

Ehrenfeucht theories approximated by countably categorical theories

Akito Tsuboi
University of Tsukuba*

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1 Preliminaries

Throughout, L denotes a countable language, and T is a complete theory formulated in L . L -formulas are denoted by φ, ψ , and so on. Variables are denoted by x, y, z and so on, and finite tuples of variables are denoted by $\bar{x}, \bar{y}, \bar{z}$ and so on. $I(\omega, T)$ denotes the number of countable models of T modulo isomorphism. $S(T)$ is the set of all complete types in T (over the empty set). ω is the set of natural numbers.

The most famous example of a theory with $1 < I(\omega, T) < \omega$ will be the following:

Example 1 (Ehrenfeucht). Let $L = \{<, 0, 1, 2, \dots\}$, and let T be the theory of the L -structure $(\mathbb{Q}, 0, 1, 2, \dots)$. Then, $I(\omega, T) = 3$, i.e., there are exactly three countable models of T : M_0 , the standard model; M_1 , the countable nonstandard model whose nonstandard part has a minimum element; and M_2 , the countable nonstandard model whose nonstandard part has no minimum element.

In this paper, a theory with $1 < I(\omega, T) < \omega$ is called an Ehrenfeucht theory. Lachlan's conjecture states that there is no stable Ehrenfeucht theory T . It was shown in [3] that by mimicking the method in Example 1, we cannot obtain a stable Ehrenfeucht theory. Indeed, it was shown that if $\{T_n : n \in \omega\}$ is an increasing sequence of \aleph_0 -categorical stable theories, then $T = \bigcup_{n \in \omega} T_n$ is non-Ehrenfeucht.

Definition 2. A type $p(\bar{x}) \in S(T)$ is called powerful, if whenever $M \models T$ realizes p , then M realizes all types in $S(T)$.

Below, we summarize the necessary facts for this paper, providing a brief explanation for each as to why it holds.

Fact 3. *Let T be an Ehrenfeucht theory.*

1. *T is small, i.e., $S(T)$ is countable. Suppose, for a contradiction, that $S(T)$ is uncountable. Choose any $p_0 \in S(T)$ and a countable model M_0 realizing p_0 . Since $S(T)$ is uncountable, there is $p_1 \in S(T)$ and a countable model M_1 realizing both p_0 and p_1 . Continuing this process, we get $p_0, \dots, p_n \in S(T)$ and countable models M_0, \dots, M_n such that M_n realizes p_0, \dots, p_n but does not realize p_{n+1} . Then, the M_n are mutually non-isomorphic, leading to a contradiction.*
2. *A powerful type exists. A similar argument as above shows that if there is no powerful type, then there would be infinitely many non-isomorphic countable models.*

Now we work in a countably saturated model \mathcal{M} of T , where T is Ehrenfeucht. For \bar{a} and A in \mathcal{M} , the type of \bar{a} over A is denoted by $\text{tp}(\bar{a}/A)$.

Fact 4. 1. *For all finite set $A \subset \mathcal{M}$, there is a prime and atomic model over A . This follows from the fact that $S(T)$ is countable.*

2. *Let $p(\bar{x}) \in S(T)$ be a powerful type. There exist realizations \bar{a} and \bar{a}' of p such that (i) $\text{tp}(\bar{a}'/\bar{a})$ is an isolated type, while (ii) $\text{tp}(\bar{a}/\bar{a}')$ is not. This can be shown as follows: Let \bar{a} be a realization of p . Choose a realization \bar{b} of p such that $\text{tp}(\bar{b}/\bar{a})$ is non-isolated. Let $q(\bar{x}, \bar{y}) = \text{tp}(\bar{a}, \bar{b})$. Since p is powerful, we can choose a realization (\bar{a}', \bar{b}') of q from a prime model $M_{\bar{a}}$ over \bar{a} . Then, it is easy to see that (\bar{a}, \bar{a}') satisfies the conditions (i) and (ii).*

2 Main Results

Continuing from the previous section, we work in \mathcal{M} .

Definition 5. We say $p(\bar{x}) = \text{tp}(\bar{a}/\bar{b})$ is semi-isolated, if there is a formula $\varphi(\bar{x}, \bar{b}) \in p$ such that $\forall \bar{x}(\varphi(\bar{x}, \bar{b}) \rightarrow \psi(\bar{x}))$ is true for all $\psi \in \text{tp}(\bar{a})$. If this situation holds, we also say $\varphi(\bar{x}, \bar{b})$ generates $p(\bar{x})$, and write $\varphi(\bar{x}, \bar{b}) \models p(\bar{x})$.

It is clear that an isolated type is semi-isolated, but the converse is not true in general. However, if $\text{tp}(\bar{b}/\bar{a})$ is isolated, then the following conditions are equivalent:

1. $\text{tp}(\bar{a}/\bar{b})$ is isolated;
2. $\text{tp}(\bar{a}/\bar{b})$ is semi-isolated.

This can be explained as follows. Since $\text{tp}(\bar{b}/\bar{a})$ is isolated, there is a formula $\varphi(\bar{x}, \bar{a})$ isolating the type.

