# FINITE TYPE INVARIANTS OF STRING LINKS AND THE HOMFLYPT POLYNOMIAL OF KNOTS

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ABSTRACT. A family of finite type invariants of string link is given by the HOM-FLYPT polynomial of knots using various closure operations on (cabled) string links. In this note we will show the following:

(1) These invariants, together with Milnor invariants of length  $\leq 5$ , give classifications of *n*-string links up to  $C_k$ -equivalence for  $k \leq 5$ , and give a complete set of finite type invariants of degree  $\leq 4$ .

(2) Any Milnor invariant of length n + 1 > 2 of a  $C_n$ -trivial string link is expressed as a linear combination of such invariants.

## 1. String links and $C_k$ -moves

The notion of string link was introduced by Le Dimet [3] and Habegger-Lin [5]. A string link is a kind of tangle without closed components in the cylinder, which generalizes pure braids.

Let D be the unit disk in the plane. Choose n points  $p_1, ..., p_n$  in the interior of D so that  $p_1, ..., p_n$  lie in order on the x-axis, see Figure 2.1. An *n*-string link  $L = K_1 \cup \cdots \cup K_n$  in  $D \times [0,1]$  is a disjoint union of oriented arcs  $K_1, ..., K_n$  such that each  $K_i$  runs from  $(p_i, 0)$  to  $(p_i, 1)$  (i = 1, ..., n). The string link  $K_1 \cup \cdots \cup K_n$ with  $K_i = \{p_i\} \times [0, 1]$  (i = 1, ..., n) is called the *trivial n*-string link and denoted by  $\mathbf{1}_n$ .

The set  $S\mathcal{L}(n)$  of isotopy classes of *n*-string links fixing the endpoints has a monoidal structure, with composition given by the *stacking product* and with the trivial *n*-string link  $\mathbf{1}_n$  as unit element.

Habiro [6] and Goussarov [4] introduced independently the notion of  $C_k$ -move as follows. (This notion can also be defined by using the theory of claspers, see Subsection 5.1.) A  $C_k$ -move is a local move on (string) links as illustrated in Figure 1.1, which can be regarded as a kind of 'higher order crossing change' (in particular, a  $C_1$ -move is a crossing change). The  $C_k$ -move generates an equivalence relation on (string) links, called  $C_k$ -equivalence, which becomes finer as k increases. Thus we have a descending filtration

$$\mathcal{SL}(n) = \mathcal{SL}_1(n) \supset \mathcal{SL}_2(n) \supset \mathcal{SL}_3(n) \supset \dots$$

where  $\mathcal{SL}_k(n)$  denotes the set of  $C_k$ -trivial n-string links, i.e., string links which are  $C_k$ -equivalent to  $\mathbf{1}_n$ . For  $1 \leq k \leq l$ , let  $\mathcal{SL}_k(n)/C_l$  denote the set of  $C_l$ -equivalence



FIGURE 1.1. A  $C_k$ -move involves k + 1 strands of a link, labelled here by integers between 0 and k.

classes of  $C_k$ -trivial *n*-string links. This is known to be a finitely generated nilpotent group. Furthermore, if  $l \leq 2k$ , this group is abelian [6, Thm. 5.4].

# 2. FINITE TYPE INVARIANTS OF STRING LINKS

A singular n-string links is a proper immersion  $\bigsqcup_{i=1}^{n}[0,1]_{i} \rightarrow D^{2} \times [0,1]$  of the disjoint union  $\bigsqcup_{i=1}^{n}[0,1]_{i}$  of n copies of [0,1] in  $D^{2} \times [0,1]$  such that the image of  $[0,1]_{i}$  runs from  $(p_{i},0)$  to  $(p_{i},1)$   $(1 \leq i \leq n)$ , and whose singularities are transverse double points (in finite number).

Denote by  $\mathbf{ZSL}(n)$  the free abelian group generated by SL(n). A singular *n*-string link  $\sigma$  with k double points can be expressed as an element of  $\mathbf{ZSL}(n)$  using the following skein formula.

Let A be an abelian group. An n-string link invariant  $f : S\mathcal{L}(n) \to A$  is a finite type invariant of degree  $\leq k$  if its linear extension to  $\mathbb{ZSL}(n)$  vanishes on every n-string-link with (at least) k+1 double points. If f is of degree  $\leq k$  but not of degree k-1, then f is called a finite type invariant of degree k.

We recall a few classical examples of such invariants in the next two subsections.

