Certifying the hyperbolicity of knots and links

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The Geometrisation Conjecture was very difficult to prove.

But in practice, it is remarkably easy to find a hyperbolic structure on a 3-manifold.

Why?!

Finding hyperbolic structures

<u>Question</u>: What is the computational complexity of determining whether a compact 3-manifold is hyperbolic and, if it is hyperbolic, how hard is it to find the hyperbolic structure?

Previous work

[Manning, Casson] One can determine whether a closed 3-manifold M is hyperbolic and find its hyperbolic structure, as long as $\pi_1(M)$ has solvable word problem.

[Kuperberg] Gave an algorithm to determine whether M is hyperbolic and to find its hyperbolic structure that runs in time that is elementary recursive ie at most



where t is the number of tetrahedra in a given triangulation of M. [Scull] The running time is at most

 $2^{2^{t^{O(t)}}}$

Ruling out hyperbolic structures

If a closed orientable 3-manifold is not hyperbolic, then it is one of:

- Seifert fibred
- reducible
- toroidal.

<u>Theorem</u>: [Ivanov, Schleimer] S^3 recognition is in NP.

<u>Theorem</u>: [Lackenby-Schleimer] Recognition of elliptic 3-manifolds is in NP.

<u>Theorem</u>: [Jackson] Recognition of Seifert fibre spaces with non-empty boundary is in NP.

Closed Seifert fibre spaces remain a problem, particularly the small ones.

Hyperbolic structures on link complements

<u>Problem</u>: (LINK HYPERBOLICITY) Given a diagram of a link L with c crossings, is L hyperbolic?

<u>Theorem</u>: [Haraway-Hoffman, Badwin-Sivek] This problem is in co-NP.

Theorem: [Baroni, Lackenby] This problem is in NP.

Showing that a link is not hyperbolic

<u>Theorem</u>: [Thurston] Let L be a link in the 3-sphere. Then one of the following holds:

- L is the unknot;
- L is split;
- there is an essential torus in $S^3 L$;
- there is an essential annulus in $S^3 L$;
- L is hyperbolic.

Haraway-Hoffman used the following fact:

<u>Theorem</u>: [Lackenby] Deciding whether a compact orientable 3-manifold has incompressible boundary is in NP.

From now onwards, we'll focus on:

Theorem: [Baroni, Lackenby] Link hyperbolicity is in NP.

Given a hyperbolic link L, the proof of its hyperbolicity divides into two cases:

- L is a fibred knot [Baroni]
- L is not fibred or has more than one component [Lackenby].

We will start by examining the fibred case.

The Nielsen-Thurston type of a surface automorphism

Let S be an orientable surface of finite type and $\chi(S) < 0$, and let $\phi: S \to S$ be a homeomorphism. Then exactly one of the following holds:

- 1. ϕ is periodic;
- 2. ϕ is reducible;
- 3. ϕ is pseudo-anosov (\Leftrightarrow ($S \times I$)/ ϕ is hyperbolic)

Suppose that we are given ϕ as a word w in 'standard generators' in the mapping class group of S.

<u>Theorem</u>: [Bell-Webb] For a fixed surface S (with at least one puncture), there is an algorithm to determine the Nielsen-Thurston type of ϕ that runs in polynomial time in the length of w.

<u>Theorem</u>: [Baroni] There is an algorithm to determine the Nielsen-Thurston type of ϕ that runs in polynomial time in the length of w and in $|\chi(S)|$.

This uses hierarchies.

A hierarchy is a sequence of compact orientable 3-manifolds $M = M_1, \ldots, M_{\ell+1}$ and orientable surfaces S_1, \ldots, S_ℓ such that:

• each S_i is properly embedded in M_i ;

• each
$$M_{i+1} = M_i \setminus S_i$$
;

• $M_{\ell+1}$ is a collection of 3-balls.

Boundary patterns

A boundary pattern for a 3-manifold M is subset P of ∂M that is a disjoint union of simple closed curves and trivalent graphs.

The manifolds in a hierarchy naturally inherit a boundary pattern.