Theorem 6. *Let $\{T_n\}_{n \in \omega}$ be an increasing sequence of \aleph_0 -categorical theories. Suppose that $T = \bigcup_{n \in \omega} T_n$ is an Ehrenfeucht theory. Then, there exist a definable order $<$ and elements e_n ($n \in \omega$) in T^{eq} such that (i) $e_n < e_{n+1}$ for all $n \in \omega$ and (ii) $<$ is a dense order without endpoints.*

Proof. We fix a powerful type $p \in S(T)$. For simplicity in notation, we assume that p is a 1-type, i.e., $p = p(x)$ where x is a single variable. Using Fact 4.2, choose realizations a and a' of p such that $\text{tp}(a'/a)$ is isolated and $\text{tp}(a/a')$ is not isolated. Let $r(x, y) = \text{tp}(a', a)$ and choose a formula $\psi(x, y) \in r(x, y)$ witnessing that $\text{tp}(a'/a)$ is isolated. Then, $\psi(x, y)$ satisfies the following property:

- (*) For all d realizing p , $\psi(x, d)$ generates $p(x)$, and $\psi(d, y)$ does not generate $p(y)$.

Observe that the formulas satisfying (*) are closed under a finite disjunction. Choose $m \in \omega$ such that all symbols in $\psi(x, y)$ belong to $L(T_m)$, which is the language of T_m . Since T_m is \aleph_0 -categorical, there are only a finite number of formulas satisfying (*). Thus, we can assume $\psi(x, y) \in r(x, y)$ is the weakest formula possessing the property (*).

Claim A. *For $n \in \omega$, let $\psi^n(x, y)$ be the formula $\exists z_1, \dots, z_n [\psi(x, z_1) \wedge \bigwedge_{i=1, \dots, n-1} \psi(z_i, z_{i+1}) \wedge \psi(z_n, y)]$. Then, each formula $\psi^n(x, y)$ satisfies the condition (*) above.*

We assume $n = 2$ for simplicity since other cases are treated similarly. It is clear that $\psi^2(x, d) \models p(x)$ holds for all d realizing p . Suppose, for a contradiction, that $\psi^2(d, y) \models p(y)$. Choose d_1 and d_2 such that both (d, d_1) and (d_1, d_2) realize $r(x, y)$. Since $\psi(x, y)$ belongs to r , $\psi^2(d, d_2)$ holds. So, $\text{tp}(d_2/d)$ is semi-isolated. As $\text{tp}(d_1/d_2)$ is also semi-isolated, we must have that $\text{tp}(d_1/d)$ is semi-isolated. By the remark just after Definition 5, $\text{tp}(d_1/d)$ must be isolated. A contradiction. (End of Proof of Claim A)

Claim B. *Let (d, d_1) and (d_1, d_2) realize $r(x, y)$. Then, $\psi(\mathcal{M}, d)$ is a proper subset of $\psi(\mathcal{M}, d_2)$.*

By Claim A, $\psi^3(x, y)$ satisfies (*), so $\psi^3(x, d)$ proves $\psi(x, d)$. Also, by the choice of d, d_1, d_2 , we see that $\psi(x, d_2)$ proves $\psi^3(x, d)$. Thus, $\psi(\mathcal{M}, d_2)$ is a subset of $\psi(\mathcal{M}, d)$. If the inclusion is not proper, since $\psi(d_1, d_2)$ holds, we must have $\psi(d_1, d)$. This is impossible, since $\text{tp}(d_1/d)$ is not semi-isolated. (End of Proof of Claim B)

Using Claim B, we can choose a sequence $\{a_n\}_{n \in \omega}$ of realizations of p such that $\{\psi(\mathcal{M}, a_n)\}_{n \in \omega}$ forms a strictly increasing sequence of uniformly defined definable sets. Thus, T has the strict order property.

Now we define formulas:

$$\begin{aligned} y \succeq z &\equiv \forall x(\psi(x, y) \rightarrow \psi(x, z)), \\ y \sim z &\equiv y \succeq z \succeq y, \\ y \succ z &\equiv y \succeq z \wedge \neg z \succeq y. \end{aligned}$$

Clearly, \succ defines a pre-order without endpoints.

Claim C. *There is a dense pre-order without end points.*

For all $n \in \omega$, let $y \succ^n z$ be the formula $\exists x_1, \dots, x_n(y \succ x_1 \succ \dots \succ x_n \succ z)$. These formulas are all in $L(T_m)$. Since T_m is \aleph_0 -categorical, there is $k \in \omega$ such that \succ^n are equivalent for all numbers $n \geq k$. Then, \succ^k defines a dense pre-order. Indeed, if $b \succ^k d$, then $b \succ^{2k+1} d$. So, there exists e with $b \succ^k e \succ^k d$. (End of Proof of Claim C)

In the quotient structure by \sim , the order induced by \succ satisfies our requirement. \square

Question 7. The language of Example 1 is infinite. However, there exist Ehrenfeucht theories whose languages are finite. For instance, consider the language $L = \{\preceq, \sim\}$ and the L -theory T axiomatized by the following sentences:

1. \preceq is a dense linear preorder without endpoints;
2. $x \sim y$ if and only if $x \preceq y \preceq x$;
3. For every positive natural number n , there exists a unique \sim -equivalence class, denoted c_n , containing exactly n elements. Moreover, these equivalence classes form a strictly increasing chain: $c_1 \prec c_2 \prec \dots$.

Can we construct a theorem demonstrating that mimicking this example does not yield a stable Ehrenfeucht theory?

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Professor Emeritus at the University of Tsukuba,
E-mail address: tsuboi@math.tsukuba.ac.jp