2.1. Finite type knot invariants. Recall that the Conway polynomial of a knot K has the form

$$abla_K(z) = 1 + \sum_{k \ge 1} a_{2k}(K) z^{2k}.$$

It is not hard to show that the  $z^{2k}$ -coefficient  $a_{2k}$  in the Conway polynomial is a finite type invariant of degree 2k [1].

Recall also that the HOMFLYPT polynomial of a knot K is of the form

$$P(K;t,z) = \sum_{k=0}^{N} P_{2k}(K;t) z^{2k},$$

where  $P_{2k}(K;t) \in \mathbb{Z}[t^{\pm 1}]$  is called the 2kth coefficient polynomial of K. Denote by  $P_{2k}^{(l)}(K)$  the *l*th derivative of  $P_{2k}(K;t)$  evaluated at t = 1. It was proved by Kanenobu and Miyazawa that  $P_{2k}^{(l)}$  is a finite type invariant of degree 2k + l [9].

Note that both the Conway and HOMFLYPT polynomials of knots are invariant under orientation reversal, and that both are multiplicative under the connected sum [12]. 2.2. Milnor invariants of string links. Given an *n*-component oriented, ordered link L in  $S^3$ , Milnor invariants  $\overline{\mu}_L(I)$  of L are defined for each multi-index  $I = i_1 i_2 \dots i_m$  (i.e., any sequence of possibly repeating indices) among  $\{1, \dots, n\}$  [18, 19]. The number m is called the *length* of Milnor invariant  $\overline{\mu}(I)$ , and is denoted by |I|. Unfortunately, the definition of these  $\overline{\mu}(I)$  contains a rather intricate self-recurrent indeterminacy.

Habegger and Lin showed that Milnor invariants are actually well defined integervalued invariants of *string* links [5], and that the indeterminacy in Milnor invariants of a link is equivalent to the indeterminacy in regarding it as the closure of a string link.

In the unit disk  $D^2$ , we chose a point  $e \in \partial D$  and loops  $\alpha_1, ..., \alpha_n$  as illustrated in Figure 2.1. For an *n*-component string link  $L = K_1 \cup \cdots \cup K_n$  in  $D^2 \times [0,1]$  with  $\partial K_j = \{(p_j, 0), (p_j, 1)\}\ (j = 1, ..., n), \text{ set } Y = (D^2 \times [0, 1]) \setminus L, Y_0 = (D^2 \times \{0\}) \setminus L,$ and  $Y_1 = (D^2 \times \{1\}) \setminus L$ . We may assume that each  $\pi_1(Y_t)$   $(t \in \{0,1\})$  with base point (e, t) is the free group F(n) on generators  $\alpha_1, ..., \alpha_n$ . We denote the image of  $\alpha_j$ in the lower central series quotient  $F(n)/F(n)_q$  again by  $\alpha_j$ . By Stallings' theorem [23], the inclusions  $i_t : Y_t \longrightarrow Y$  induce isomorphisms  $(i_t)_* : \pi_1(Y_t)/\pi_1(Y_t)_q \longrightarrow$  $\pi_1(Y)/\pi_1(Y)_q$  for any positive integer q. Hence the induced map  $(i_1)_*^{-1} \circ (i_0)_*$  is an automorphism of  $F(n)/F(n)_q$  and sends each  $\alpha_j$  to a conjugate  $l_j\alpha_j l_j^{-1}$  of  $\alpha_j$ , where  $l_j$  is the longitude of  $K_j$  defined as follows. Let  $\gamma_j$  be a zero framed parallel of  $K_j$ such that the endpoints  $(c_j, t) \in D^2 \times \{t\}$  (t = 0, 1) lie on the x-axis in  $\mathbb{R}^2 \times \{t\}$ . The longitude  $l_j \in F(n)/F(n)_q$  is an element represented by the union of the arc  $\gamma_j$  and the segments  $e \times [0,1]$ ,  $c_j e \times \{0,1\}$  under  $(i_1)_*^{-1}$ . The coefficient  $\mu_L(i_1 i_2 \dots i_{k-1} j)$   $(k \leq q)$ of  $X_{i_1} \cdots X_{i_{k-1}}$  in the Magnus expansion  $E(l_j)$  is well-defined invariant of L, and it is called a *Milnor*  $\mu$ -invariant of length k. It is known that Milnor  $\mu$ -invariants of length k are finite type invariants of degree k-1 for string links [2, 13].



