

# Some Binary Minimal Clones on a Finite Set

Hajime Machida\*

Tokyo, Japan

## 1 Introduction

According to Type Theorem of I. G. Rosenberg ([Ro86]), minimal functions, i.e., generators of minimal clones with minimal arity, on a finite set are classified into five types.

- |                               |                                |
|-------------------------------|--------------------------------|
| (1) unary function            | (2) binary idempotent function |
| (3) ternary majority function | (4) ternary minority function  |
| (5) semiprojection            |                                |

While the minimal functions of type (1) and type (4) have been completely determined, the problem of characterizing minimal functions of types (2), (3) and (5) still remain unsolved. In fact, the problem of determining all minimal clones is considered to be one of the most difficult problems in clone theory.

In ([BM20]), we considered some specific kind of minimal functions of type (2); binary idempotent functions. This article is a brief report of the results obtained there.

Behrisch, M. and Machida, H., “On minimality of some binary clones related to unary functions”, to appear in *Proceedings 50th International Symposium on Multiple-Valued Logic*, IEEE, 2020.

This paper will appear in a symposium proceedings, which is not easily accessible to those who do not attend the symposium. So, it would be of some value to present the results here.

## 2 Preliminaries

For  $k > 1$ , let  $E_k = \{0, 1, \dots, k-1\}$ . Denote by  $\mathcal{O}_k^{(n)}$ , for any  $n > 0$ , the set of  $n$ -variable functions on  $E_k$  and by  $\mathcal{O}_k$  the set of all functions on  $E_k$ , i.e.,  $\mathcal{O}_k = \bigcup_{n>0} \mathcal{O}_k^{(n)}$ . The  $n$ -variable  $i$ -th projection  $e_i^n$ ,  $1 \leq i \leq n$ , is the function in  $\mathcal{O}_k^{(n)}$  defined by  $e_i^n(x_1, \dots, x_n) = x_i$  for all  $x_1, \dots, x_n \in E_k$ . Denote by  $\mathcal{J}_k$  the set of projections on  $E_k$ .

A subset  $C$  of  $\mathcal{O}_k$  is a clone on  $E_k$  if  $C$  contains all the projections, i.e.,  $\mathcal{J}_k \subseteq C$ , and is closed under (functional) composition. The set  $\mathcal{L}_k$  of clones on  $E_k$  forms a lattice with respect to inclusion. It is well known that the cardinality of  $\mathcal{L}_k$  is of continuum for every  $k \geq 3$ .

---

\*machida.zauber@gmail.com

For a clone  $C \in \mathcal{L}_k$  and a subset  $F$  of  $C$ ,  $F$  is said to *generate*  $C$  if  $C$  is the smallest clone containing  $F$ . When  $F$  generates  $C$ , we write  $C = \langle F \rangle$ . If  $F$  is a singleton, i.e.,  $F = \{f\}$ , we simply write  $\langle f \rangle$  instead of  $\langle F \rangle$ .

A clone  $C$  is called a *binary clone* if it is generated by a binary function, i.e., if  $C = \langle f \rangle$  for some  $f \in \mathcal{O}_k^{(2)}$ .

A clone  $C$  in  $\mathcal{L}_k$  is a *minimal clone* if it is an atom of  $\mathcal{L}_k$ . Equivalently,  $C (\neq \mathcal{J}_k)$  is a minimal clone if  $\mathcal{J}_k \subset C' \subseteq C$  implies  $C' = C$  for any  $C'$  in  $\mathcal{L}_k$ . A minimal clone  $C$  is generated by a single function, i.e.,  $C = \langle f \rangle$  for some  $f$  in  $\mathcal{O}_k \setminus \mathcal{J}_k$ . A function  $f$  is a *minimal function* if it generates a minimal clone  $C$  and its arity is minimum among functions generating  $C$ .