Essential boundary patterns

A boundary pattern P is essential if, for any properly embedded disc D that intersects P at most three times, ∂D bounds a disc D' in ∂M that intersects P in one of the following:



- an arc,
- a tripod.



A hierarchy $M = M_1, \ldots, M_{\ell+1}$ is essential if the final manifold $M_{\ell+1}$ inherits an essential boundary pattern.

<u>Theorem</u>: [Waldhausen, Johansson] Let M be a compact orientable 3-manifold with non-empty boundary and empty boundary pattern. Then the following are equivalent:

- ∂M is incompressible and M is irreducible;
- ▶ *M* has an essential hierarchy.

An example: the knot 5_2



(i) The knot 52



(iii) The exterior of this surface



(v) The pattern of one of the balls



(ii) The first surface in the hierarchy



(iv) The second surface in the hierarchy



(vi) A simplified copy of the pattern

So, 5_2 is non-trivial.

Low genus hierarchies

<u>Theorem</u>: Let M be a compact orientable irreducible 3-manifold with an essential boundary pattern P, and let \mathcal{H} be a handle structure for (M, P). Then (M, P) admits a hierarchy

$$(M,P) = (M_1,P_1) \xrightarrow{S_1} (M_2,P_2) \xrightarrow{S_2} \cdots \xrightarrow{S_n} (M_{n+1},P_{n+1})$$

and each (M_i, P_i) has a handle structure \mathcal{H}_i such that the following hold:

- 1. each S_i is normal and fundamental in \mathcal{H}_i ;
- 2. complexity(\mathcal{H}_{i+2}) < complexity(\mathcal{H}_i).

<u>Moreover</u>: When M is the exterior of a link L and $P = \emptyset$, and \mathcal{H} is the handle structure arising from a diagram D, then $|\chi(S_i)|$ and $|S_i \cap P_i|$ are both $O(c(D)^2)$, where c(D) is the crossing number of D.

The JSJ surfaces for a manifold with pattern

There is an analogue of the JSJ decomposition for a manifold M with essential pattern P.

The decomposing surfaces are tori, annuli (disjoint from P) and squares (ie discs intersecting P four times).

<u>Theorem</u>: The JSJ surfaces decompose (M, P) into the following pieces:

- 'simple' manifolds (ie they contain no essential tori, annuli disjoint from P, or squares);
- Seifert fibre manifolds,
- patterned *I*-bundles.

The *I*-bundles \mathcal{B} determine a transfer map $\tau : \partial_h \mathcal{B} \to \partial_h \mathcal{B}$.

Determining the JSJ by induction

Let $(M, P) \xrightarrow{S} (M', P')$ be a decomposition with gluing map ϕ .

Any essential torus/annulus/square in (M, P) is cut up by S into a torus/annuli/squares in (M', P').

So, the *I*-bundle pieces of the JSJ for (M, P) are obtained by patching together the *I*-bundle pieces \mathcal{B}' of the JSJ for (M', P'). But $\partial_h \mathcal{B}'$ might not patch together precisely under ϕ . So one shrinks \mathcal{B}' , forming a smaller *I*-bundle \mathcal{B}'_2 . Again, this might not patch together correctly under ϕ . So one shrinks \mathcal{B}'_2 , forming a smaller *I*-bundle \mathcal{B}'_3 , etc. This process stabilises after $O(|\chi(S)|)$ steps. So in this way we end with the JSJ for (M, P). (In fact, one keeps track of just its transfer map τ .) Note that this does not work when S is a fibre.

Completing the proof

We are given a diagram D for the non-fibred hyperbolic link L.

The hierarchy for the exterior of L is given to us as part of the certificate.

This shows that L is not the unknot and not split.

We compute the JSJ of the manifolds, working backwards along the hierarchy.

So if L is hyperbolic and non-fibred, we may certify that $S^3 - L$ has no JSJ tori and is not Seifert fibred.

In other words, we have a polynomial time certificate that L is hyperbolic.

Further questions

- What about general Haken 3-manifolds?
- What about general 3-manifolds?
- Finding the hyperbolic structure: is this in FNP?
- Do hyperbolic quantities (eg volume) relate to the combinatorics of an essential hierarchy?