FIGURE 2.1

Convention 2.1. As said above, each Milnor invariant  $\mu(I)$  for *n*-string links is indexed by a sequence I of possibly repeating integers in  $\{1, ..., n\}$ . In the following, when denoting indices of Milnor invariants, we will always let distinct letters denote distinct integers, unless otherwise specified. For example,  $\mu(iijk)$   $(1 \le i, j, k \le n)$  stands for all Milnor invariants  $\mu(iijk)$  with  $i, j, k \in \{1, ..., n\}$  pairwise distincts. 2.3. Closure invariants. Given an *n*-string link L and a sequence  $I = i_1 i_2 \cdots i_m$  of m elements in  $\{1, ..., n, \overline{1}, ..., \overline{n}\}$ , we will construct in the next subsection an oriented knot K(L; I) in  $S^3$  as a closure of L with respect to I. Roughly speaking, we build the knot in  $S^3$  by connecting the endpoints of the  $i_j$ th components of L (j = 1, ..., m) so that, when running along the knot following the orientation, we meet these components in the order  $i_1, i_2, \cdots, i_m$ . Indices contained in  $\{1, ..., n\}$ , resp. in  $\{\overline{1}, ..., \overline{n}\}$ , correspond to components whose orientation agree, resp. disagree, with the orientation of the knot. If some index appears more than one in I, then we properly take parallels of the corresponding component of L.

2.3.1. Definition of the knot K(L; I) for a sequence I without repetition. Let  $I = i_1 i_2 \cdots i_m$  be a sequence of m elements in  $\{1, ..., n, \overline{1}, ..., \overline{n}\}$  without repeated number, i.e., for each i = 1, ..., n, the number of times that i or  $\overline{i}$  appears in I is at most one. Let  $L = K_1 \cup ... \cup K_n$  be an n-string link in  $D^2 \times [0, 1] \subset S^3$ .

Suppose that  $\partial K_i = p_i \times \{0,1\} \subset D^2 \times \{0,1\}$ . For each *I*, we choose a tangle  $T_I$  in  $S^3 \setminus (D^2 \times [0,1])$  as follows:

- If  $i_k$  and  $i_{k+1}$  are in  $\{1, ..., n\}$  then connect  $p_{i_k} \times \{1\}$  and  $p_{i_{k+1}} \times \{0\}$  in  $S^3 \setminus (D^2 \times [0, 1])$ .
- If  $i_k$  is in  $\{1, ..., n\}$  and  $i_{k+1}$  is in  $\{\overline{1}, ..., \overline{n}\}$  then connect  $p_{i_k} \times \{1\}$  and  $p_{\overline{i_{k+1}}} \times \{1\}$  in  $S^3 \setminus (D^2 \times [0, 1])$ .
- If  $i_k$  and  $i_{k+1}$  are in  $\{\overline{1}, ..., \overline{n}\}$  then connect  $p_{\overline{i_k}} \times \{0\}$  and  $p_{\overline{i_{k+1}}} \times \{1\}$  in  $S^3 \setminus (D^2 \times [0, 1])$ .
- If  $i_k$  is in  $\{\overline{1}, ..., \overline{n}\}$  and  $i_{k+1}$  is in  $\{1, ..., n\}$  then connect  $p_{\overline{i_k}} \times \{0\}$  and  $p_{i_{k+1}} \times \{0\}$  in  $S^3 \setminus (D^2 \times [0, 1])$ .

Here we implicitely mean that  $\overline{\overline{i}} = i$  and  $i_{m+1} = i_1$  in our notation. Let  $L_I$  be the *m*-string link obtained from *L* by deleting all components  $K_j$  of *L* such that neither *j* nor  $\overline{j}$  appears in *I*. Then we have a knot  $K(L; I) := L_I \cup T_I$  in  $S^3$ . See Figure 2.2 for an example. For each *I*, we choose  $T_I$  so that  $K(\mathbf{1}_n; I)$  is the trivial knot. While there are several choices of  $T_I$  tangles for each *I*, we choose one and fix it.