As stated in Section 1, a minimal function is of one of the five types, and minimal functions of type (1) and (4) have been completely characterized. The case of type (1) is easy; a unary function  $s$  is minimal if and only if either it is a permutation of prime order or a non-surjective function satisfying  $s^2 = s$ . The case of type (4) was solved in [Ro86]; a minority function  $m$  is minimal if and only if  $k = 2^r$  for some  $r > 0$  and  $m(x, y, z) = x + y + z$  for all  $x, y, z \in E_k$  where  $(E_k^r; +)$  is the elementary 2-group. (The operation  $+$  is the component-wise addition over  $\text{GF}(2)$ .)

For  $k = 2$ , i.e., the case of Boolean functions, there are 7 minimal clones, which are easily obtained from the Post lattice ([Po41]). For  $k = 3$ , it is known ([Cs83]) that there are 84 minimal clones. Among them, 48 are generated by binary idempotent functions.

In this article, we consider binary minimal clones, i.e., minimal clones generated by binary idempotent functions. A binary function  $f \in \mathcal{O}_k^{(2)}$  is *idempotent* if  $f(x, x) = x$  for all  $x \in E_k$ .

A unary function  $s$  on  $E_k$  naturally induces a directed graph  $(V, A)$  where  $V = E_k$  and  $A = \{(x, s(x)) \mid x \in E_k\}$ . This graph will be denoted by  $\Gamma(s)$ . For every vertex  $x \in V$ , either  $x$  is in a cycle or there is a simple path connecting  $x$  and some vertex in some cycle.

### 3 Two Classes of Binary Minimal Functions

Two particular classes  $\mathcal{F}_1$  and  $\mathcal{F}_2$  of binary idempotent functions are introduced, and we shall consider minimal functions in these classes.

A unary function  $s \in \mathcal{O}_k^{(1)}$  is *reflective* (or *retractive*) if it satisfies  $s^2 (= s \circ s) = s$ , equivalently,  $s$  is the identity on  $\text{Im}(s)$ . (Note: A more standard term may be idempotent which, however, is used for another meaning in this article.)

#### 3.1 The Class $\mathcal{F}_1$

For a unary function  $s \in \mathcal{O}_k^{(1)}$ , we define a binary function  $\varphi_s \in \mathcal{O}_k^{(2)}$  by

$$\varphi_s(x, y) = \begin{cases} s(y) & \text{if } x = 0, \\ x & \text{if } x \in E_k \setminus \{0\}. \end{cases}$$

When  $s$  is 0-preserving, i.e.,  $s(0) = 0$ ,  $\varphi_s$  is idempotent. The Cayley table of  $\varphi_s$  is shown below.

$x \setminus y$	0	1	2	$\dots$	$k-1$
0	0	$s(1)$	$s(2)$	$\dots$	$s(k-1)$
1	1	1	1	$\dots$	1
2	2	2	2	$\dots$	2
$\vdots$	$\vdots$	$\vdots$	$\vdots$		$\vdots$
$k-1$	$k-1$	$k-1$	$k-1$	$\dots$	$k-1$

Now,  $\mathcal{F}_1$  is defined to be the set of binary idempotent functions  $\varphi_s$  with 0-preserving unary functions  $s$ , i.e.,

$$\mathcal{F}_1 = \{ \varphi_s \in \mathcal{O}_k^{(2)} \mid s \in \mathcal{O}_k^{(1)}, s(0) = 0 \}.$$

For any  $f \in \mathcal{F}_1$ , if  $f = \varphi_s$  we denote  $s$  by  $s_f$ . In other words,  $s_f$  is the unary function satisfying  $s_f(y) = f(0, y)$  for all  $y \in E_k$ .

The following characterizes the minimal functions in  $\mathcal{F}_1$ . Here,  $c_0 (\in \mathcal{O}_k^{(1)})$  denotes the unary constant function taking the value 0.

**Proposition 3.1** *For any  $f \in \mathcal{F}_1$ ,  $f$  is a minimal function if and only if either of the following conditions is satisfied.*

- (1)  $s_f$  is reflective, i.e.,  $s_f^2 = s_f$ , and  $s_f \neq c_0$ ,
- (2)  $s_f^2 = c_0$  and  $s_f \neq c_0$ .

