i-1,2}, z_{i-1,3}\}$  for each  $i \in [t+1]$ . Now we order the vertices of  $D_i$  as follows:

$$a_1, z_{i-1,1}, z_{i-1,2}, z_{i-1,3}, x_{i1}, x_{i2}, x_{i3}, y_{i1}, y_{i2}, y_{i3}, z_{i1}, z_{i2}, z_{i3}.$$

Then we can easily check that  $u \rightarrow v$  in  $D_i$  only if  $v$  is on the left hand side of  $u$  in the above sequence. Thus  $D_i$  is acyclic for each  $i \in [t+1]$ . Furthermore, an arc goes from a vertex in the  $j$ th partite set  $Q_j$  to the  $i$ th partite set  $Q_i$  in  $D$  only if  $i \leq j$ . Therefore  $D$  is acyclic.

Two nonadjacent vertices in  $G$  belong to  $Q_i$  for a unique  $i \in [t+1]$ . Moreover they cannot belong to the same clique in  $\mathcal{S}_i$ . However, no two vertices from distinct cliques in  $\mathcal{S}_i$  prey on a common vertex in  $D$  and so they are nonadjacent in  $C(D)$ . Therefore  $E(C(D)) \subset E(K_3^m)$ .

We show that  $E(C(D)) \supset E(K_3^m)$  in the following. We note that for each vertex  $v$  in  $D_i$ ,  $N_{D_i}^-(v)$  is either a clique in  $\mathcal{S}_i$  or an empty set and that each clique in  $\mathcal{S}_i$  is the neighborhood of a vertex in  $D_i$  for each  $i \in [t]$ . Thus, the competition graph of  $D_i$  is  $K_3^3 \cup \{a_1, z_{i-1,1}, z_{i-1,2}, z_{i-1,3}\}$  if  $i \in [t]$ . Similarly,  $C(D_{t+1}) = K_3^r \cup \{a_1, z_{t,1}, z_{t,2}, z_{t,3}\}$  if  $r > 0$ .

It remains to show that a vertex in  $Q_i$  and a vertex in  $Q_j$  are adjacent in  $C(D)$  for  $i, j \in [t]$  with  $i < j$ . We note that for any  $i \in [t]$  and  $k = 1, 2, 3$ , it holds that  $Q_i$  is the disjoint union of  $S(x_{ik}), S(y_{ik})$  and  $S(z_{ik})$ . Now we take a vertex  $u \in Q_j$ . Then  $u \in S(x_{j1})$  or  $u \in S(y_{j1})$  or  $u \in S(z_{j1})$ .

Firstly consider the case  $u \in S(x_{j1})$ . Then in  $D$ ,

$$S(x_{i1}) \cup \{u\} \rightarrow a_1, \quad S(y_{i1}) \cup \{u\} \rightarrow x_{i3}, \quad \text{and} \quad S(z_{i1}) \cup \{u\} \rightarrow z_{i-1,1}.$$

Thus  $u$  is adjacent to every vertex in  $Q_i$ .

Now suppose that  $u \in S(y_{j1})$ . Then in  $D$ ,

$$S(x_{i2}) \cup \{u\} \rightarrow x_{i1}, \quad S(y_{i2}) \cup \{u\} \rightarrow y_{i1}, \quad \text{and} \quad S(z_{i2}) \cup \{u\} \rightarrow z_{i-1,2}.$$

Thus  $u$  is adjacent to every vertex in  $Q_i$ .

Finally suppose that  $u \in S(z_{j1})$ . Then in  $D$ ,

$$S(x_{i3}) \cup \{u\} \rightarrow x_{i2}, \quad S(y_{i3}) \cup \{u\} \rightarrow y_{i2}, \quad \text{and} \quad S(z_{i3}) \cup \{u\} \rightarrow z_{i-1,3}.$$

Thus  $u$  is adjacent to every vertex in  $Q_i$ . Therefore,  $E(C(D)) \supset E(K_3^m)$ . This completes the proof that  $C(D) = K_3^m \cup \{a_1, a_2, a_3, a_4\}$ . Hence  $k(K_3^m) \leq 4$ . From Theorem 6, it follows that  $k(K_3^m) = 4$ .  $\square$

## 5 Concluding remarks

In this paper, we give bounds for  $K_n^m$  and computed the competition numbers of  $K_3^m$  and  $K_2^m$ . Note that  $k(K_3^m) = k(K_3^3) = 4$  for  $m \geq 3$  and  $k(K_2^m) = k(K_2^2) = 2$  for  $m \geq 2$ . We conjecture that  $k(K_n^m) = k(K_n^n)$  for  $m \geq n$ .

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