FIGURE 2.2

2.3.2. Definition of the knot K(L; I) for an arbitrary sequence I. Let  $L = K_1 \cup \cdots \cup K_n$  be an *n*-string link. Let  $I = i_1 i_2 \cdots i_m$  be a sequence of m elements of  $\{1, ..., n, \overline{1}, ..., \overline{n}\}$ , where for each number i(=1, ..., n), the number of times that i or  $\overline{i}$  appears in I is  $r_i$ . Let  $m = \sum_i r_i$ . Denote by  $D_I(L)$  the *m*-string link obtained from L as follows:

- Replace each string  $K_i$  by  $r_i$  zero-framed parallel copies of it, labeled from  $K_{(i,1)}$  to  $K_{(i,r_i)}$  according to the natural order induced by the orientation of the diametral axis in  $D^2$ . If  $r_i = 0$  for some index *i*, simply delete  $K_i$ .
- Let  $D_I(L) = K'_1 \cup \cdots \cup K'_m$  be the *m*-string link  $\bigcup_{i,j} K_{(i,j)}$  with the order induced by the lexicographic order of the index (i, j). This ordering defines a bijection

$$\varphi: \{(i,j) \mid 1 \le i \le n, 1 \le j \le r_i\} \to \{1, ..., m\}.$$

We also define a sequence D(I) of elements of  $\{1, ..., m\}$  without repeated number as follows. First, consider a sequence of elements of  $\{(i, j); 1 \leq i \leq n, 1 \leq j \leq r_i\}$ by replacing each number *i* in *I* with  $(i, 1), ..., (i, r_i)$  in this order. For example if  $I = 12\overline{2}31$ , we obtain the sequence  $(1, 1), (2, 1), (\overline{2}, 2), (3, 1), (1, 2)$ . Next replace each term (i, j) of this sequence with  $\varphi((i, j))$ . Hence we have  $D(12\overline{2}31) = 13\overline{4}52$ . Since D(I) does not contain repeated number, we have a closure  $K(D_I(L); D(I))$  of *L* with respect to the sequence D(I). We call the knot  $K(D_I(L); D(I))$  the closure knot with respect to *I*.

It is not hard to show the following proposition.

**Proposition 2.2.** Let I be a sequence of elements in  $\{1, ..., n, \overline{1}, ..., \overline{n}\}$ , and let  $v_m$  be a finite type invariant of degree m for knots. Then the assignment  $L \mapsto v_m(K(D_I(L); D(I)))$  defines a finite type invariant of degree m for n-string links.

Convention 2.3. Let  $v_m$  be a finite type invariant of degree  $m \geq 2$  for knots. For an *n*-string link *L* and a sequence  $I = i_1 i_2 \cdots i_m$  of *m* elements of  $\{1, ..., n, \overline{1}, ..., \overline{n}\}$ , we denote  $v_m(K(D_I(L); D(I)))$  by  $v_m(L; I)$  or  $v_m(D_I(L); D(I))$ . For example, we denote  $P_0^{(m)}(K(D_I(L); D(I)))$  and  $a_m(K(D_I(L); D(I)))$  by  $P_0^{(m)}(L; I)$  and  $a_m(L; I)$ respectively. We call  $P_0^{(m)}(L; I)$  and  $a_m(L; I)$  the  $P_0^{(m)}$ -closure invariant and the  $a_m$ -closure invariant respectively.

# 3. $C_k$ -moves and finite type invariants

3.1. The Goussarov-Habiro Conjecture. Goussarov and Habiro showed independently the following.

**Theorem 3.1** ([4, 6]). Two knots (1-string links) cannot be distinguished by any finite type invariant of degree  $\leq k$  if and only if they are  $C_k$ -equivalent.

It is known that the 'if' part of the statement holds for links as well, but explicit examples show that the 'only if' part of Theorem 3.1 does not hold for links in general, see  $[6, \S7.2]$ .

However, Theorem 3.1 may generalize to string links.

**Conjecture** (Goussarov-Habiro; [4, 6]). Two string links of the same number of components share all finite type invariant of degree  $\leq k - 1$  if and only if they are  $C_k$ -equivalent.

As in the link case, the 'if' part of the conjecture is always true. The 'only if' part is also true for k = 1 (in which case the statement is vacuous) and k = 2; the only finite type string link invariant of degree 1 is the linking number, which is known to classify

string links up to  $C_2$ -equivalence [21]. (Note that this actually also applies to links). The Goussarov-Habiro conjecture was (essentially) proved for k = 3 by the first author in [15]. Massuyeau gave a proof for k = 4, but it is mostly based on algebraic arguments and thus does not provide any information about the corresponding finite type invariants [14]. In [16], we classify *n*-string links up to  $C_k$ -move for  $k \leq 5$ , by explicitly giving a complete set of low degree finite type invariants. In addition to Milnor invariants, these include several closure invariants of string links. In the next subsection, we give the statements of these results. As a consequence, we show that the Goussarov-Habiro Conjecture is true for  $k \leq 5$ .