Let us describe the conditions (1) and (2) in terms of the associated graph  $\Gamma(s_f)$  of  $s_f$ . Let  $f \in \mathcal{F}_1$  be a function satisfying (1) or (2). Since  $\Gamma(s_f)$  has a loop leaving and entering the vertex 0, we have  $\#\text{loop}(\Gamma(s_f)) \geq 1$ , where  $\#\text{loop}(\Gamma(s_f))$  denotes the number of loops in  $\Gamma(s_f)$ . Moreover, it is easy to see that  $\Gamma(s_f)$  is cycle-free, except loops. Define the *height* of  $\Gamma(s_f)$ , denoted by  $\text{height}(\Gamma(s_f))$ , as the length of a longest simple path connecting some vertex and some loop-vertex.

The conditions given in Proposition 3.1 can be expressed with respect to the graph  $\Gamma(s_f)$  as follows.

- (1)  $\Gamma(s_f)$  has no cycles except loops,  $\#\text{loop}(\Gamma(s_f)) \geq 2$  and  $\text{height}(\Gamma(s_f)) \leq 1$ .
- (2)  $\Gamma(s_f)$  has no cycles except loops,  $\#\text{loop}(\Gamma(s_f)) = 1$  and  $\text{height}(\Gamma(s_f)) = 2$ .

Examples of graphs satisfying these conditions are depicted in Figure 1.

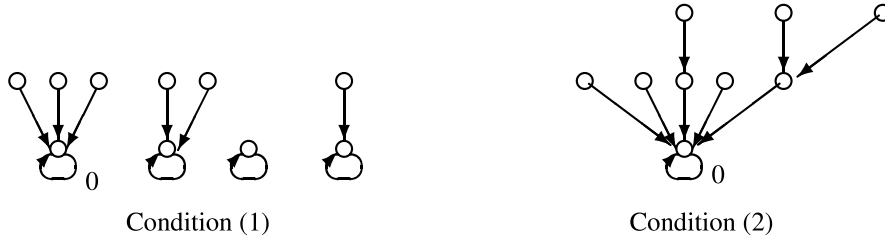


Figure 1: Examples of  $\Gamma(s_f)$  for Conditions (1) and (2)

## 4 The Class $\mathcal{F}_2$

For a unary function  $s \in \mathcal{O}_k^{(1)}$  we define a binary function  $\psi_s \in \mathcal{O}_k^{(2)}$  by

$$\psi_s(x, y) = \begin{cases} s(y) & \text{if } (x, y) \in \{0\} \times E_k, \\ s(x) & \text{if } (x, y) \in E_k \times \{0\}, \\ \min\{x, y\} & \text{if } (x, y) \in (E_k \setminus \{0\})^2. \end{cases}$$

Furthermore, we assume  $s$  to be 0-preserving, i.e.,  $s(0) = 0$ . As  $\psi_s$  is commutative, the Cayley table of  $\psi_s$  is “symmetric” as shown below.

$x \setminus y$	0	1	2	3	...	$k-1$
0	0	$s(1)$	$s(2)$	$s(3)$	...	$s(k-1)$
1	$s(1)$	1	1	1	...	1
2	$s(2)$	1	2	2	...	2
3	$s(3)$	1	2	3	...	3
$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$		$\vdots$
$k-1$	$s(k-1)$	1	2	3	...	$k-1$

We define  $\mathcal{F}_2$  to be the collection of such binary idempotent commutative functions  $\psi_s$ .

$$\mathcal{F}_2 = \{ \psi_s \in \mathcal{O}_k^{(2)} \mid s \in \mathcal{O}_k^{(1)}, s(0) = 0 \}$$

For any  $f \in \mathcal{F}_2$ , the unary function  $s \in \mathcal{O}_k^{(1)}$  such that  $f = \psi_s$  will be denoted by  $s_f$ . In other words,  $s_f$  is a function which satisfies  $s_f(z) = f(0, z) (= f(z, 0))$  for all  $z \in E_k$ .