3.2. Invariants of degree  $\leq 4$ . In this subsection, we give a  $C_k$ -classification of string links for  $k \leq 5$ . While the statements here look different from the statements in [16], they are essentially the same (we just use a different notation for closure invariants).

Recall that there is essentially only one finite type knot invariant of degree 2, namely  $a_2$ , and that there is essentially only one finite type knot invariant of degree 3, namely  $P_0^{(3)}$ . There are essentially two linearly independent finite type knot invariants of degree 4, namely  $a_4$  and  $P_0^{(4)}$ . We will use these knot invariants to define a number of finite type string links invariants of degree  $\leq 4$  by using some closure. These various invariants, together with Milnor invariants of length  $\leq 5$ , give the following classification of *n*-string links up to  $C_k$ -equivalence for  $k \leq 5$ .

**Theorem 3.2** ([15]). Let  $L, L' \in SL(n)$ . Then the following assertions are mutually equivalent:

- (1) L and L' are  $C_3$ -equivalent,
- (2) L and L' share all finite type invariants of degree  $\leq 2$ ,
- (3)  $a_2(L;i) = a_2(L';i) \ (1 \le i \le n),$  $a_2(L;i\overline{j}) = a_2(L';i\overline{j}) \ (1 \le i < j \le n),$  $\mu_L(ij) = \mu_{L'}(ij) \ (1 \le i < j \le n) \ and$  $\mu_L(ijk) = \mu_{L'}(ijk) \ (1 \le i < j < k \le n).$

**Theorem 3.3** ([16]). Let  $L, L' \in SL(n)$ . Then the following assertions are mutually equivalent:

- (1) L and L' are  $C_4$ -equivalent,
- (2) L and L' share all finite type invariants of degree  $\leq 3$ ,
- (3) L and L' share all finite type invariants of degree  $\leq 2$ , and  $P_0^{(3)}(L;i) = P_0^{(3)}(L';i) \ (1 \leq i \leq n),$   $P_0^{(3)}(L;i\overline{j}) = P_0^{(3)}(L';i\overline{j}) \ (1 \leq i < j \leq n),$   $P_0^{(3)}(L;ik\overline{j}) = P_0^{(3)}(L';ik\overline{j}) \ (1 \leq i < j < k \leq n),$   $\mu_L(ijj) = \mu_{L'}(iijj) \ (1 \leq i < j \leq n),$   $\mu_L(ijkl) = \mu_{L'}(ijkl) \ (1 \leq i, j < k < l \leq n) \ and$  $\mu_L(ijkk) = \mu_{L'}(ijkk) \ (1 \leq i, j, k \leq n, \ i < j).$

**Theorem 3.4** ([16]). Let  $L, L' \in S\mathcal{L}(n)$ . Then the following assertions are equivalent:

(1) L and L' are  $C_5$ -equivalent,

(2) L and L' share all finite type invariants of degree  $\leq 4$ ,

3) L and L' share all finite type invariants of degree 
$$\leq 3$$
, and  
 $a_4(L;i) = a_4(L';i), P_0^{(4)}(L;i) = P_0^{(4)}(L';i) (1 \leq i \leq n),$   
 $a_4(L;i\overline{j}) = a_4(L';i\overline{j}), P_0^{(4)}(L;i\overline{j}) = P_0^{(4)}(L';i\overline{j}),$   
 $a_4(L;i\overline{j}) = a_4(L';ii\overline{j}), P_0^{(4)}(L;i\overline{j}) = P_0^{(4)}(L';i\overline{j}),$   
 $P_0^{(4)}(K(L;i\overline{j}\overline{j})) = P_0^{(4)}(K(L';i\overline{j}\overline{j})) (1 \leq i < j \leq n),$   
 $a_4(L;i\overline{k}) = a_4(L';i\overline{k}), P_0^{(4)}(L;i\overline{k}) = P_0^{(4)}(L';i\overline{j}\overline{k}),$   
 $a_4(L;i\overline{k}) = a_4(L';i\overline{k}), P_0^{(4)}(L;i\overline{k}) = P_0^{(4)}(L';i\overline{k}),$   
 $a_4(L;i\overline{k}) = a_4(L';i\overline{k}), P_0^{(4)}(L;i\overline{k}) = P_0^{(4)}(L';i\overline{k}),$   
 $a_4(L;i\overline{k}) = a_4(L';i\overline{k}), P_0^{(4)}(L;i\overline{k}) = P_0^{(4)}(L';i\overline{k}),$   
 $P_0^{(4)}(L;i\overline{j}\overline{k}) = P_0^{(4)}(L';i\overline{j}\overline{k}), P_0^{(4)}(L;i\overline{j}\overline{k}) = P_0^{(4)}(L';i\overline{j}\overline{k}),$   
 $P_0^{(4)}(L;i\overline{k}\overline{j}\overline{l}) = P_0^{(4)}(L';i\overline{k}\overline{j}) (1 \leq i < j < k \leq n),$   
 $P_0^{(4)}(L;i\overline{k}\overline{j}\overline{l}) = P_0^{(4)}(L';i\overline{k}\overline{j}\overline{l}) (1 \leq i < j < k < l \leq n),$   
 $\mu_L(ijklm) = \mu_{L'}(ijklm) (1 \leq i, j, k < l < m \leq n),$   
 $\mu_L(ijkll) = \mu_{L'}(jikll) (1 \leq i, j, k, l \leq n, j < k).$ 

Remark 3.5. A complete set of finite type link invariant of degree  $\leq 3$  has been computed in [10] using weight systems and chord diagrams. For 2-component links, this has been done for degree  $\leq 4$  invariants in [11]. All invariants are given by coefficients of the Conway and HOMFLYPT polynomials of sublinks.

# 4. Milnor invariants and $P_0^{(m)}$ -closure invariants

We start by expressing Milnor's link homotopy invariants, i.e., Milnor invariants  $\mu(I)$  with a sequence I without repeated number, in terms of the closure invariants defined in Subsection 2.3.

**Theorem 4.1** ([17]). Let  $m \ge 2$ . Let L be a  $C_m$ -trivial n-string link  $(m + 1 \le n)$ . Let I be a sequence of m + 1 elements of  $\{1, ..., n\}$  without repeated number. Then

$$\mu_L(I) = \frac{\pm 1}{m! 2^m} \sum_{J \subset I, J \neq \emptyset} (-1)^{m-|J|} P_0^{(m)}(L; J),$$

where the sum runs over all nonempty subsequences J of I.

Remark 4.2. (1) By [6], the fact that L is  $C_m$ -trivial implies that  $\mu_L(I) = 0$  for any sequence I of length  $|I| \leq m$ .

(2) Any link-homotopically trivial Brunnian *n*-string link is  $C_n$ -trivial [8, 20], and any Brunnian *n*-string link whose Milnor invariants of length  $\leq n+1$  vanish is  $C_{n+1}$ trivial [16]. Since a Brunnian *n*-string link whose Milnor invariants with length  $\leq n$ vanish is link-homotopically trivial [18], for m = n + 1 or n, a Brunnian *n*-string link whose Milnor invariants with length  $\leq m$  vanish is  $C_m$ -trivial. Moreover, any Brunnian *n*-string link is  $C_{n-1}$ -trivial [7, 20] and has vanishing Milnor invariants with length  $\leq n - 1$ , so this holds for m = n - 1 as well.

(3) Since there exists no degree one invariant of knots, such a formula does not hold for the linking number, hence the assumption  $m \ge 2$  is needed. In order to give such

a formula one should consider 'closure links', that is more general closure operations on string links that can produce links with several components.

By combining [19, Thm. 7] and Theorem 4.1, we have the following theorem.

**Theorem 4.3** ([17]). Let  $m \ge 2$ . Let L be a  $C_m$ -trivial n-string link. Let I be a sequence of m + 1 elements of  $\{1, ..., n\}$ . Then

$$\mu_L(I) = \mu_{D_I(L)}(D(I)) = \frac{\pm 1}{m! 2^m} \sum_{J \subset D(I), J \neq \emptyset} (-1)^{m-|J|} P_0^{(m)}(D_I(L); J),$$

where the sum runs over all nonempty subsequences J of D(I).