Assuming  $E_k$  as the initial segment of the set  $\mathbb{N}$  of natural numbers, we introduce the order  $\leq$  on  $E_k$  induced from the natural order on  $\mathbb{N}$ , i.e.,  $0 < 1 < 2 < \dots < k-1$  is assumed. With respect to this ordering, we say a unary function  $s \in \mathcal{O}_k^{(1)}$  is *intensional* if  $s(z) \leq z$  for all  $z \in E_k$ .

The main result for  $f \in \mathcal{F}_2$  is the following.

**Proposition 4.1** *For any  $f \in \mathcal{F}_2$ , if  $s_f$  is reflective and intensional, then  $f$  is a minimal function. On the other hand, if  $f$  is a minimal function, then  $s_f$  is intensional.*

As an example, take a unary function  $s \in \mathcal{O}_{10}^{(1)}$  given by the following table.

$x$	0	1	2	3	4	5	6	7	8	9
$s(x)$	0	0	0	3	4	0	3	7	7	3

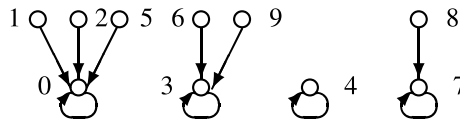


Figure 2: An example of  $\Gamma(s)$  for Proposition 4.1

The graph  $\Gamma(s)$  is shown in Figure 2. Clearly,  $s$  is reflective and intensional, implying that  $\psi_s \in \mathcal{O}_{10}^{(2)}$  is a minimal function.

**Remark** Let  $\widehat{\mathcal{F}}_2$  be the set of binary idempotent functions  $f$  satisfying  $f(x, y) = \min\{x, y\}$  for all  $(x, y) \in (E_k \setminus \{0\})^2$ . The values  $f(0, y)$  and  $f(x, 0)$  for  $x, y \in E_k \setminus \{0\}$  are arbitrary. Clearly,  $\mathcal{F}_2$  is a proper subset of  $\widehat{\mathcal{F}}_2$ . For each  $f \in \widehat{\mathcal{F}}_2$ , define unary functions  $r_f, c_f \in \mathcal{O}_k^{(1)}$  by

$$\begin{cases} r_f(y) = f(0, y) & \text{for all } y \in E_k, \\ c_f(x) = f(x, 0) & \text{for all } x \in E_k. \end{cases}$$

Let  $\widehat{\mathcal{F}}_2^\nabla$  be the subset of  $\widehat{\mathcal{F}}_2$  consisting of functions  $f \in \widehat{\mathcal{F}}_2$  for which both  $r_f$  and  $c_f$  are intensional.

Now, in general, reflectivity is not “preserved” by composition. However, it is proved that, for any  $f \in \mathcal{F}_2$ , there exists  $g \in \langle f \rangle \cap \widehat{\mathcal{F}}_2^\nabla$  for which  $r_g$  and  $c_g$  are reflective. Hence, every minimal clone generated by a function in  $\mathcal{F}_2$  contains a generator  $g \in \widehat{\mathcal{F}}_2^\nabla$  for which  $r_g$  and  $c_g$  are reflective.

## References

- [BM19] Behrisch, M. and Machida, H., An approach toward classification of minimal groupoids on a finite set, *Proc. 49th ISMVL*, IEEE, 2019, 164–169.
- [BM20] Behrisch, M. and Machida, H., On minimality of some binary clones related to unary functions, to appear in *Proc. 50th ISMVL*, IEEE, 2020.
- [Cs83] Csákány, B., All minimal clones on the three element set, *Acta Cybernet.*, **6**, 1983, 227–238.
- [Cs83a] Csákány, B., Three-element groupoids with minimal clones, *Acta Sci. Math.*, (Szeged) **45**, 1983, 111–117.
- [Po41] Post, E. L., The two-valued iterative systems of mathematical logic, *Ann. Math. Studies*, **5**, Princeton Univ. Press, 1941.
- [Ro86] Rosenberg, I. G., Minimal clones I: The five types, *Colloq. Math. Soc. J. Bolyai*, **43**, North Holland, 1986, 405–427.