K. Habiro has pointed out the following remark.

Remark 4.4. It is not hard to see that the 6-string link L illustrated in Figure 4.1 is  $C_5$ -trivial and satisfies  $\mu_L(123456) = \pm 1$ . By Theorem 4.1,  $\mu_L(123456)$  can be expressed as a linear combination of  $P_0^{(5)}$ -closure invariants of L. (By applying the theorem, we have  $\mu_L(123456) = (\pm 1/5!2^5)P_0^{(5)}(L; 123456)$ .) In contrast, since  $a_5$  of knots always vanish, it is impossible to express  $\mu_L(123456)$  by any linear combination of  $a_5$ -closure invariants of L. Moreover we notice that L is equivalent to  $\mathbf{1}_6$  up to doubled-delta move, which is a local move on links defined by Naik and Stanford [22]. Hence any closure knot, and more generally any closure link (see Remark 4.2(3)) obtained from L is equivalent to a trivial knot or link up to doubled-delta moves. Since the doubled-delta move preserves the Alexander invariant, the Conway polynomial of any closure link obtained from L vanishes.



FIGURE 4.1

5. Claspers and  $P_0^{(m)}$ -closure invariants

5.1. Claspers. For a general definition of claspers, we refer the reader to [6]. Let L be a (string) link. A surface G embedded in  $D^2 \times (0, 1)$  is called a graph clasper for L if it satisfies the following three conditions:

- (1) G is decomposed into disks and bands, called *edges*, each of which connects two distinct disks.
- (2) The disks have either 1 or 3 incident edges, and are called *leaves* or *nodes* respectively.
- (3) G intersects L transversely, and the intersections are contained in the union of the interiors of the leaves.

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In particular, if a connected graph clasper G is simply connected, we call it a *tree* clasper.

A graph clasper for a (string) link L is simple if each of its leaves intersects L at one point. The degree of a connected graph clasper G is defined as half of the number of nodes and leaves. We call a degree k connected graph clasper a  $C_k$ -graph. A tree clasper of degree k is called a  $C_k$ -tree.

Given a graph clasper G for a (string) link L, there is a procedure to construct a framed link, in a regular neighbourhood of G. There is thus a notion of surgery along G, which is defined as surgery along the corresponding framed link. In particular, surgery along a simple  $C_k$ -tree is a local move as illustrated in Figure 5.1, which is equivalent to a  $C_k$ -move as defined in Section 1 (Figure 1.1).



FIGURE 5.1. Surgery along a simple  $C_5$ -tree.

The  $C_k$ -equivalence (as defined in Section 1) coincides with the equivalence relation on string links generated by surgeries along  $C_k$ -graphs and isotopies. In particular, it is known that two links are  $C_k$ -equivalent if and only if they are related by surgery along simple  $C_k$ -trees [6, Thm. 3.17].

For  $k \geq 3$ , a  $C_k$ -tree G having the shape of the tree clasper in Figure 5.1 is called a *linear*  $C_k$ -tree. The left-most and right-most leaves of G in Figure 5.1 are called the *ends* of G, and the remaining (k-1) leaves are called the *internal leaves* of G.

Suppose that the two ends of a linear  $C_k$ -tree are denoted by f and f'. Let S be a nonempty subset of the set of all internal leaves of T. We have a labeling from 1 to |S| of the leaves in S by travelling along the boundary of the disk<sup>1</sup> T from f to f' so that all leaves are visited. We call this labeling the *linear labeling of* S, from f to f'.

5.2. Generators of  $\mathcal{SL}_m(n)/C_{m+1}$ . Let  $m \geq 3$  be an integer. In this section we find generators for the abelian group  $\mathcal{SL}_m(n)/C_{m+1}$  and show that for each of these generators, there is a  $P_0^{(m)}$ -closure invariant which detects it.

For a simple tree clasper  $\Gamma$  for a string link, let  $r_i(\Gamma)$  denote the number of leaves intersecting the *i*th component of the string link.

Let  $L \in S\mathcal{L}_m(n)$  be a  $C_m$ -trivial *n*-string link. By Calculus of Claspers [16, Lem.3.2] and the AS and IHX relations [16, Lem.3.3], L is  $C_{m+1}$ -equivalent to a product  $\prod T_i$  of *n*-string links  $T_1, ..., T_l$ , where each  $T_k$  is obtained from  $\mathbf{1}_n$  by surgery along a simple linear  $C_m$ -tree  $\Gamma_k$ . Actually, by the IHX relation we may assume that each  $\Gamma_k$  satisfies one of the following;

- (1) all leaves of  $\Gamma_k$  intersect a single component of  $\mathbf{1}_n$ ,
- (2)  $|\{i \mid r_i(\Gamma_k) = 1\}| \ge 2$ , and the ends intersect the *p*th and *q*th components of  $\mathbf{1}_n$ , where  $p = \min\{i \mid r_i(\Gamma_k) = 1\}$  and  $q = \min\{i \mid r_i(\Gamma_k) = 1, i \neq p\}$ ,

<sup>&</sup>lt;sup>1</sup>Recall that a clasper is an embedded surface: in particular, since T is a tree clasper, the underlying surface is isotopic to a disk.

- (3)  $r_i(\Gamma_k) = 2$  for some i,  $|\{i \mid r_i(\Gamma_k) = 1\}| < 2$ , and the ends intersect the pth component of  $\mathbf{1}_n$ , where  $p = \min\{i \mid r_i(\Gamma_k) = 2\}$ ,
- (4)  $\Gamma_k$  is not of type (1),  $r_i(\Gamma_k) \neq 2$  for any i,  $|\{i \mid r_i(\Gamma_k) = 1\}| < 2$ , and the ends intersect the *p*th component of  $\mathbf{1}_n$ , where  $(r_p(\Gamma_k), p)$  is the minimum among  $\{(r_i(\Gamma_k), i) \mid i = 1, ..., n, r_i(\Gamma_k) \geq 3\}$  with respect to the lexicographic order.

This implies that  $\mathcal{SL}_m(n)/C_{m+1}$  is generated by all string links obtained from  $\mathbf{1}_n$  by surgery along a  $C_m$ -tree of one of the 4 types above.

Let us reduce the number of generators of type (4). Let  $\mathcal{T}_p$  be the set of linear  $C_m$ -trees of type (4) with ends intersecting the *p*th component of  $\mathbf{1}_n$ . Each tree in  $\mathcal{T}_n$  has a unique leaf not intersecting the *n*th component of  $\mathbf{1}_n$ . By [16, Lem.3.6], the case reduces to trees of type (3). Hence we may assume that  $p \neq n$ . By the IHX relation, we may assume that the two ends are the 'top', resp. 'bottom', leaves on the pth component of  $\mathbf{1}_n$ , which are defined as the last, resp. first, leaf we meet while traveling along this component from the initial point to the terminal point. For a  $C_m$ -tree  $\Gamma \in \mathcal{T}_p$  with top end f and bottom end f', we consider the linear labeling (from 1 to m-1) of the set of all internal leaves of  $\Gamma$ , from f' to f (see Section 5.1). Suppose that while traveling along the *p*th component from f in this order. We say that  $\Gamma$  is flat (on the pth component of  $\mathbf{1}_n$ ) if  $i_1 < i_2 < \cdots < i_s$ . Let  $\mathcal{F}_p$  be the set of flat trees in  $\mathcal{T}_p$ .

Define  $\mathcal{F}_p^0$  as set of  $C_k$ -trees in  $\mathcal{F}_p$  which do not contain a *fork*. Here we say that a tree clasper T for  $\mathbf{1}_n$  contains a fork if there exists a 3-ball that intersects  $\mathbf{1}_n \cup T$  as represented in Figure 5.2



FIGURE 5.2

**Proposition 5.1** ([17]). For an integer  $m \ge 3$ ,  $SL_m(n)/C_{m+1}$  is generated by string links obtained from  $\mathbf{1}_n$  by surgery along linear trees of type (1), (2), (3) or in  $\mathcal{F}_p^0$  (p = 1, ..., n-1).

The abelian group  $\mathcal{SL}_m(n)/C_{m+1}$  can be decomposed into a direct sum  $G_1 \oplus G_2$ , where  $G_1$  (resp.  $G_2$ ) is the subgroup generated by string links obtained from  $\mathbf{1}_n$ surgery along a linear  $C_m$ -tree of type (1) (resp. of type (2), (3) or in  $\mathcal{F}_p^0$  (p = 1, ..., n-1)). By the Goussarov-Habiro Theorem [4, 6],  $G_1$  is classified by finite type invariants. For the group  $G_2$ , we have the following

**Theorem 5.2** ([17]). Let  $m \geq 3$  be an integer. For any simple linear  $C_m$ -tree  $\Gamma$  for  $\mathbf{1}_n$  of type (2), (3) or in  $\mathcal{F}_p^0$  (p = 1, ..., n - 1), there is a sequence I of elements of  $\{1, ..., n\}$  such that  $P_0^{(m)}((\mathbf{1}_n)_{\Gamma}; I) = \pm m! 2^m$ . Hence  $(\mathbf{1}_n)_{\Gamma}$  has infinite order in  $\mathcal{SL}_m(n)/C_{m+1}$ .

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